# Synchronised Hyperedge Replacement as a Model for Service Oriented Computing

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International PhD School in Theory and Practice of
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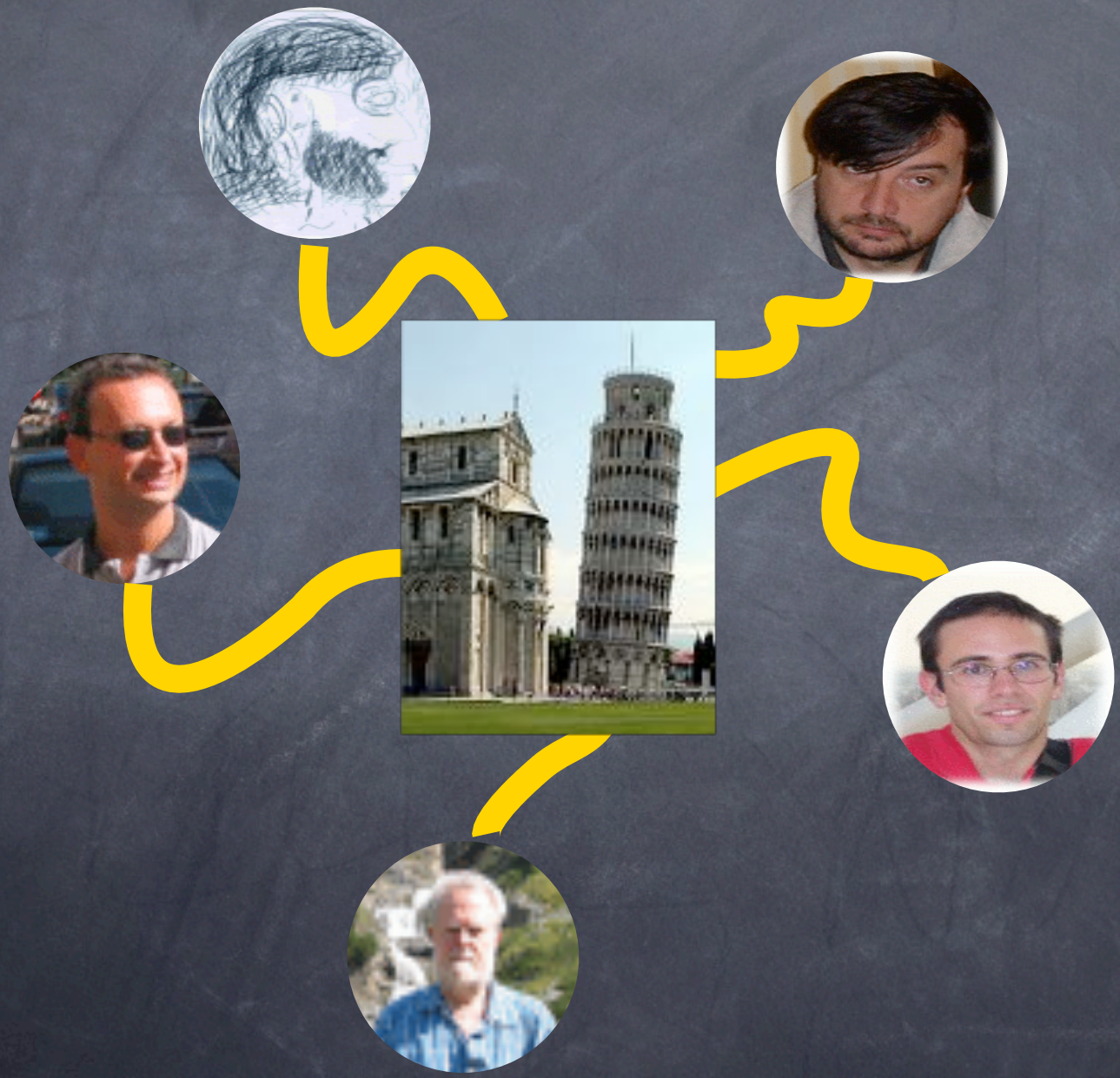
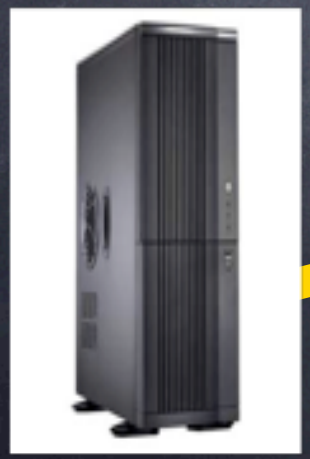
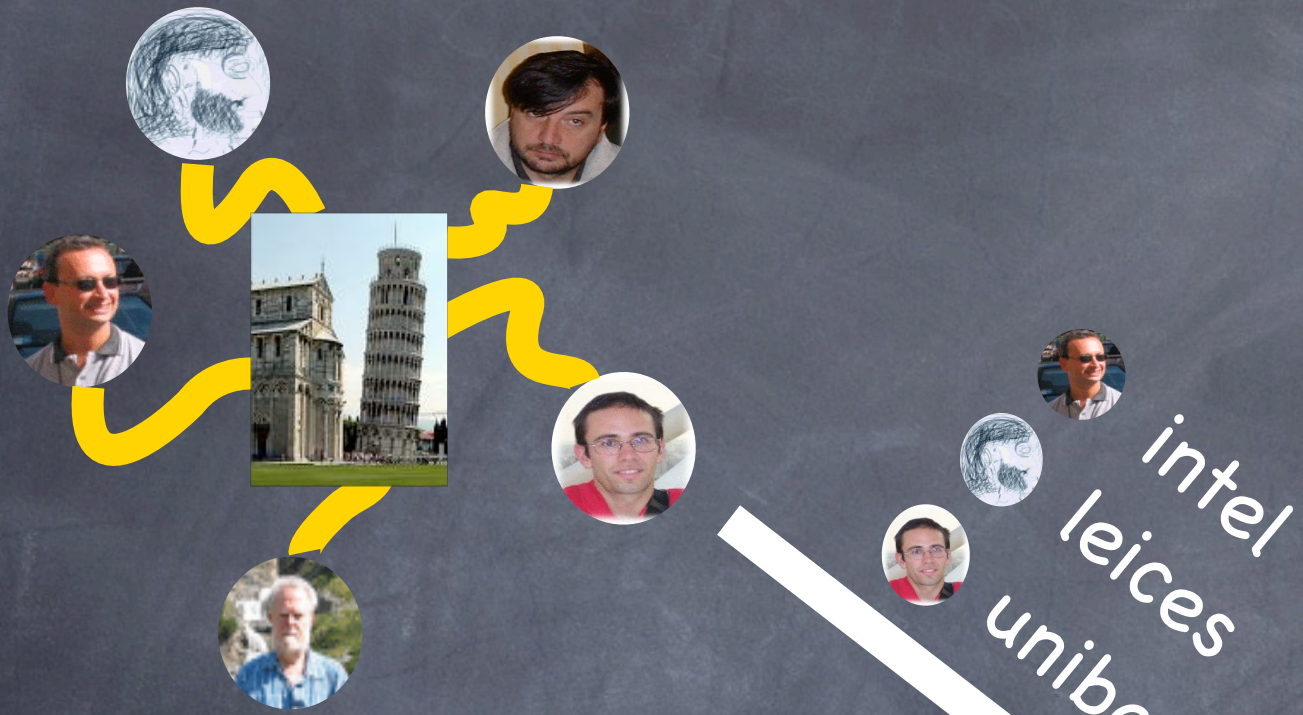
# SHR framework

- SHR as a uniform framework for (non-)functional aspects of SOC
  - Context-free flavour
  - "SOC systems as Hypergraphs" & "SOC computations as SHR"
    - Components = hyperedges
    - Systems = bunches of hyperedges
    - Computing = rewrite hypergraphs...(distributed constraint solving)
    - ...using "some" (parameterisable) synchronisation policy



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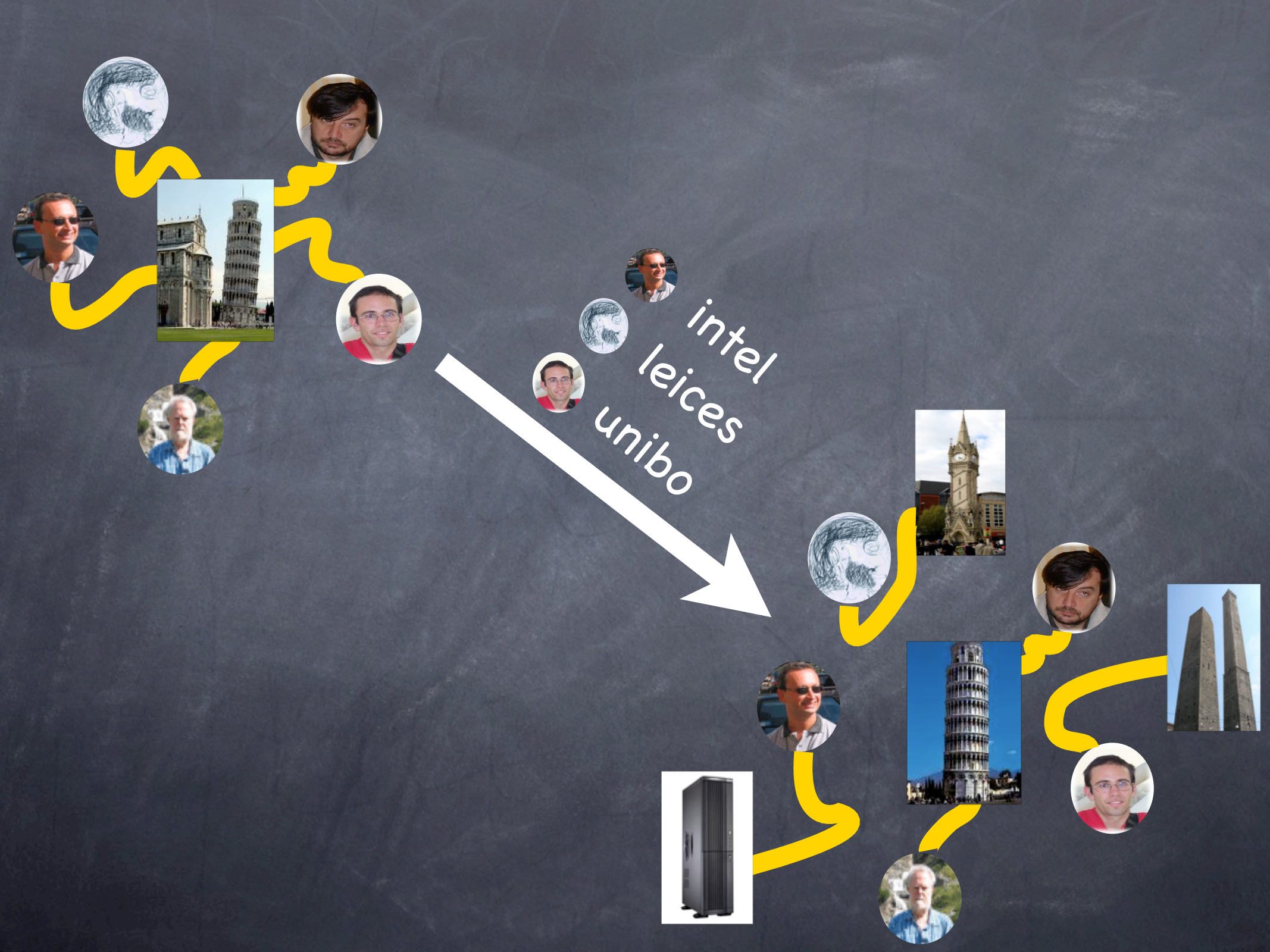
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# Models

- Process calculi
  - CSP, CCS and π-calculus...
- Graph-based models
  - Synchronised Hyperedge Replacement (SHR)
  - Originally, SHR as a model of distributed systems and software architectures but
  - expressive enough to model many process calculi
- ...

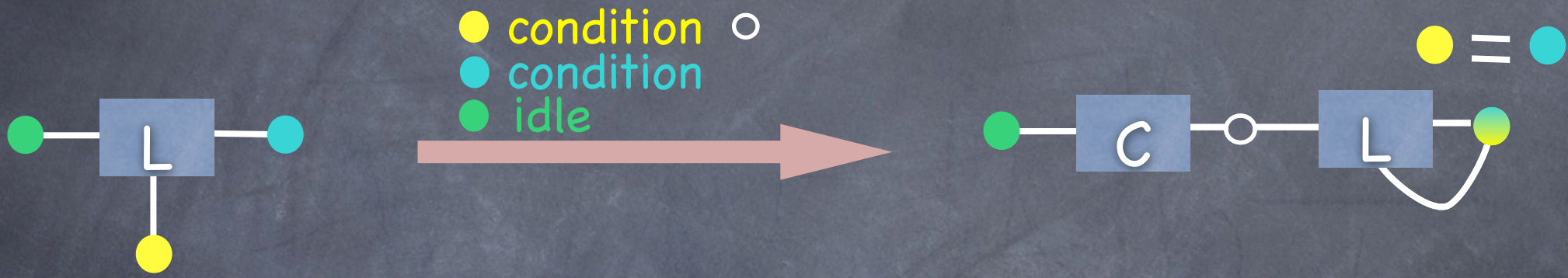
# SHR features

- can express many forms of synchronisation
- constraint satisfaction guide rewriting by synchronising "context-free productions"
  - components' behaviour independently specified by productions
  - productions impose conditions on adjacent nodes
  - global transitions as application of "compatible" productions
- QoS mechanism for driving the rewritings

# SOC

- Modern distributed systems
  - complex and heterogeneous
  - many architectural levels
  - many communication infrastructures
  - geographically distributed
  - highly dynamic
- SOC as modelling paradigm
  - Services are
    - independently specified/ published
    - searched/discovered and dynamically assembled
    - dynamically reconfiguration
    - mobile and requiring complex synchronisations
  - "QoS aware"

# SHR rewriting: in a nutshell



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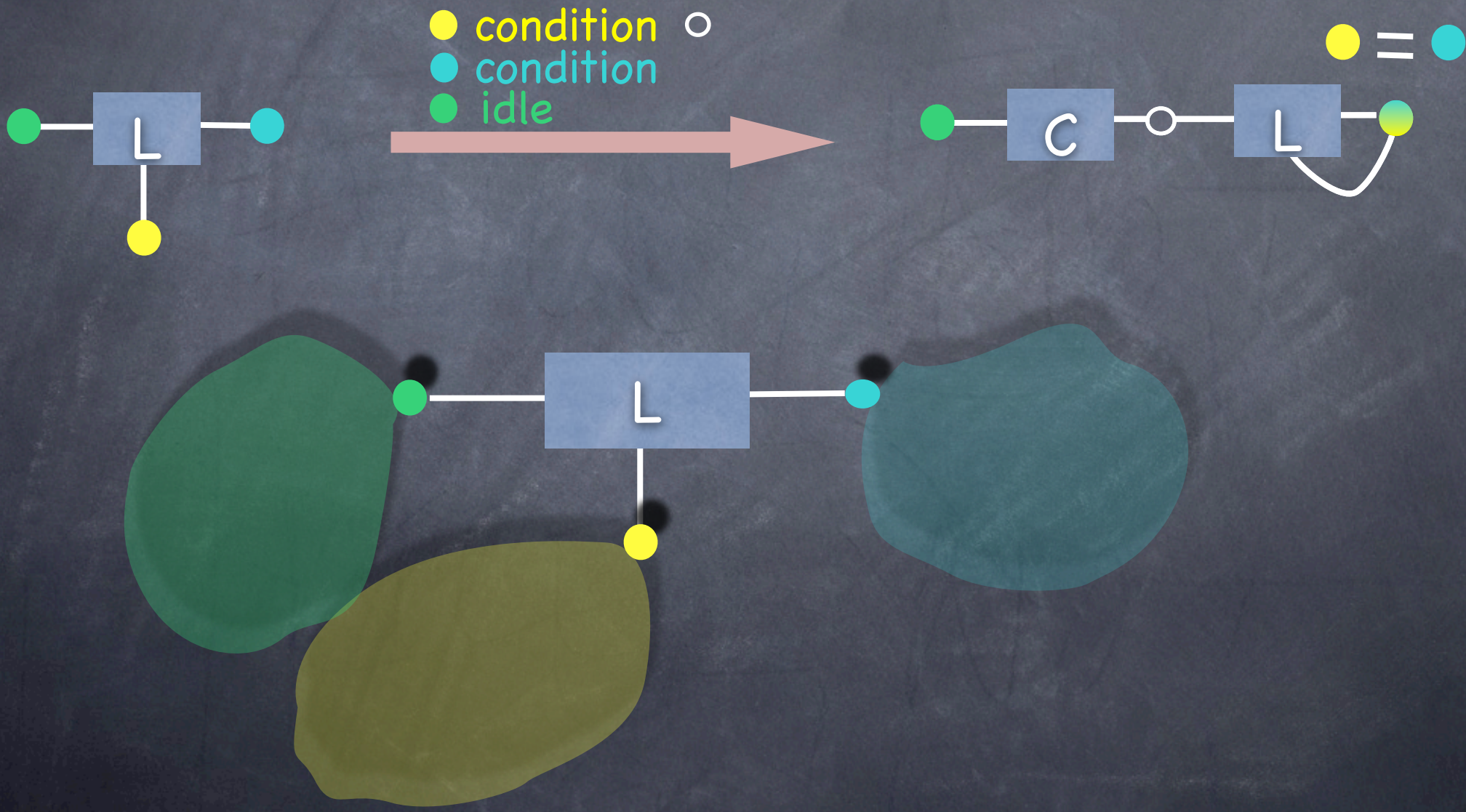

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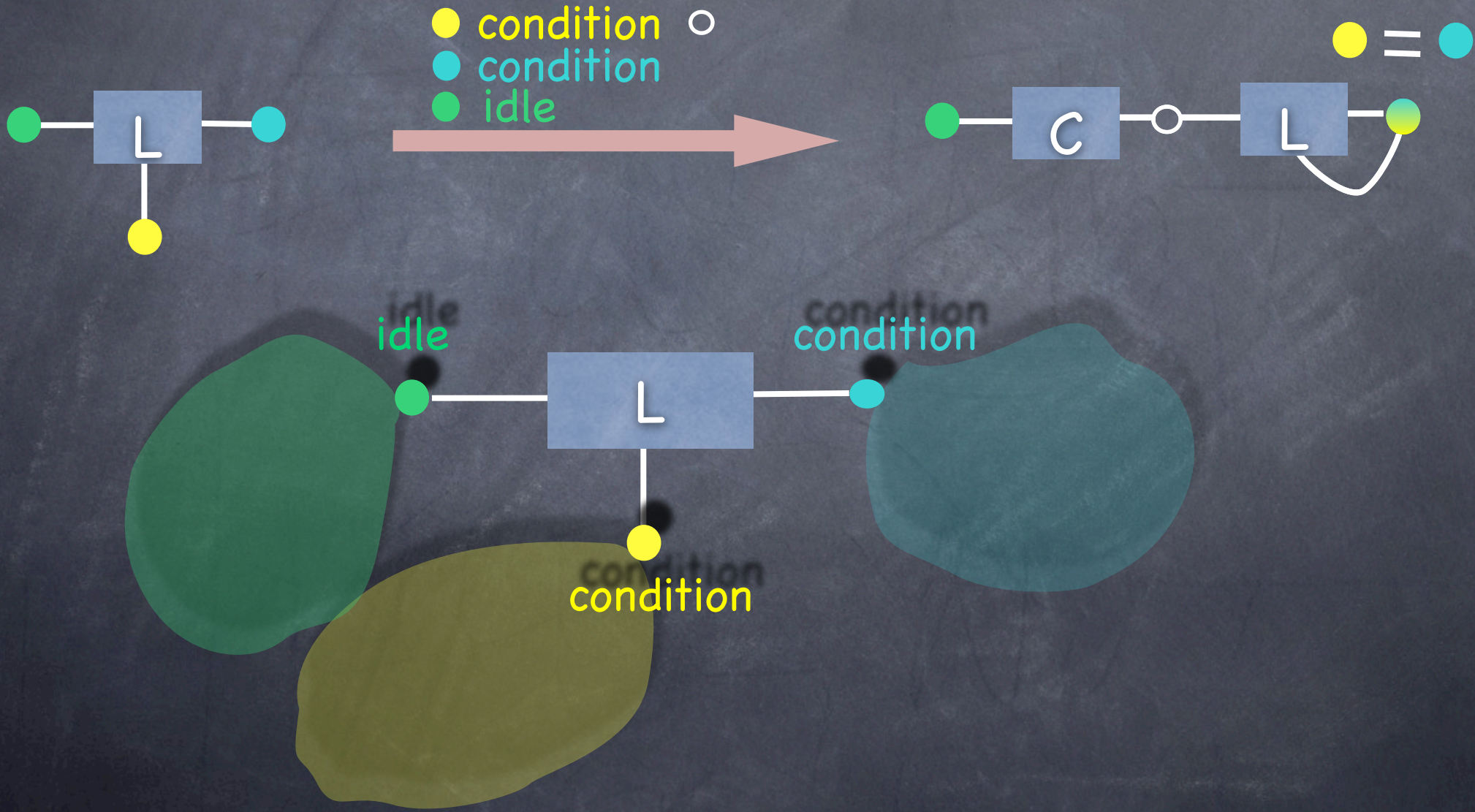

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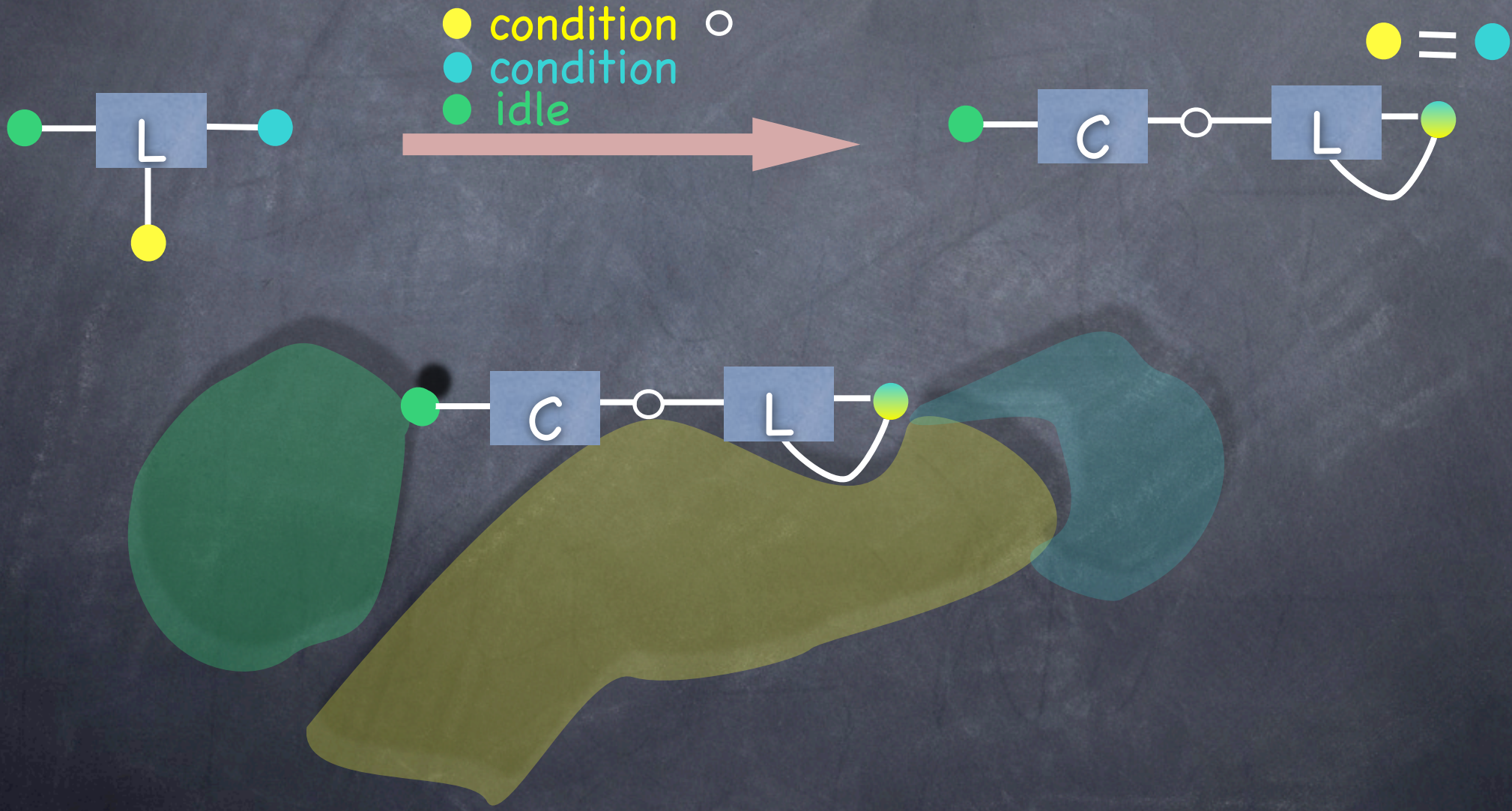

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  - application level QoS
  - sophisticated synchronisations

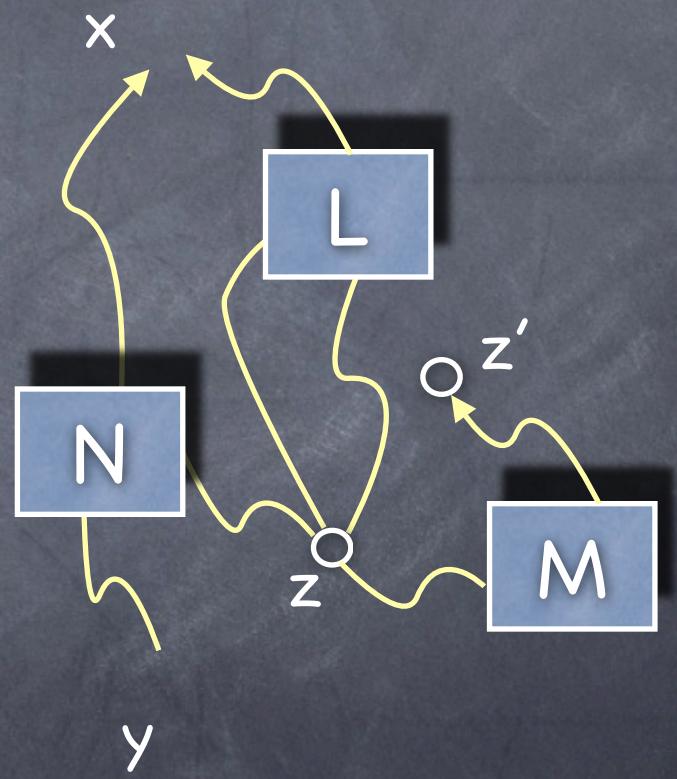
# Plan

- Give the basic definitions for SHR
- Analise 2 specific cases:
  - Milner synchronisation (with(out) mobility)
  - SHReQ
- ADR (if time allows)

# Hypergraphs

Syntax

# Exercises



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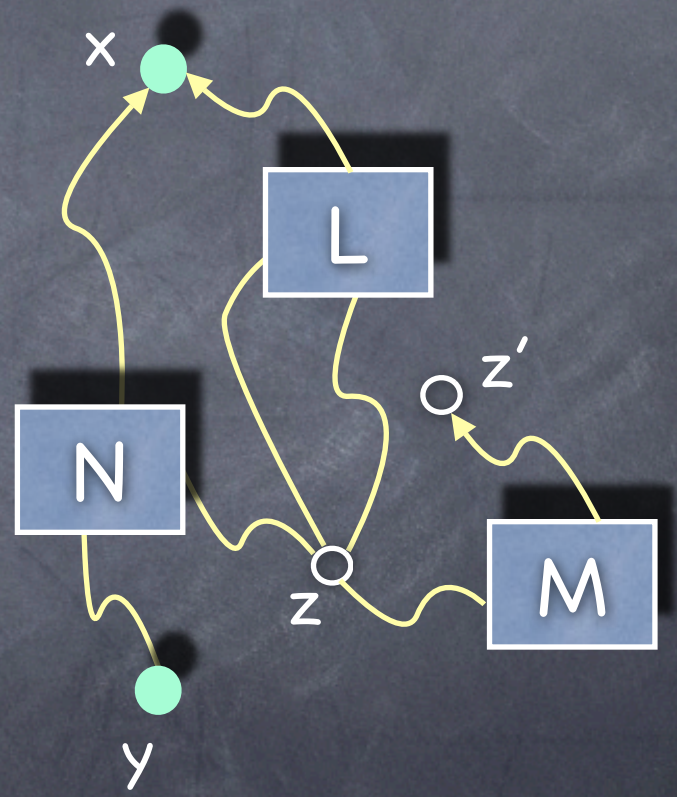

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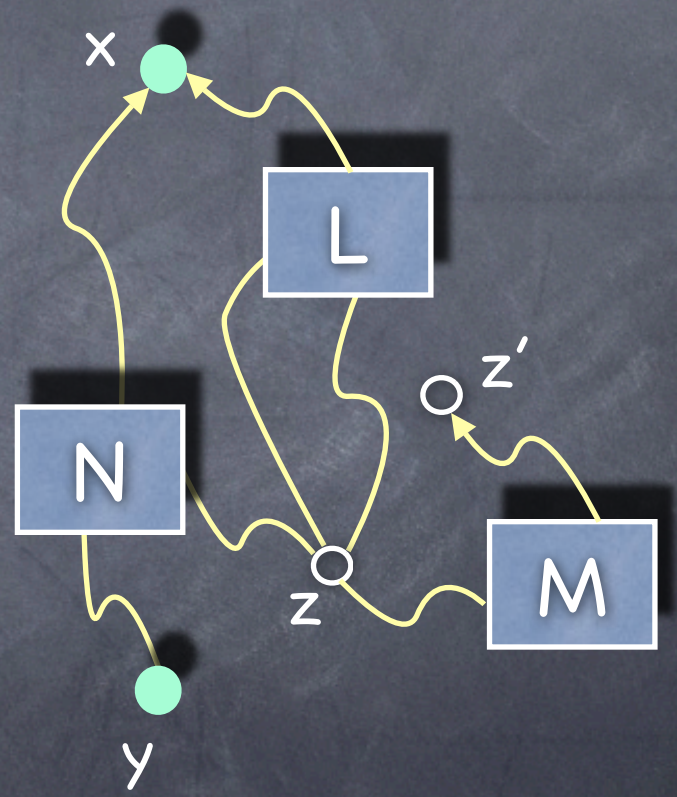

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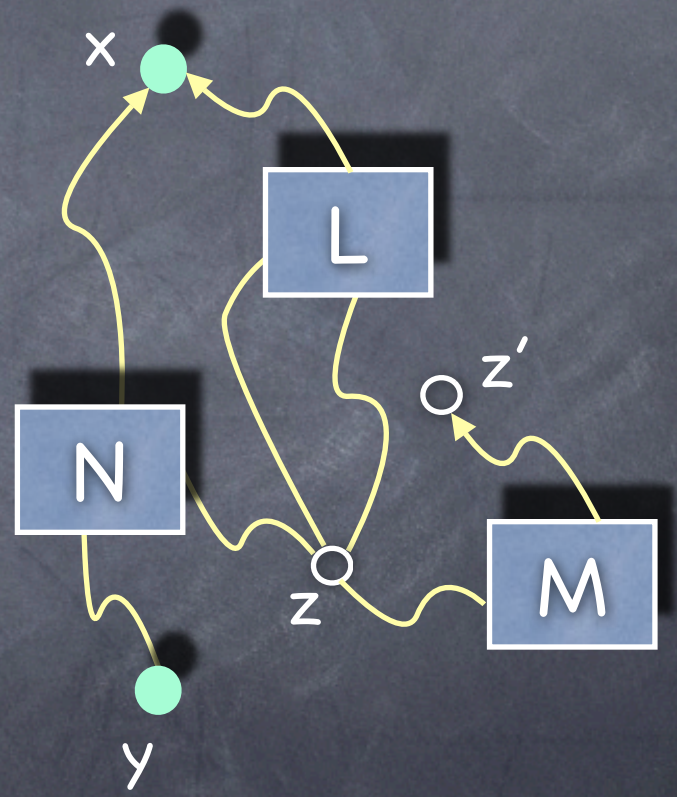

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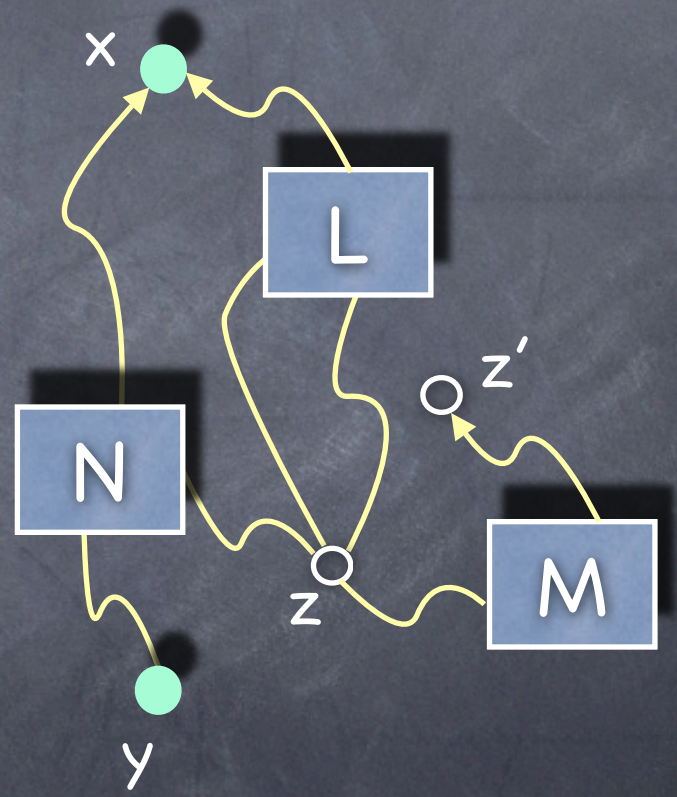

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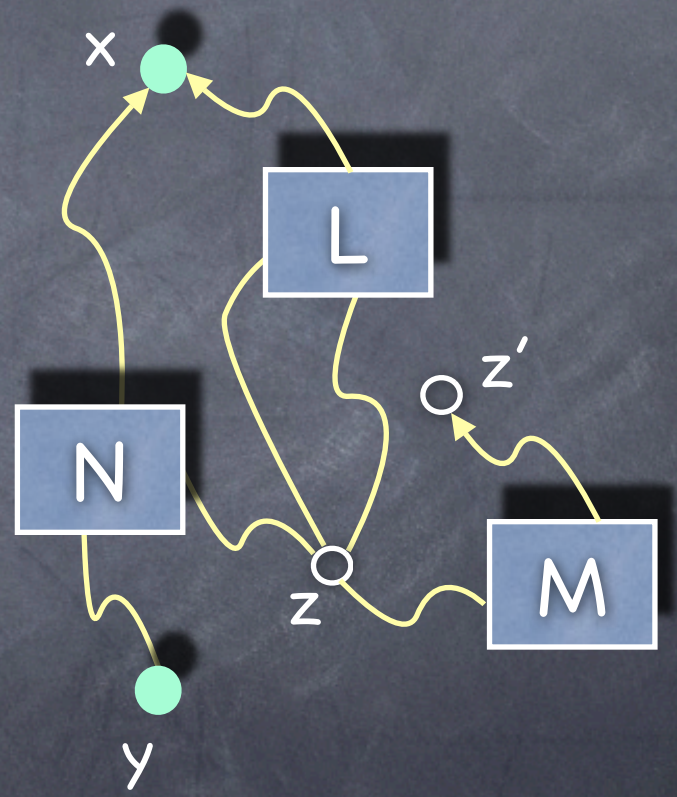
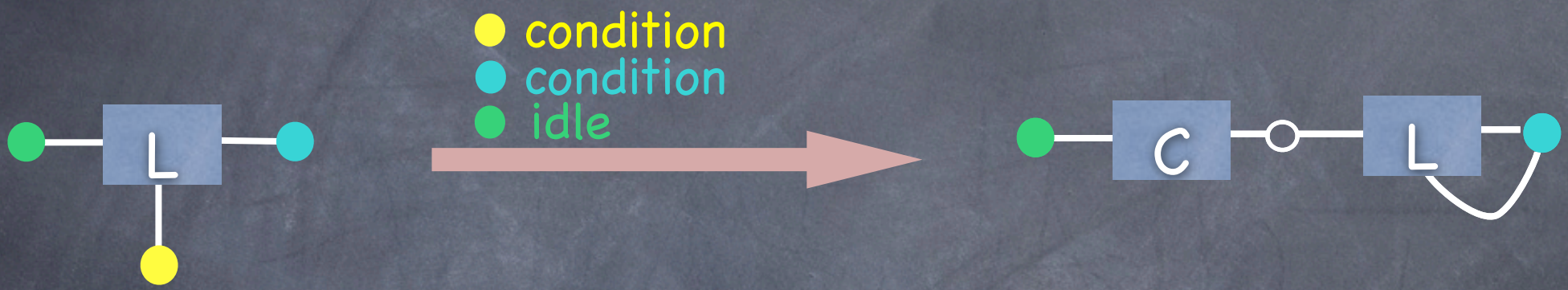
  - $x,y \vdash L(x,y) \mid L(x,z)$



# The simplest SHR: Basic Milner SHR

"Milner" synchronisation
without mobility [fhlmt05]

# bMSHR: in a nutshell



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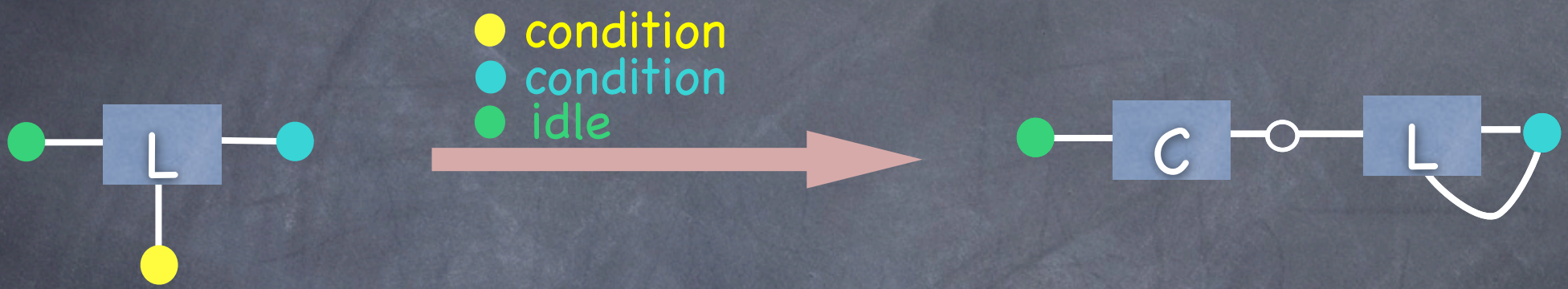


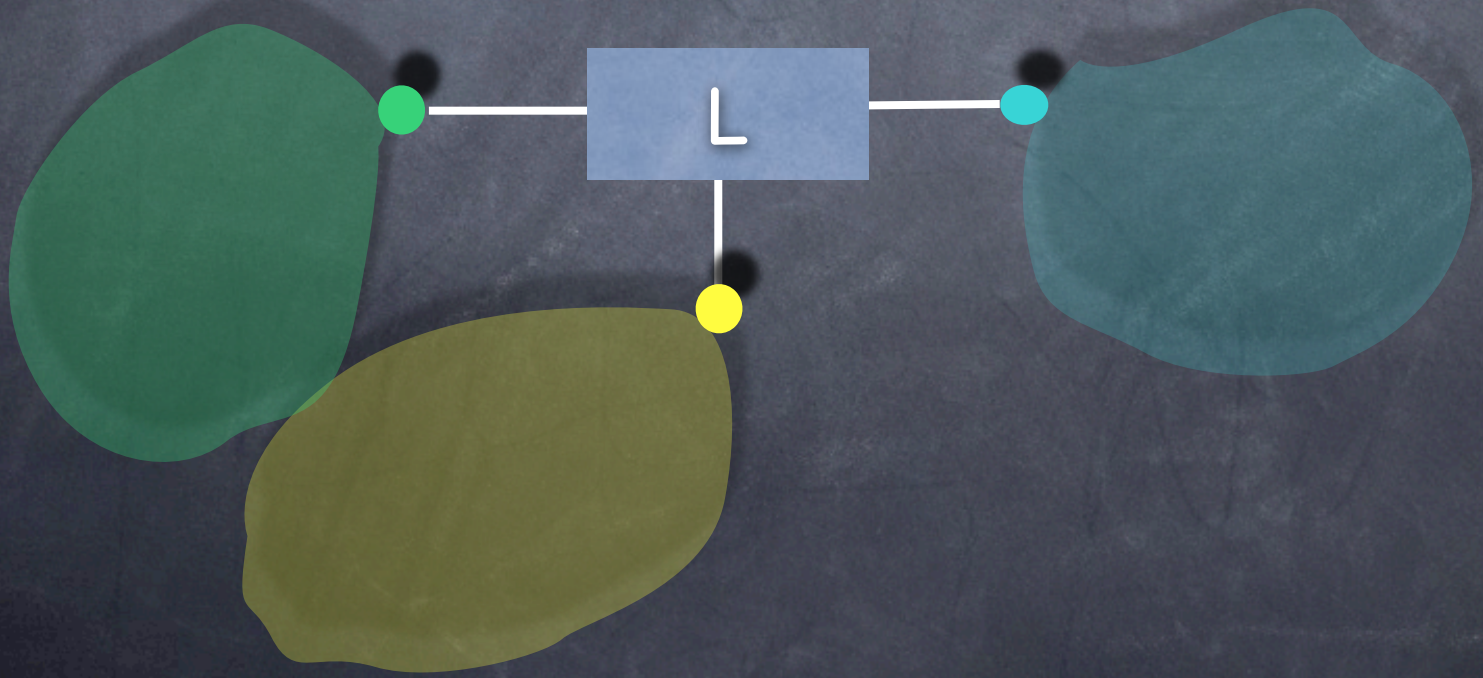

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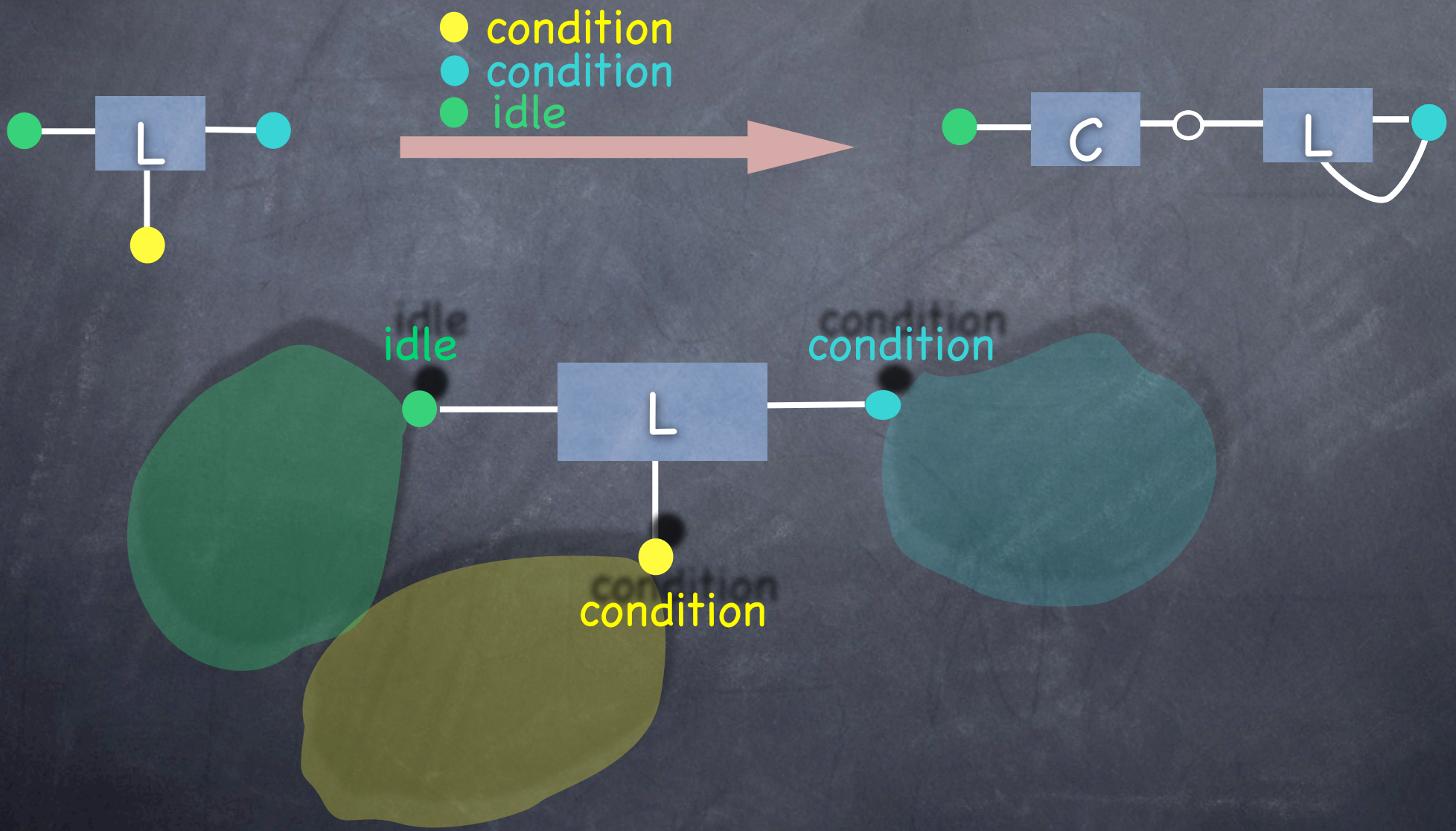

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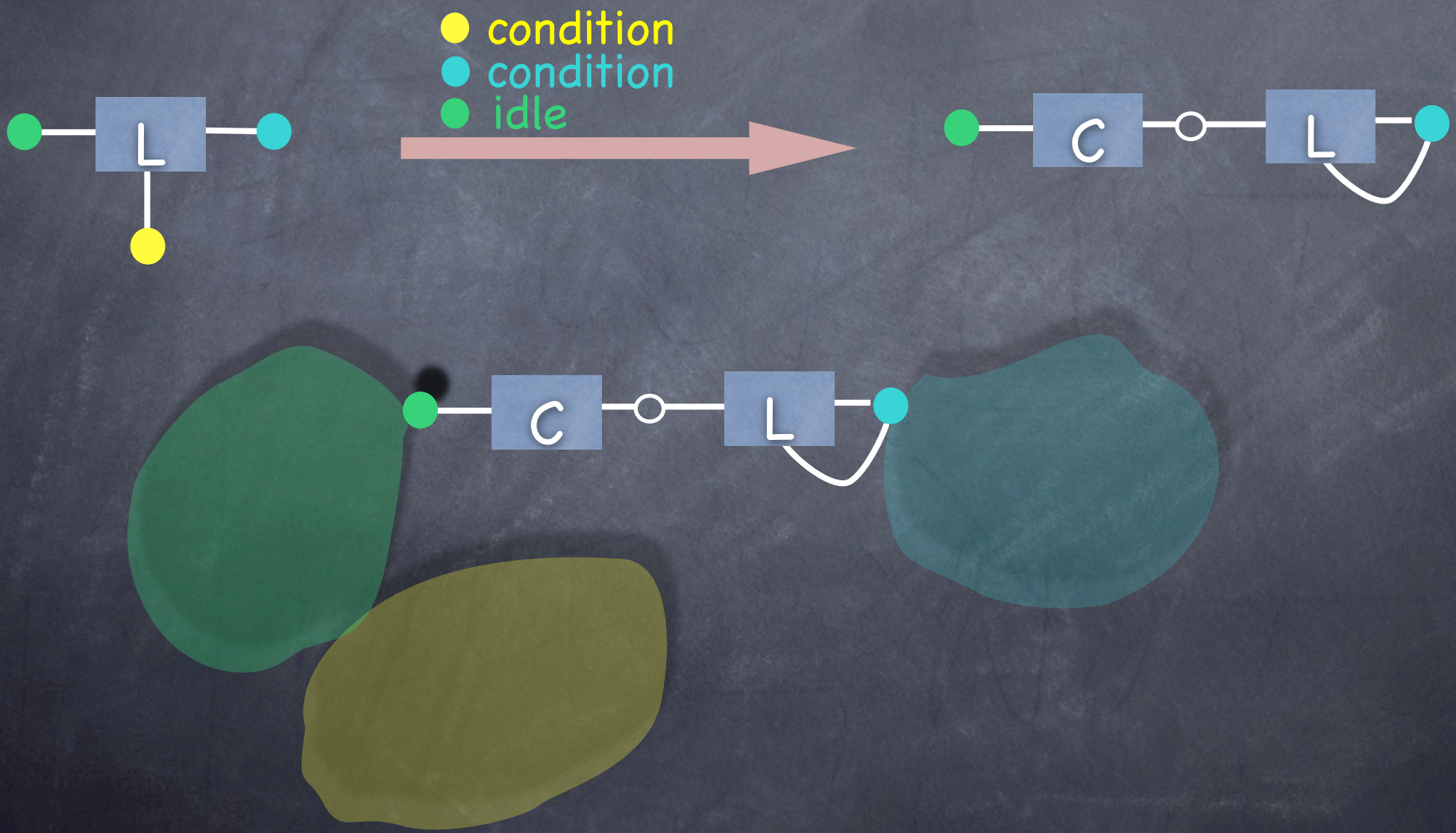
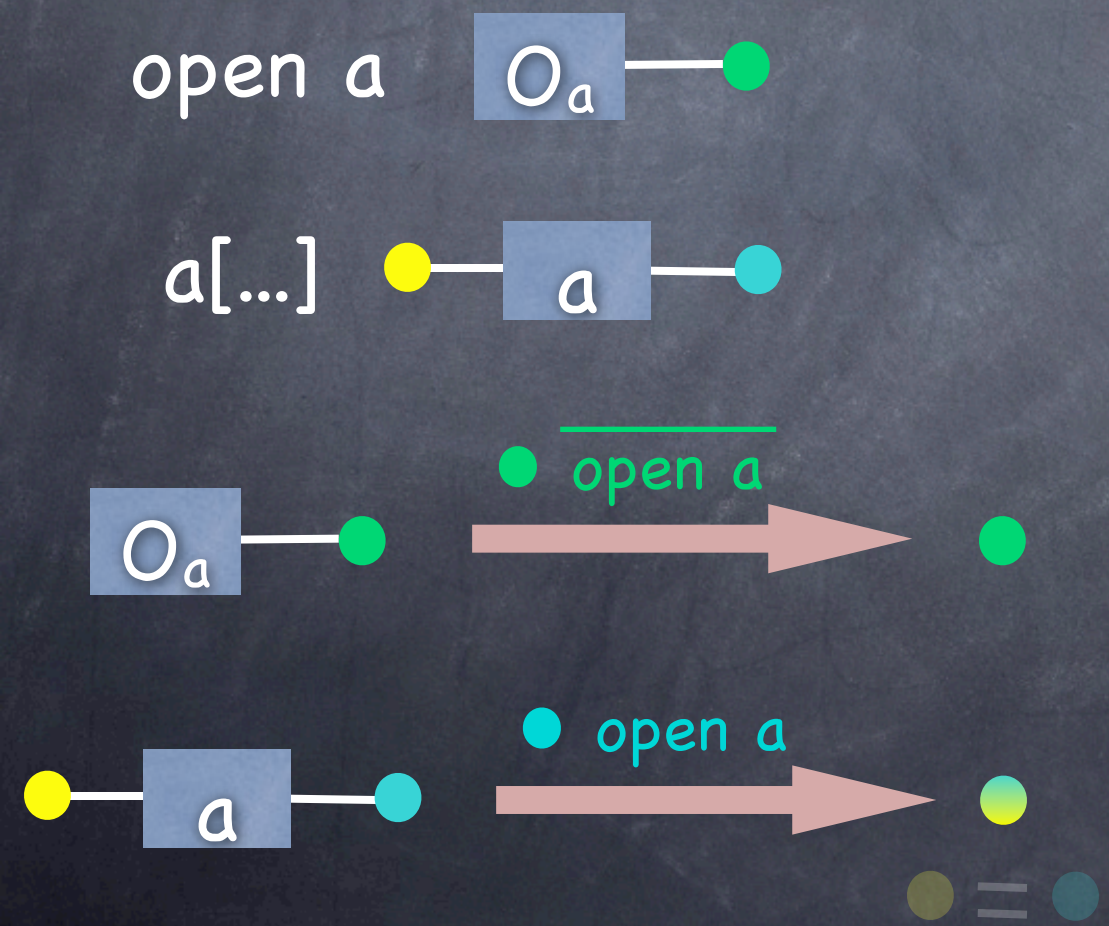


# SHR & mobility

"Milner" synchronisation
mobility [fmt01]

# Quest for mobility...

In Ambient: open a | a[...] → ...



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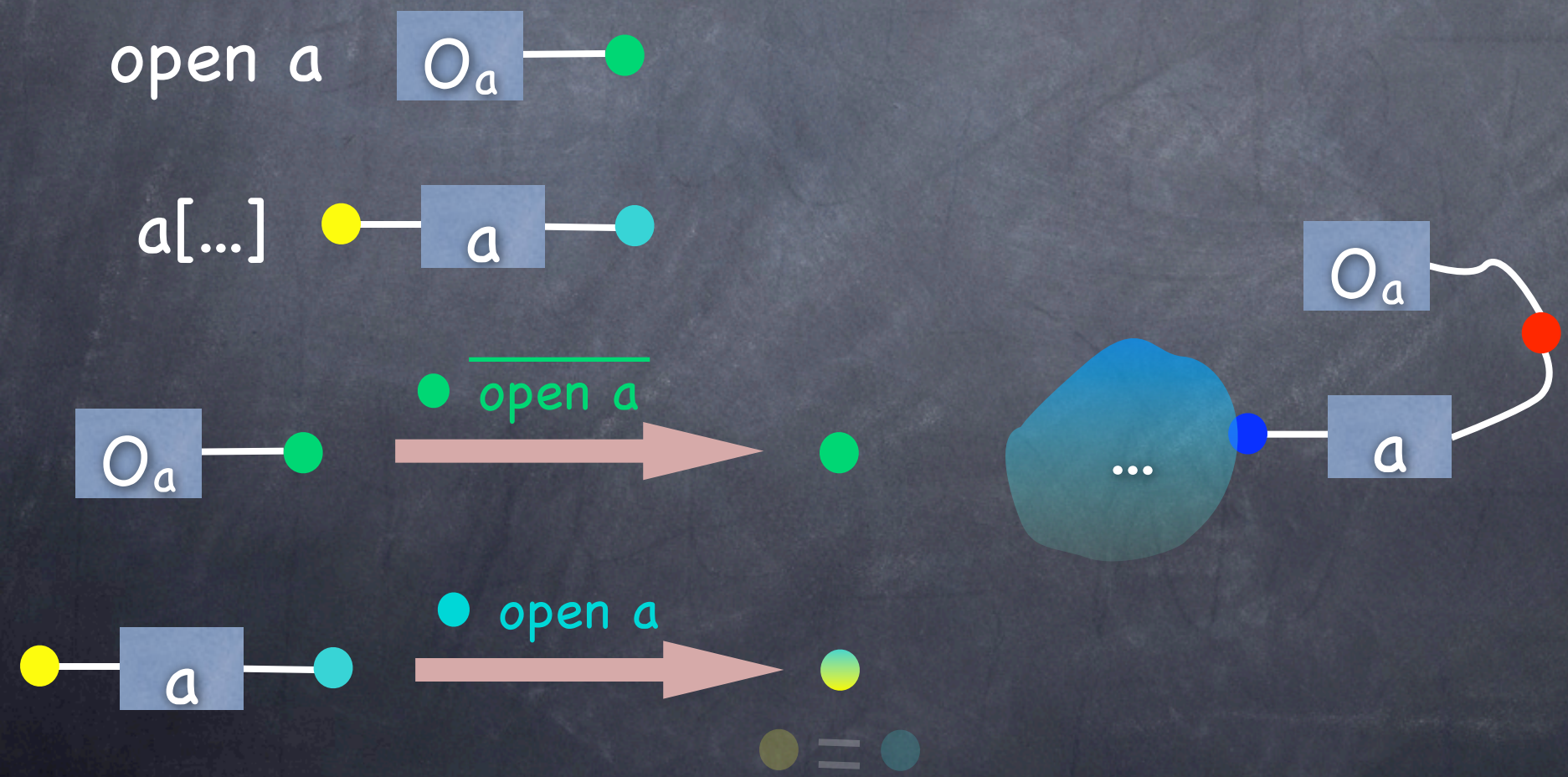

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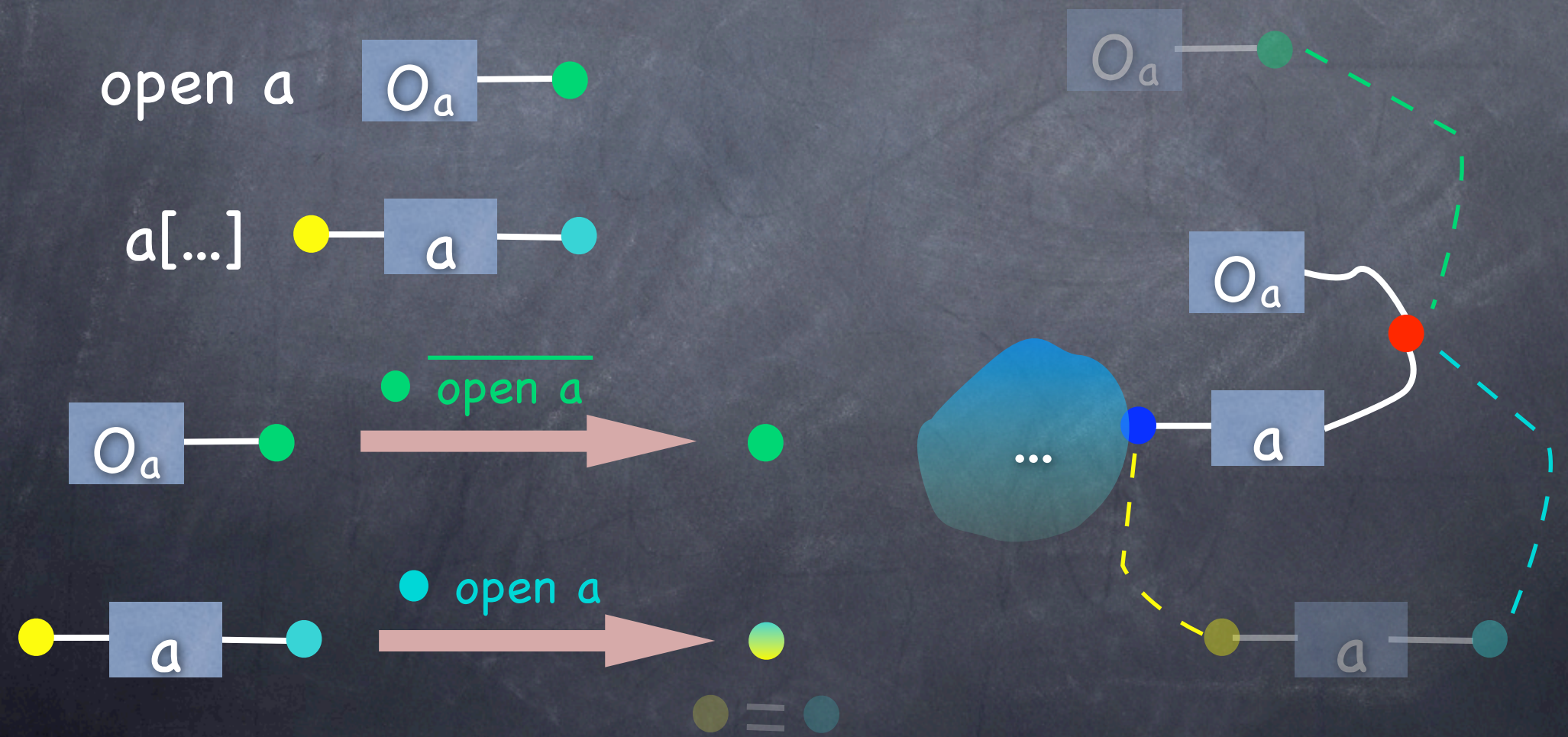

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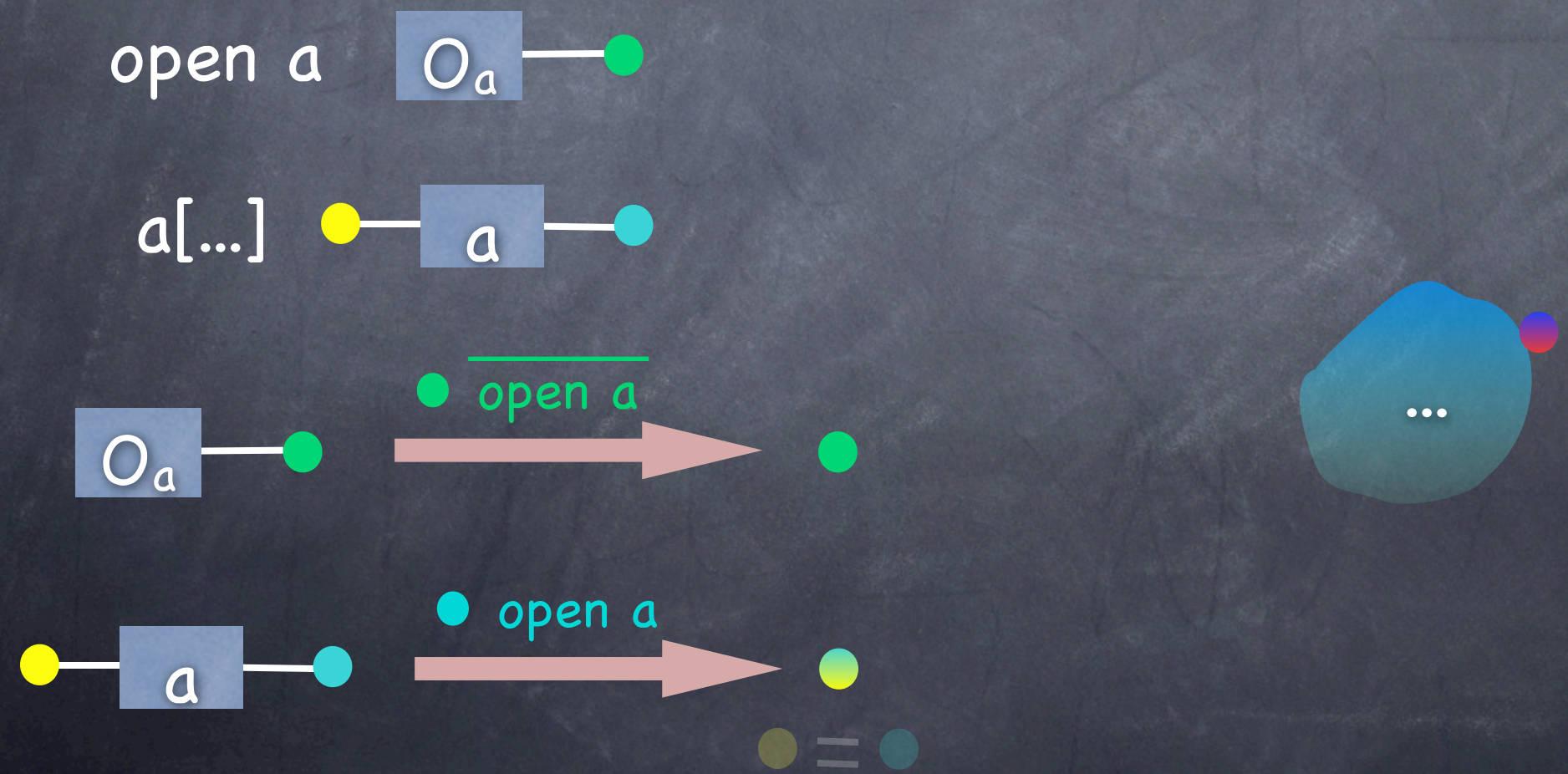

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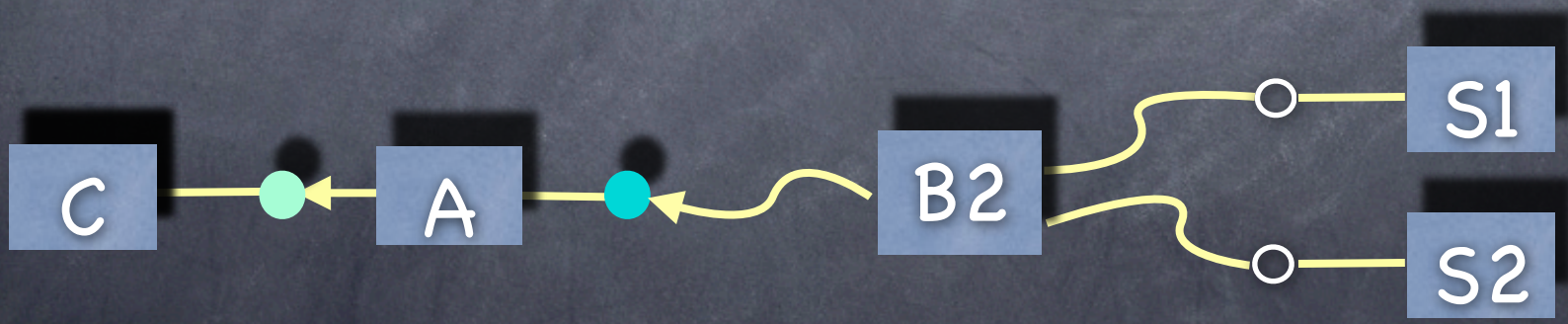

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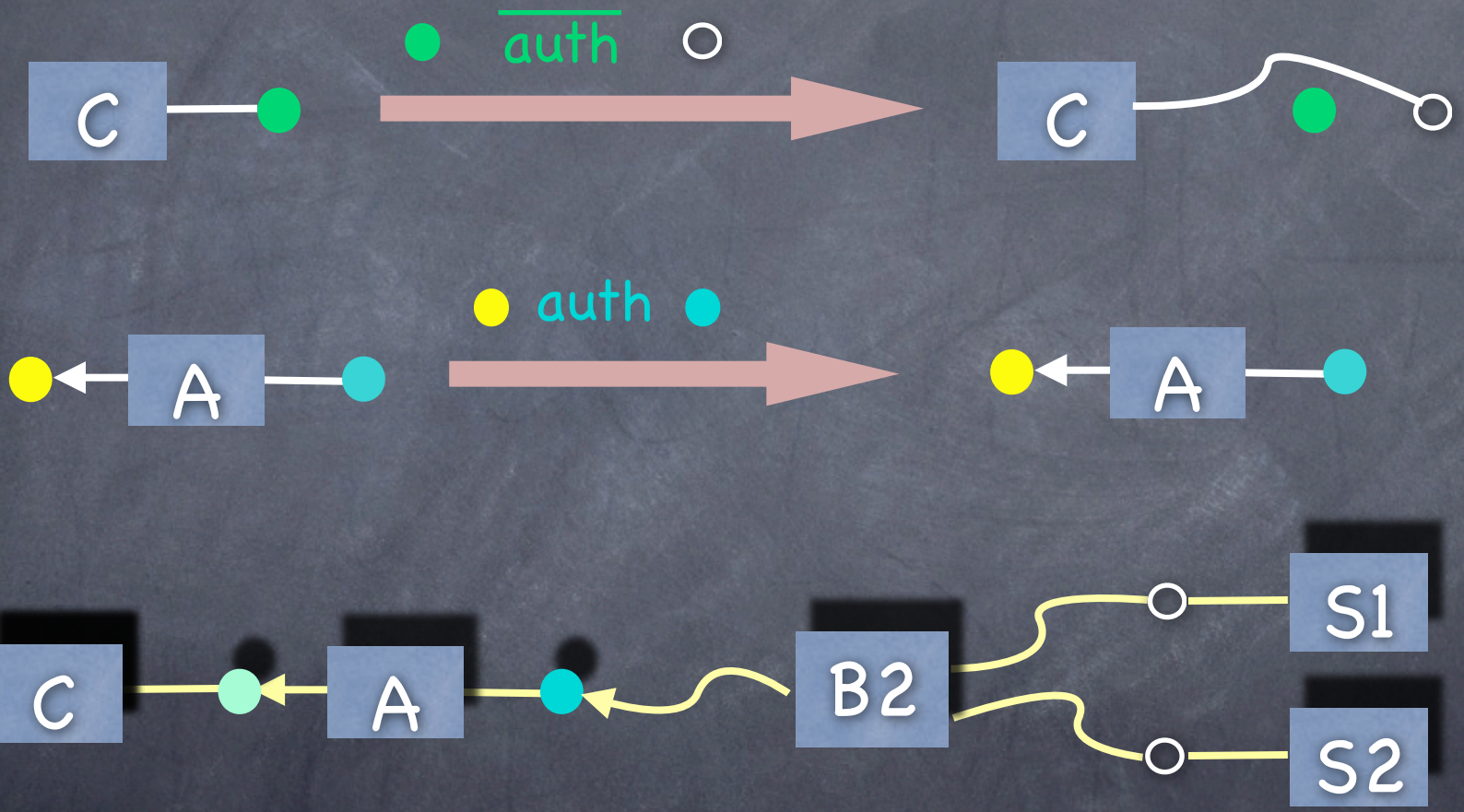

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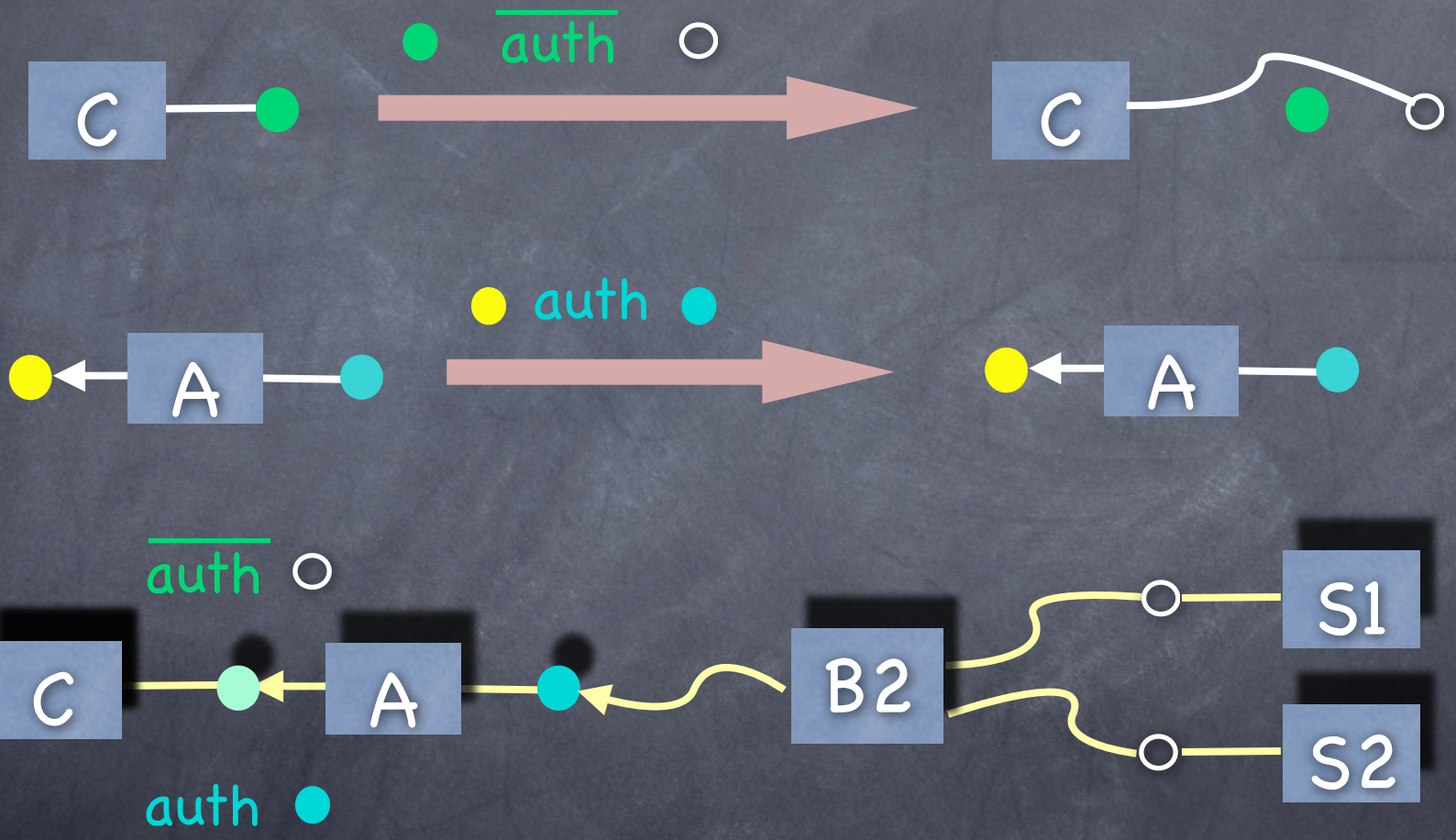


mgu( ● = ○ ) yields { ●/○ }

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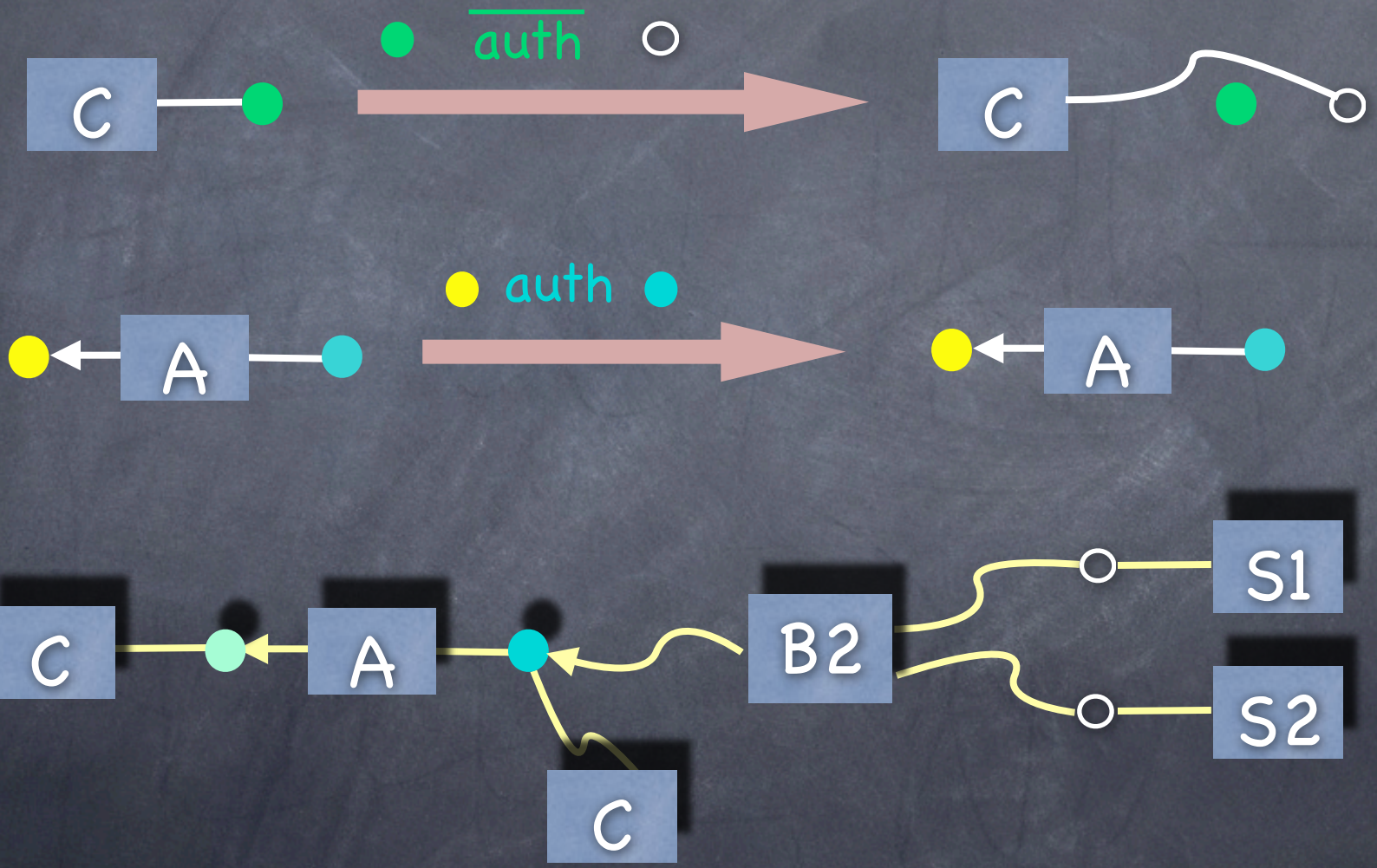

# Synchronisation algebras with mobility

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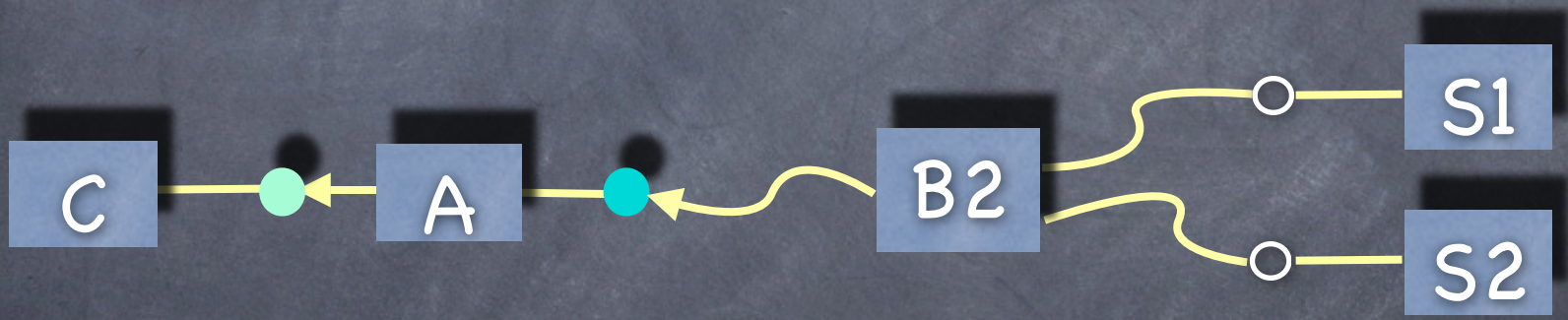

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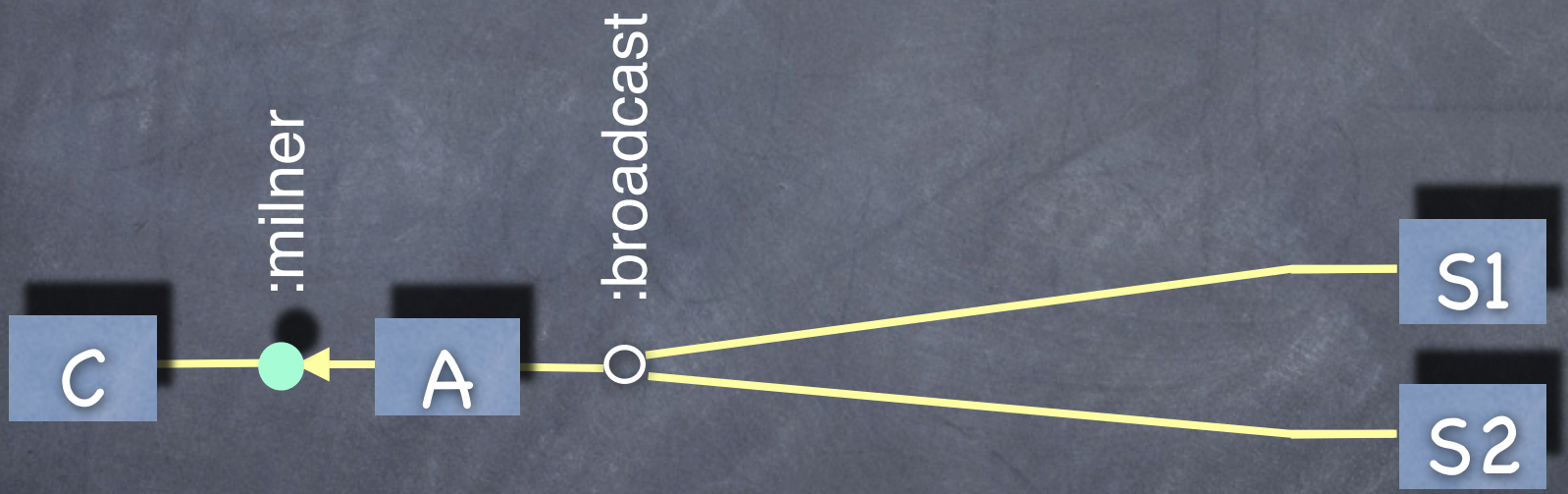
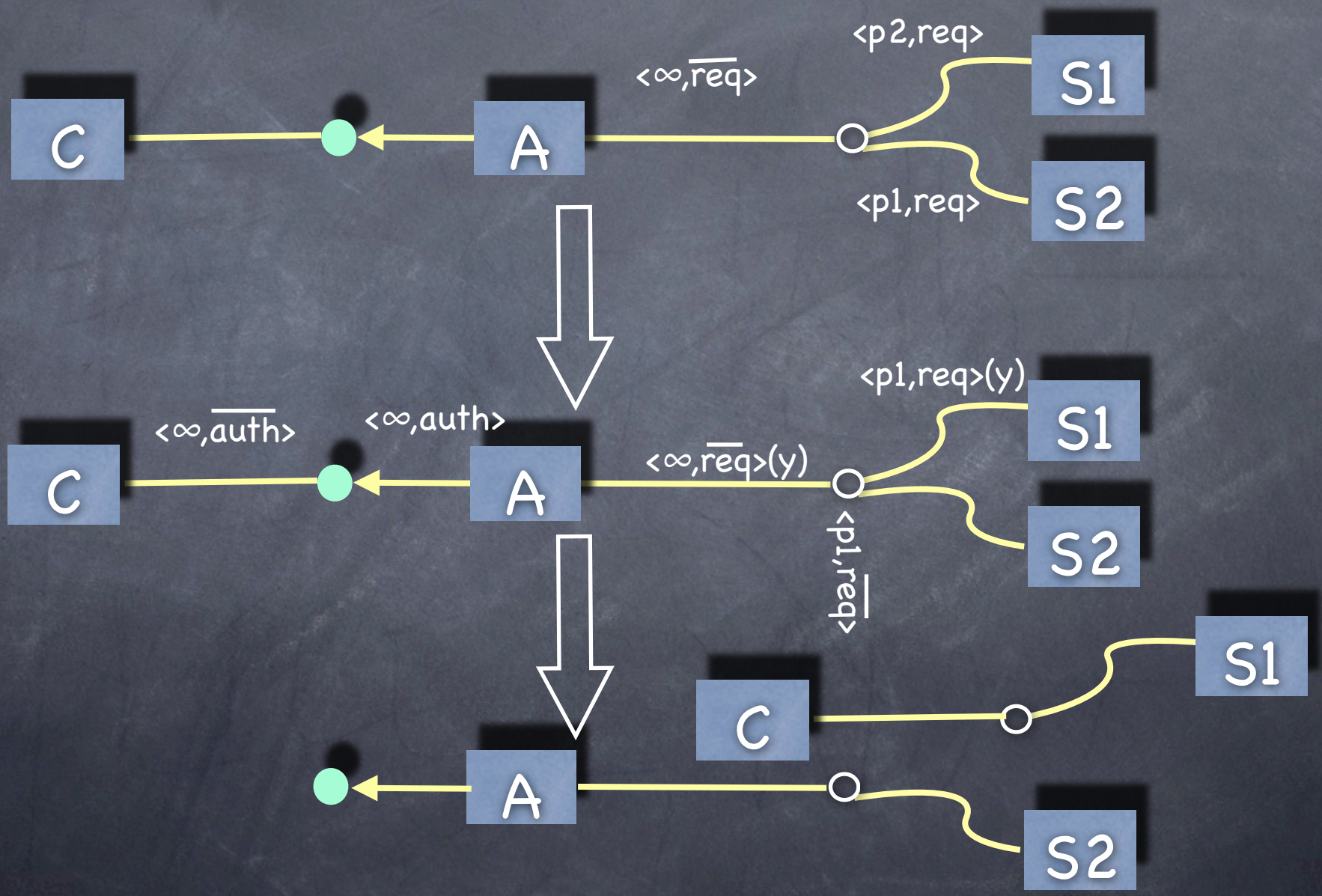


# SHReQ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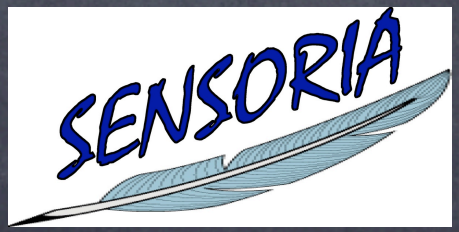
# Dealing with quality

# Watch...


...for more fun :)

# Service Oriented Architectural Design

with
R. Bruni, A. Lluch Lafuente, and U. Montanari
Dipartimento di Informatica
Universita' di Pisa



# Motivations

SEnSOria aims to develop an approach for engineering SOCs

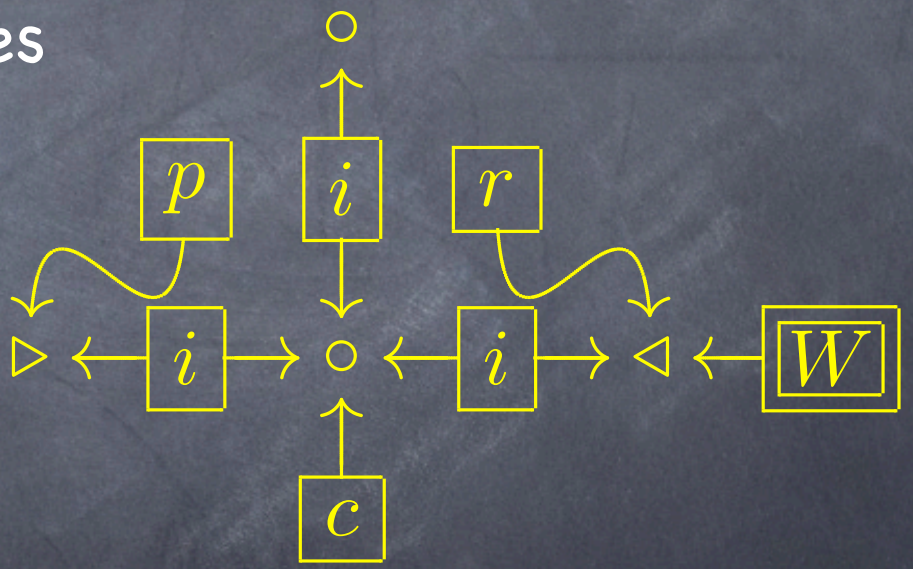
- Key issues of service-based architectures:
  - design
  - reconfiguration
- Styles for reusing existing design patterns
- Run-time changes (e.g., dynamic binding)
  - require reconfigurations of architectures
  - complement their static reconfigurations
  - driven by architectural information specified during design
- Often, architectural styles must be preserved or consistently changed

# ADR principles

- Architectures are modelled as suitable graphs
- Hierarchical architectural designs
  - style preserving rules (not original)
  - algebraic presentation (original)
- Reconfigurations defined over style proofs instead of actual architectures
  - exploits the algebraic presentation
  - straightforward definition of hierarchical and inductive reconfigurations (ordinary term rewriting and SOS)
  - only valid contexts considered (not all concrete designs)
  - matching is simpler during reconfigurations (design driven)

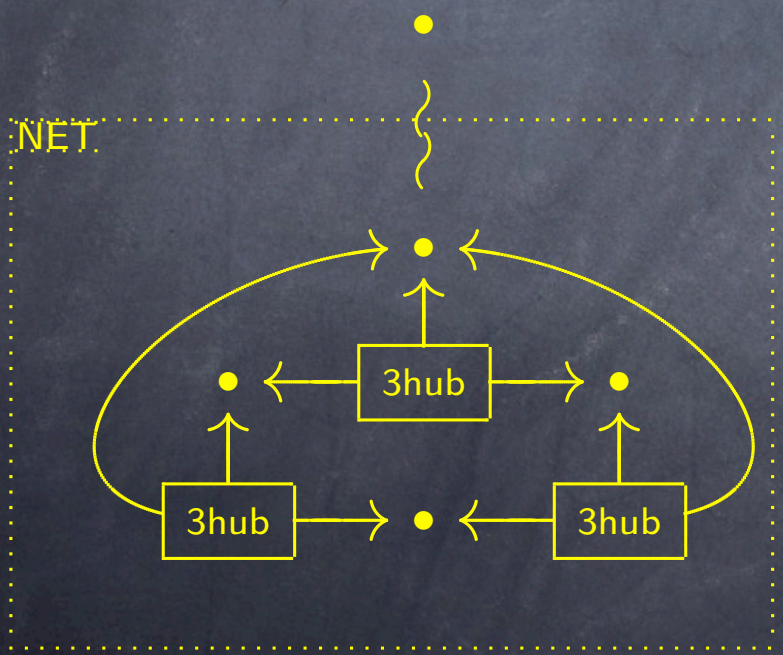
# ADR ingredients

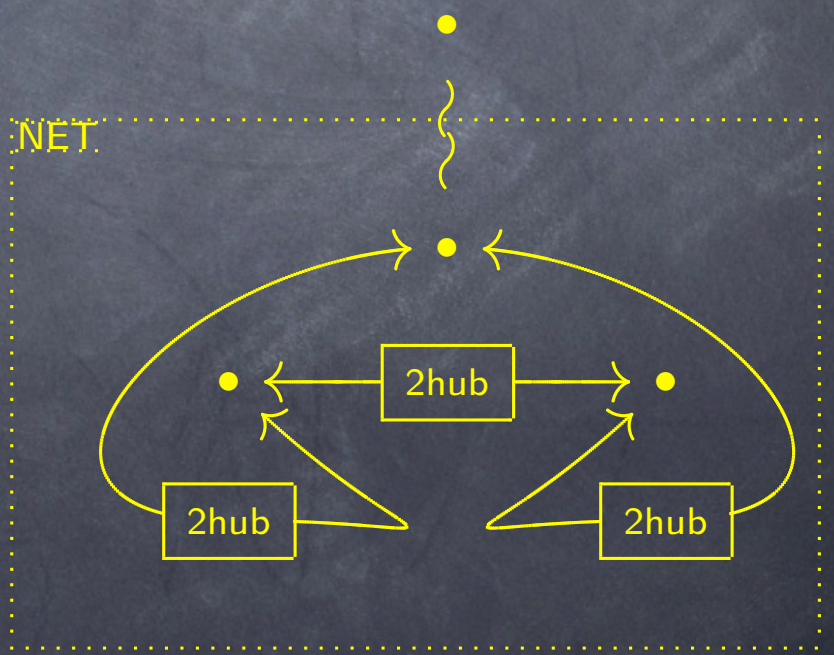
- Hypergraphs
  - edges model components: can be terminal and non-terminal edges
  - nodes model connecting ports
- Type-(hyper)graphs
- Productions
  - rules like L ::= R
  - specify how non-terminals should be replaced



# ADR by example

- A local networking architecture
- 2 styles where each network hub has degree of connectivity 2 or 3
- Connections between hubs are also driven by the style



# Designs and productions

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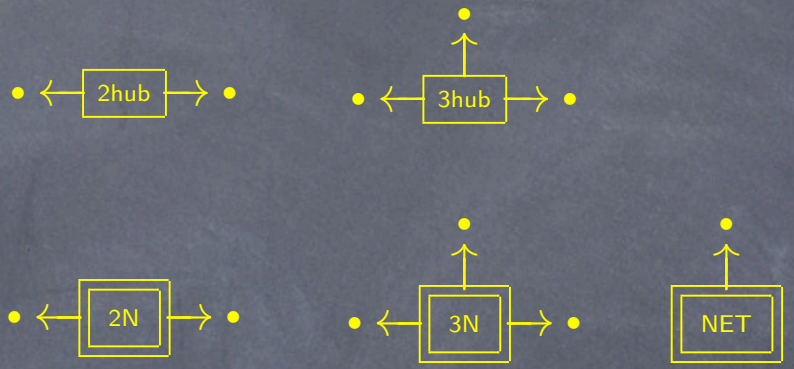
- Edges for the network example



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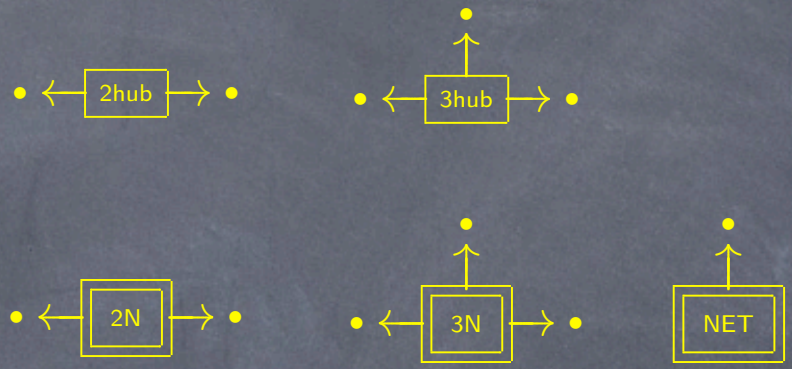
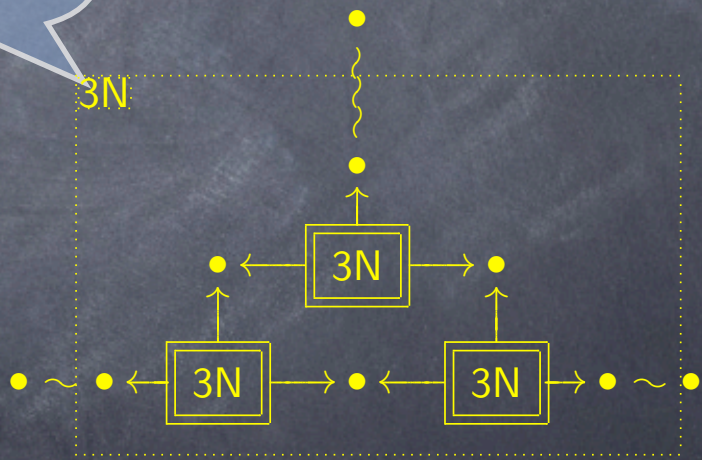



- A design consists of
  - a lhs L which is a graph made of a single non-terminal edge
  - a rhs R graph possibly containing non-terminal edges
  - a map from the nodes of L to the nodes of R
- A production is a design where the occurrences of non-terminal are distinguished

represents the abstract class of the component




$3N ::= \texttt{link3}(3N, 3N, 3N)$

$\texttt{link3} : 3N \times 3N \times 3N \to 3N$

type of the production

# ADR metaphor

- A term of a grammar is an instance of a design
- Terms with variables are partial designs
- Replacing variables corresponds to refinement
- Replacing subterms with variables corresponds to abstraction
- Replacements are driven by term rewriting rules, namely reconfiguration rules t -> t'
  - style is preserved if t and t' have the same abstract class
  - otherwise styles change...in a consistent way

# Design rewritings

$$\texttt{link3to2}: \frac{x_1 \xrightarrow{\texttt{3to2}} x_1' \quad x_2 \xrightarrow{\texttt{3to2}} x_2' \quad x_3 \xrightarrow{\texttt{3to2}} x_3'}{\texttt{link3}(x_1, x_2, x_3) \xrightarrow{\texttt{3to2}} \texttt{link2}(\texttt{link2}(x_2', x_1'), x_3')}$$

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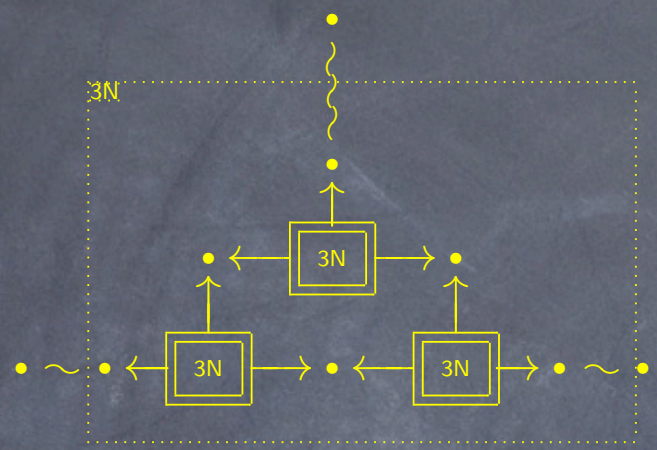
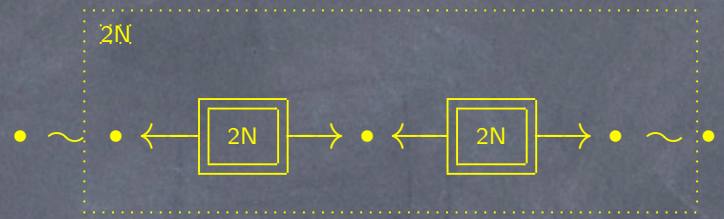



$\texttt{link2} : 2N \times 2N \to 2N$




$$\texttt{link3to2} : \frac{x_1 \xrightarrow{\texttt{3to2}} x'_1 \quad x_2 \xrightarrow{\texttt{3to2}} x'_2 \quad x_3 \xrightarrow{\texttt{3to2}} x'_3}{\texttt{link3}(x_1, x_2, x_3) \xrightarrow{\texttt{3to2}} \texttt{link2}(\texttt{link2}(x'_2, x'_1), x'_3)}$$

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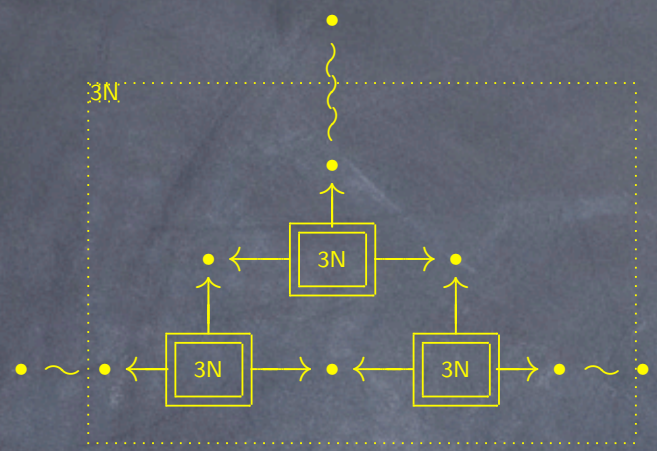


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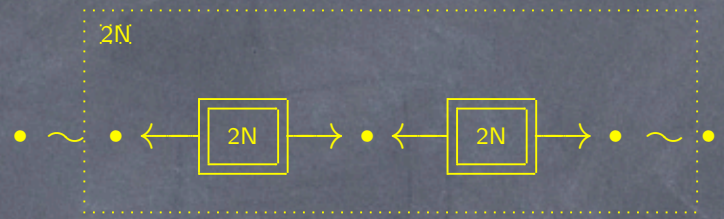


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