HD-automata with distinctions

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joint work with Kidane Yemane & Marino Miculan

Menu

A glimpse of HD-automata
Why distinctions?
HD-automata with distinctions
Conclusions

Ingredients bought from...

Ghani, Yemane, Victor [CMCS04]
Cattani, Sewell [lics00]
Fiore, Turi, Plotkin [lics99]
Gabbay, Pitts [lics99]
Fiore, Moggi, Sangiorgi [lics96]
Stark [lics96]

time

0

for a comparison, see Fiore-Staton [CMCS04] and Gadducci-Miculan-Montanari [HO Sym. Comp. 06]

An Operational Model for HD Formalisms

HD formalisms can express computations where ø new "events" can be generated ø behaviour depends on events generated in the past Examples: Petri nets VP-CCS ø nominal calculi Ø ...





Solution Names explicitly used in the model

Solution Names are local to states and transitions

"[...] identity of names does not affect the behaviour of a process[...]" [fms96]

Operations on names can be modelled (creation/deallocation)





`a' is exposed as 'v' during the transition
`c' is called 'y', afterward
`u' is freshly generated and identified as 'x'
`b' is discharged along the transition

Behavioural Minimisation

 HDA aim to yield minimal representation of process behaviour

 The minimisation procedure must preserve behaviour of processes

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Need of Symmetries

Basic HDA do not have canonical minimal representatives

P(x,y) = out x y.P(x,y) + out y x.P(x,y)

Q(x,y) = out x y.Q(y,x) + out y x.Q(y,x)





 Both minimal but not isomorphic (Pistore's thesis)

Use of Symmetries

- Symmetries are enough for obtaining a minimal realisation [MP00]
- Symmetries can model early/late pi-calculus [FMP02,FMT05] or fusion (hyperbisimulation) [FMTYV05]

ø but...

Other ingredients

[mathcall] and that equality and inequality conditions on names may affect process bisimilarity."

[fms96]

"[...] finding a mathematical formalism to ensure that extruded names are renamed injectively while other names may be renamed non-injectively is the key to understanding open bisimulation."

[gyv04]

HDA with Distinctions

Def. A distinction relation on \mathcal{N} is a pair (n, d), denoted by $n^{(d)}$, where $n \in \wp_{\text{fin}}(\mathcal{N})$ and $d \subseteq n \times n$ is a symmetric relation such that $(x, x) \notin d$, for all $x \in n$.

We equip HDA with distinctions

symmetries must
 "respect" distinctions



Nfs in Pictures



symmetries are just the identities
in general, they must respect distinctions

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HDA on DNSet

Def. An *HD*-automaton with distinctions over $L \in \text{obj}(\text{DNSet})$ is a coalgebra $K : E \to \wp_{\text{fin}}(L \otimes E)$ for the functor $T_L(_-) = \wp_{\text{fin}}(L \otimes _-)$.

The minimisation alg. on hda is a variant of Ferrari, Montanari and Pistore's one

it constructs (an approx. of)
 the final coalgebra

using an explicit
 normalisation step

 $h_{H_{(0)}}(q) \stackrel{\text{def}}{=} \emptyset, \quad \forall q \in Q_E$ $H_{(i+1)} \stackrel{\text{def}}{=} N_{i+1}(K; T(H_{(i)}))$

Concluding remarks

Classical HDA extended (?) with distinctions
Minimisation algorithm re-shaped in a more modular form
HDA with distinctions can model open pi

To be done: work out the relationships between

Set and DNSet

DNSet and pullback-preserving functors over index category of distinctions



Thank You