

Coinduction - An Introductory Example

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based on: Jacobs and Rutten, A Tutorial on (Co)Algebras and (Co)Induction. EATCS Bulletin 62, 1997, p.222-259

coinduction

How to reason about infinite data/processes? An example:

Given $\text{zip} : A^{\mathbb{N}} \times A^{\mathbb{N}} \rightarrow A^{\mathbb{N}}$, $\text{even} : A^{\mathbb{N}} \rightarrow A^{\mathbb{N}}$, $\text{odd} : A^{\mathbb{N}} \rightarrow A^{\mathbb{N}}$

defined by

$$\text{head}(\text{zip}(l_1, l_2)) = \text{head}(l_1)$$

$$\text{tail}(\text{zip}(l_1, l_2)) = \text{zip}(\text{tail}(l_2), \text{tail}(l_1))$$

$$\text{head}(\text{even}(l)) = \text{head}(l)$$

$$\text{tail}(\text{even}(l)) = \text{even}(\text{tail}(\text{tail}(l)))$$

$$\text{odd}(l) = \text{even}(\text{tail}(l))$$

show that

$$\text{zip}(\text{even}(x), \text{odd}(x)) = x$$

example, cont'd

Given

$$\begin{aligned} \text{head}(\text{zip}(l_1, l_2)) &= \text{head}(l_1) \\ \text{tail}(\text{zip}(l_1, l_2)) &= \text{zip}(l_2, \text{tail}(l_1)) \end{aligned}$$

$$\begin{aligned} \text{head}(\text{even}(l)) &= \text{head}(l) \\ \text{tail}(\text{even}(l)) &= \text{even}(\text{tail}(\text{tail}(l))) \\ \text{odd}(l) &= \text{even}(\text{tail}(l)) \end{aligned}$$

show

$$\text{zip}(\text{even}(x), \text{odd}(x)) = x$$

$$\text{head}(\text{zip}(\text{even}(x), \text{odd}(x))) = \text{head}(\text{even}(x)) = \text{head}(x)$$

example, cont'd

Given

$$\text{head}(\text{zip}(l_1, l_2)) = \text{head}(l_1)$$

$$\text{tail}(\text{zip}(l_1, l_2)) = \text{zip}(l_2, \text{tail}(l_1))$$

$$\text{head}(\text{even}(l)) = \text{head}(l)$$

$$\text{tail}(\text{even}(l)) = \text{even}(\text{tail}(\text{tail}(l)))$$

$$\text{odd}(l) = \text{even}(\text{tail}(l))$$

show

$$\text{zip}(\text{even}(x), \text{odd}(x)) = x$$

$$\text{head}(\text{zip}(\text{even}(x), \text{odd}(x))) = \text{head}(\text{even}(x)) = \text{head}(x)$$

example, cont'd

Given

$$\text{head}(\text{zip}(l_1, l_2)) = \text{head}(l_1)$$

$$\text{tail}(\text{zip}(l_1, l_2)) = \text{zip}(l_2, \text{tail}(l_1))$$

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$$\text{odd}(l) = \text{even}(\text{tail}(l))$$

show

$$\text{zip}(\text{even}(x), \text{odd}(x)) = x$$

$$\text{head}(\text{zip}(\text{even}(x), \text{odd}(x))) = \text{head}(\text{even}(x)) = \text{head}(x)$$

example, cont'd

$$\text{head}(\text{zip}(l_1, l_2)) = \text{head}(l_1)$$

$$\text{tail}(\text{zip}(l_1, l_2)) = \text{zip}(l_2, \text{tail}(l_1))$$

$$\text{head}(\text{even}(l)) = \text{head}(l)$$

$$\text{tail}(\text{even}(l)) = \text{even}(\text{tail}(\text{tail}(l)))$$

$$\text{odd}(l) = \text{even}(\text{tail}(l))$$

show

$$\text{zip}(\text{even}(x), \text{odd}(x)) = x$$

$$\begin{aligned}\text{tail}(\text{zip}(\text{even}(x), \text{odd}(x))) &= \text{zip}(\text{odd}(x), \text{tail}(\text{even}(x))) \\ &= \text{zip}(\text{odd}(x), \text{even}(\text{tail}(\text{tail}(x)))) \\ &= \text{zip}(\text{even}(\text{tail}(x)), \text{odd}(\text{tail}(x))) = \text{tail}(x)\end{aligned}$$

example, cont'd

$$\text{head}(\text{zip}(l_1, l_2)) = \text{head}(l_1)$$

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show

$$\text{zip}(\text{even}(x), \text{odd}(x)) = x$$

$$\begin{aligned}\text{tail}(\text{zip}(\text{even}(x), \text{odd}(x))) &= \text{zip}(\text{odd}(x), \text{tail}(\text{even}(x))) \\ &= \text{zip}(\text{odd}(x), \text{even}(\text{tail}(\text{tail}(x)))) \\ &= \text{zip}(\text{even}(\text{tail}(x)), \text{odd}(\text{tail}(x))) = \text{tail}(x)\end{aligned}$$

example, cont'd

$$\text{head}(\text{zip}(l_1, l_2)) = \text{head}(l_1)$$

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$$\text{odd}(l) = \text{even}(\text{tail}(l))$$

show

$$\text{zip}(\text{even}(x), \text{odd}(x)) = x$$

$$\begin{aligned}\text{tail}(\text{zip}(\text{even}(x), \text{odd}(x))) &= \text{zip}(\text{odd}(x), \text{tail}(\text{even}(x))) \\ &\stackrel{\textcolor{red}{=}}{=} \text{zip}(\text{odd}(x), \text{even}(\text{tail}(\text{tail}(x)))) \\ &= \text{zip}(\text{even}(\text{tail}(x)), \text{odd}(\text{tail}(x))) = \text{tail}(x)\end{aligned}$$

example, cont'd

$$\text{head}(\text{zip}(l_1, l_2)) = \text{head}(l_1)$$

$$\text{tail}(\text{zip}(l_1, l_2)) = \text{zip}(l_2, \text{tail}(l_1))$$

$$\text{head}(\text{even}(l)) = \text{head}(l)$$

$$\text{tail}(\text{even}(l)) = \text{even}(\text{tail}(\text{tail}(l)))$$

$$\textcolor{red}{\text{odd}(l) = \text{even}(\text{tail}(l))}$$

show

$$\text{zip}(\text{even}(x), \text{odd}(x)) = x$$

$$\begin{aligned}\text{tail}(\text{zip}(\text{even}(x), \text{odd}(x))) &= \text{zip}(\text{odd}(x), \text{tail}(\text{even}(x))) \\ &= \text{zip}(\textcolor{red}{\text{odd}(x)}, \textcolor{purple}{\text{even}(\text{tail}(\text{tail}(x)))}) \\ &\textcolor{red}{=} \text{zip}(\textcolor{red}{\text{even}(\text{tail}(x))}, \textcolor{purple}{\text{odd}(\text{tail}(x))}) = \text{tail}(x)\end{aligned}$$

example, cont'd

$$\text{head}(\text{zip}(l_1, l_2)) = \text{head}(l_1)$$

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show

$$\text{zip}(\text{even}(x), \text{odd}(x)) = x$$

$$\begin{aligned}\text{tail}(\text{zip}(\text{even}(x), \text{odd}(x))) &= \text{zip}(\text{odd}(x), \text{tail}(\text{even}(x))) \\ &= \text{zip}(\text{odd}(x), \text{even}(\text{tail}(\text{tail}(x)))) \\ &= \text{zip}(\text{even}(\text{tail}(x)), \text{odd}(\text{tail}(x))) = \text{tail}(x)\end{aligned}$$

explanation (bisimulation)

Coinduction Proof Principle: Two streams $x, x' \in A^{\mathbb{N}}$ are equal iff there is a relation R with xRx' and for all y, y'

$$\begin{aligned} yRy' &\Rightarrow \text{head}(y) = \text{head}(y') \\ yRy' &\Rightarrow \text{tail}(y) R \text{tail}(y') \end{aligned}$$

Remark: A relation R satisfying the two conditions above is called a stream-bisimulation.

Example: Put $\text{zip}(\text{even}(x), \text{odd}(x))) R x$ for all $x \in A^{\mathbb{N}}$

$$\begin{aligned} \text{tail}(\text{zip}(\text{even}(x), \text{odd}(x))) &= \text{zip}(\text{odd}(x), \text{tail}(\text{even}(x))) \\ &= \text{zip}(\text{odd}(x), \text{even}(\text{tail}(\text{tail}(x)))) \\ &= \text{zip}(\text{even}(\text{tail}(x)), \text{odd}(\text{tail}(x))) R \text{tail}(x) \end{aligned}$$

explanation (coinduction via finality)

Def: An object Z is **final** if for all objects X there is a unique arrow $X \rightarrow Z$.

Observation: $A^{\mathbb{N}} \rightarrow A \times A^{\mathbb{N}}$ is the final coalgebra (for the functor $TX = A \times X$).

Fact: To say that $X \rightarrow A \times X$ satisfies the coinduction proof principle is equivalent to saying that $X \rightarrow A \times X$ is the final coalgebra.