

An overview

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Agenda

- ④ Formal Methods and Engineering Practice
(and some motivation)
- ④ An overview of SRML as related with other languages
 - ④ SCA (Service Deployment)
 - ④ PEPA (Quantitative analysis)
 - ④ StPowla (Workflows and Business Policies)
 - ④ BPEL (Web services orchestration)

Formal Methods

- ➊ Different languages and levels of abstractions for different purposes
 - ➋ e.g., Petri-Nets is a graphical formal language to model/analyze concurrent processes
 - ➌ e.g., Pi-calculus is a process calculus to model/analyze concurrent communicating processes (it has primitives for modelling reconfiguration)
 - ➍ e.g., Z is a logic notation to specify systems ...
- ➎ Why many? Because each targets a specific aspect and aim. Focusing on a specific target and aim limits the complexity of models.
- ➏ Why formal methods? Specification, transformation, analysis, understanding the semantics of “real” languages or real problems or real scenarios

One Example: the semantics of BPEL

- BPEL is a standard for Web Service orchestration inspired by pi-calculus. It does not come with a formal semantics (although some proposals have been done in academia)
- A number of engines support BPEL. Each BPEL engine may implement a slightly different semantics
- E.g., if you have two parallel threads in a transaction and one fails: do you have to interrupt the other one? do you have to wait for its completion and then compensate it? In which order should the compensations be executed?
- Background knowledge about orchestration semantics enables BPEL programmers to update and adapt to evolving technologies (e.g., update or change).
- In general, the understanding of the key problems of a scenario is a know-how which makes it easier to familiarise and cope with evolving technologies

What about SRML?

- Service-oriented modelling is an emerging discipline and issue, supported and encouraged by major companies (Microsoft, IBM, BEA, etc)
- There are a number of formal languages that address, very specifically, the different aspects of service-oriented engineering
- There is no standard language for the modelling of service oriented systems (yet)
- SRML is the only language (at the moment) that address architectural and behavioural modelling in a whole formal framework
 - algebraic semantics of reconfiguration, logics of interaction, c-semirings
- Will I use SRML in industry? Probably not, but you will possibly use (or even define!) some language that models some of the features modelled by SRML.
- Learning an reference language which is general enough, because defined for being such, provides a basis to cope with an evolving set of technologies in the context of distributed dynamically reconfiguring systems

A view of ensemble

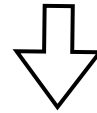
use case diagrams

from scratch
use case -> SRML

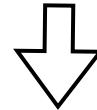
from scratch
CBD

from scratch
reuse
refinement

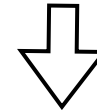
Define requirements



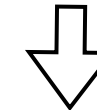
Define modules + EX-Is



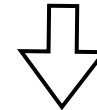
Define internal structure



Define behaviours



Analyse

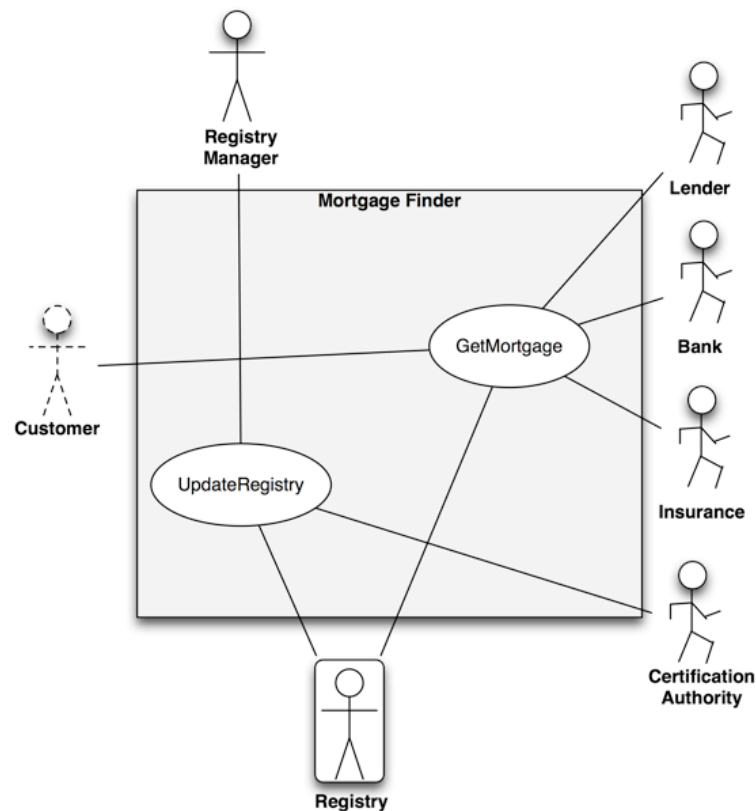


Deploy

The mortgage example

use case diagrams

Define requirements

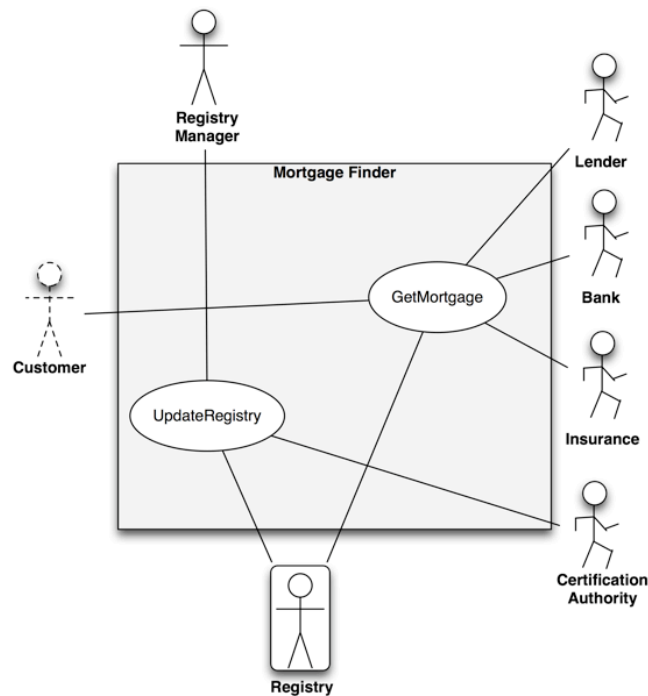


the Mortgage scenario

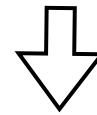
The mortgage example

use case diagrams

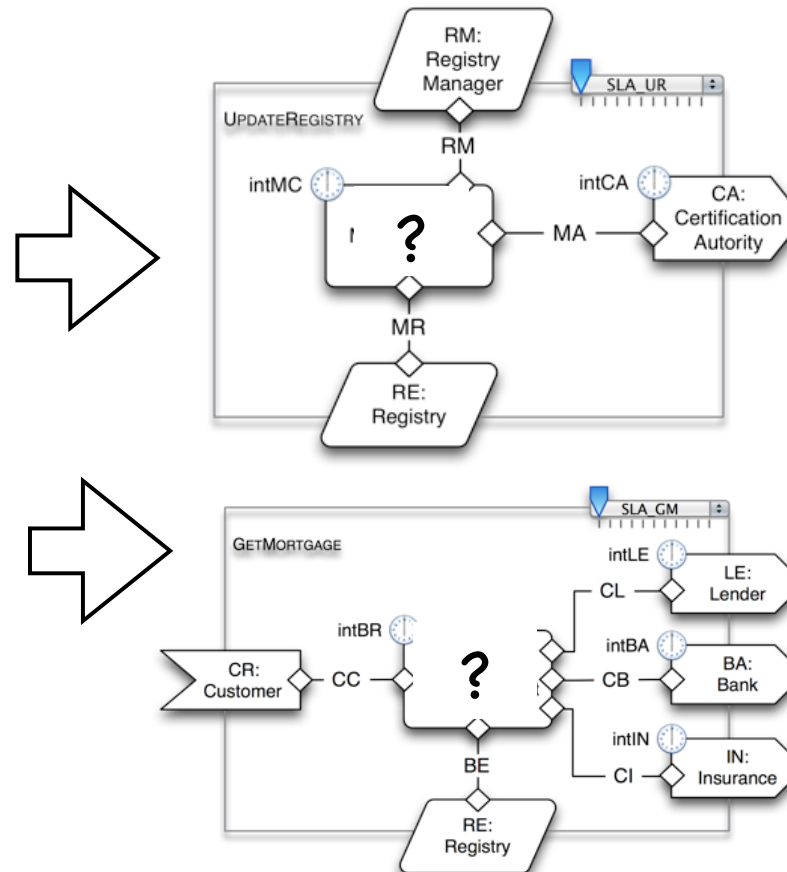
from scratch
use case -> SRML



Define requirements



Define modules + EX-Is



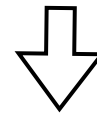
A view of ensemble

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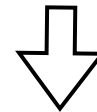
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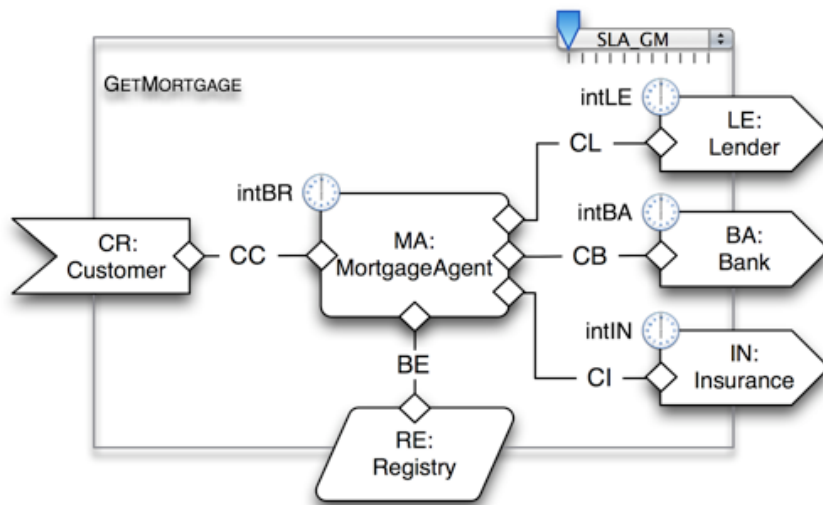
Define requirements



Define modules + EX-Is



Define internal structure



more component could be defined
for modular development, reuse,
limiting the complexity etc.

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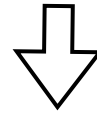
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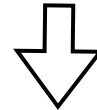
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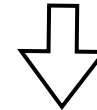
Define requirements



Define modules + EX-Is



Define internal structure



Define behaviours

The specifications (interactions)

BUSINESS ROLE MortgageAgent is

INTERACTIONS

r&s getProposal

⌚ idData:usrdata,
income:moneyvalue,
preferences:prefdata,
✉ proposal:mortgageproposal
cost:moneyvalue

s&r askProposal

⌚ idData:usrdata,
income:moneyvalue,
✉ proposal:mortgageproposal
loanData:loandata,
accountIncluded:bool,
insuranceRequired:bool

s&r getInsurance

⌚ idData:usrdata,
loanData:loandata,
✉ insuranceData:insurancedata

s&r openAccount

⌚ idData:usrdata,
loanData:loandata,
✉ accountData:accountdata

s&r signOutLoan

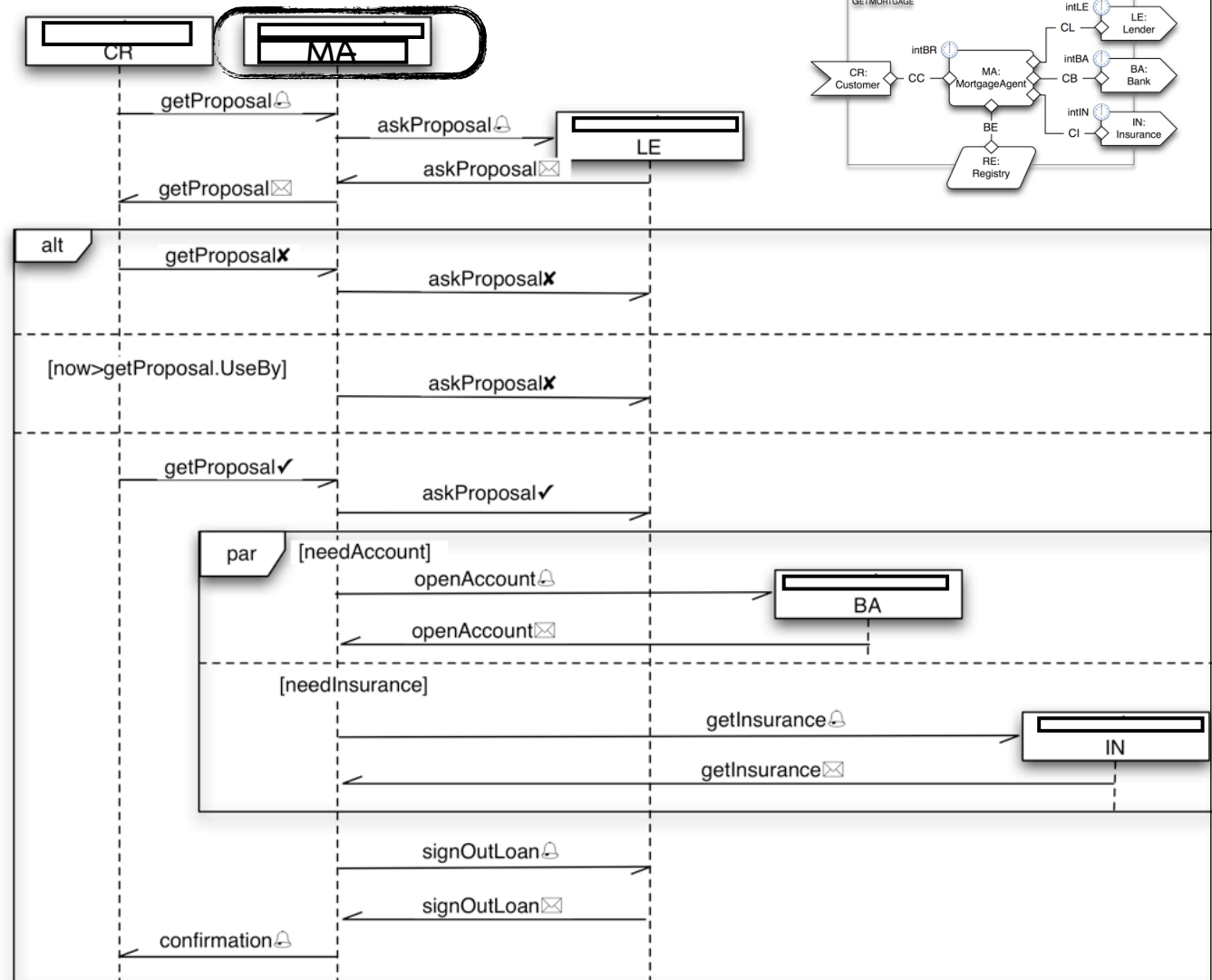
⌚ insuranceData:insurancedata,
accountData:accountdata,
✉ contract:loancontract

snd confirmation

⌚ contract:loancontract

ask getLenders(prefdata):setids

tll regContract(loandata,loancontract



The specifications (textual notation)

SPECIFICATIONS

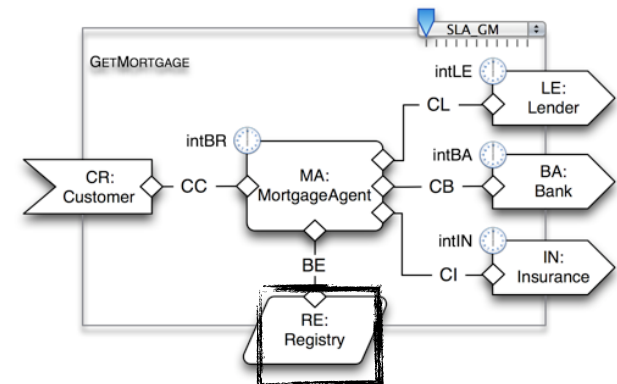
LAYER PROTOCOL Registry is

INTERACTIONS

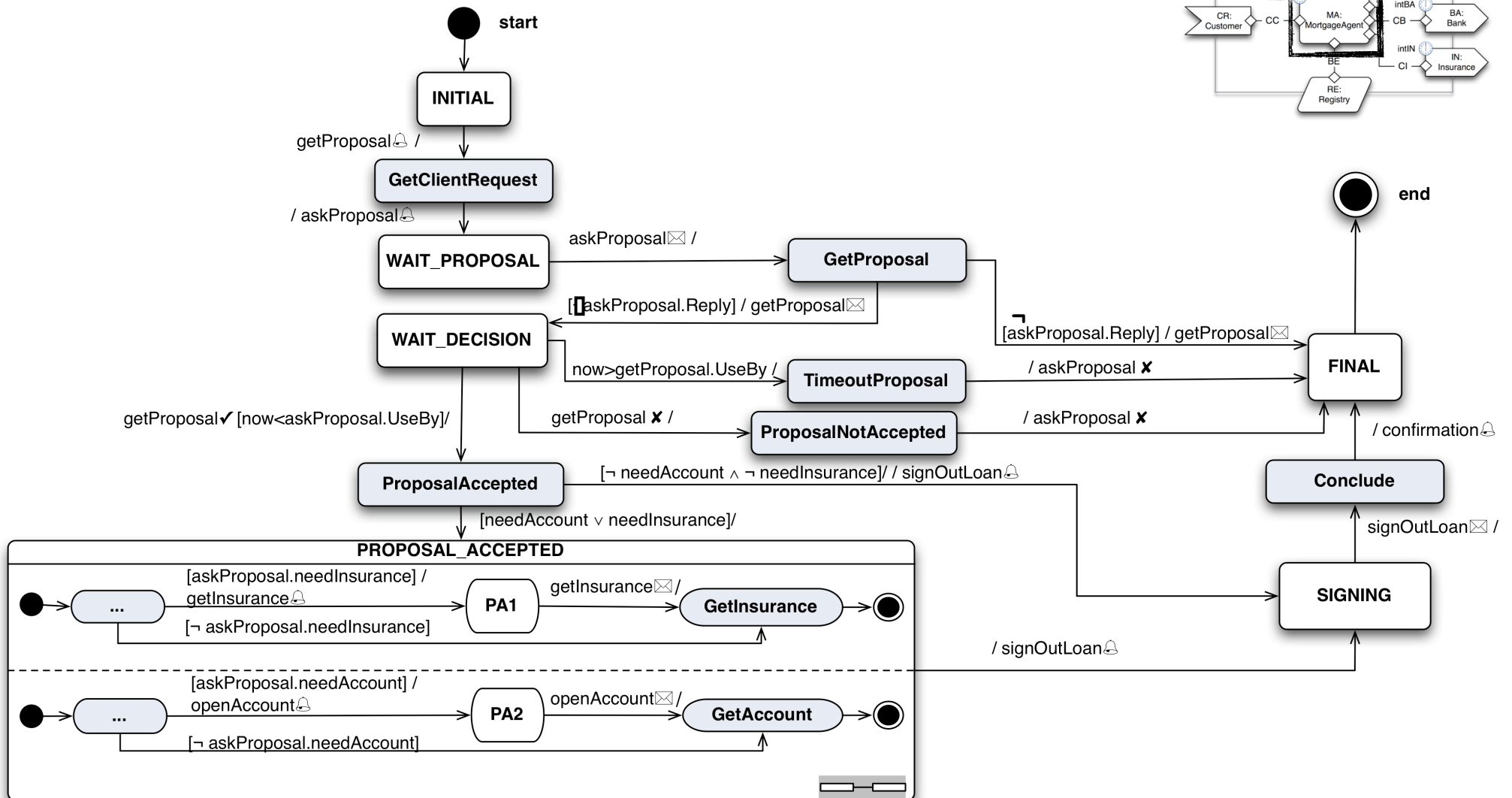
rpl getLenders(prefdata):setids

prf registerContract(loandata,loancontract)

BEHAVIOUR



The orchestration as a statechart



The specifications (the business role – interactions)

BUSINESS ROLE MortgageAgent is

INTERACTIONS

```

r&s getProposal
  Ⓐ idData:usrdata,
    income:moneyvalue,
    preferences:prefdata,
  ☒ proposal:mortgageproposal
    cost:moneyvalue
s&r askProposal
  Ⓐ idData:usrdata,
    income:moneyvalue,
  ☒ proposal:mortgageproposal
    loanData:loandata,
    accountIncluded:bool,
    insuranceRequired:bool
s&r getInsurance
  Ⓐ idData:usrdata,
    loanData:loandata,
  ☒ insuranceData:insurancedata
s&r openAccount
  Ⓐ idData:usrdata,
    loanData:loandata,
  ☒ accountData:accountdata
s&r signOutLoan
  Ⓐ insuranceData:insurancedata,
    accountData:accountdata,
  ☒ contract:loancontract
snd confirmation
  Ⓐ contract:loancontract
ask getLenders(prefdata):setids
tll regContract(loandata,loancontract)
  
```

SLA VARIABLES

CHARGE:[0..100]

ORCHESTRATION

```

local  s:[INITIAL, WAIT_PROPOSAL, WAIT_DECISION,
        PROPOSAL_ACCEPTED, SIGNING, FINAL],
        lenders:setids,
        needAccount, needInsurance:bool,
        insuranceData:insurancedata, accountData:accountdata
  
```

transition GetClientRequest

```

triggeredBy getProposalⒶ
guardedBy s=INITIAL
effects s'=WAIT_PROPOSAL
  ⌞ lenders'=getLenders(prefdata)
sends askProposalⒶ
  ⌞ askProposal.idData=getProposal.idData
  ⌞ askProposal.income=getProposal.income
  
```

transition GetProposal

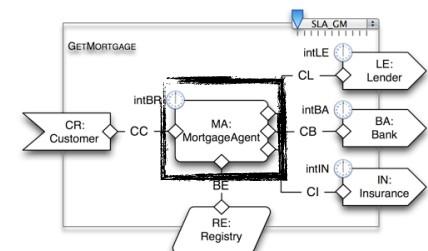
```

triggeredBy askProposal☒
guardedBy s=WAIT_PROPOSAL
effects needAccount'=askProposal.accountIncluded
  ⌞ needInsurance'=askProposal.insuranceRequired
  ⌞ askProposal.Reply ⊃ s'=WAIT_DECISION
  ⌞ ¬askProposal.Reply ⊃ s'=FINAL
sends getProposal☒
  ⌞ getProposal.Reply=askProposal.Reply
  ⌞ getProposal.proposal=askProposal.proposal
  ⌞ getProposal.cost=(CHARGE/100+1)*750
  
```

transition TimeoutProposal

```

triggeredBy now>getProposal.UseBy
guardedBy s=WAIT_DECISION
effects s'=FINAL
sends askProposal*
  
```



The specifications (the business role – interactions)

BUSINESS ROLE MortgageAgent is

INTERACTIONS

```

r&s getProposal
  Ⓐ idData:usrdata,
    income:moneyvalue,
    preferences:prefdata,
  ☒ proposal:mortgageproposal
    cost:moneyvalue
s&r askProposal
  Ⓐ idData:usrdata,
    income:moneyvalue,
  ☒ proposal:mortgageproposal
    loanData:loandata,
    accountIncluded:bool,
    insuranceRequired:bool
s&r getInsurance
  Ⓐ idData:usrdata,
    loanData:loandata,
  ☒ insuranceData:insurancedata
s&r openAccount
  Ⓐ idData:usrdata,
    loanData:loandata,
  ☒ accountData:accountdata
s&r signOutLoan
  Ⓐ insuranceData:insurancedata,
    accountData:accountdata,
  ☒ contract:loancontract
snd confirmation
  Ⓐ contract:loancontract
ask getLenders(prefdata):setids
tll regContract(loandata,loancontract)
  
```

SLA VARIABLES

CHARGE:[0..100]

ORCHESTRATION

```

local s:[INITIAL, WAIT_PROPOSAL, WAIT_DECISION,
    PROPOSAL_ACCEPTED, SIGNING, FINAL],
    lenders:setids,
    needAccount, needInsurance:bool,
    insuranceData:insurancedata, accountData:accountdata
  
```

...

transition ProposalNotAccepted

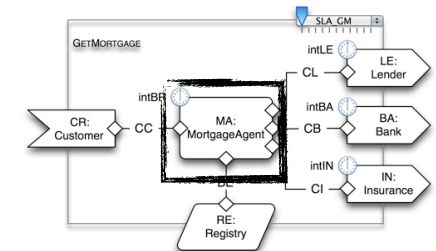
```

triggeredBy getProposal*
guardedBy s=WAIT DECISION
  ∧ now<askProposal.UseBy
effects s'=FINAL
sends askProposal*
  
```

transition ProposalAccepted

```

triggeredBy getProposal✓
guardedBy s=WAIT DECISION
  ∧ now<deadline
effects needAccount ∨ needInsurance ⊃ s'=PROPOSAL ACCEPTED
  ∧ ¬needAccount ∧ ¬needInsurance ⊃ s'=SIGNING
sends askProposal✓
  ∧ needAccount ⊃ openAccountⒶ
    ∧ openAccount.idData=getProposal.idData
    ∧ openAccount.loanData=getProposal.loanData
  ∧ needInsurance ⊃ getInsuranceⒶ
    ∧ getInsurance.idData=getProposal.idData
    ∧ getInsurance.loanData=getProposal.loanData
  ∧ ¬needAccount ∧ ¬needInsurance ⊃ signOutLoanⒶ
    ∧ signOutLoan.insuranceData=insuranceData
    ∧ signOutLoan.accountData=accountData
  
```



The specifications (the business role – interactions)

BUSINESS ROLE MortgageAgent is

INTERACTIONS

```

r&s getProposal
  ⚙ idData:usrdata,
  income:moneyvalue,
  preferences:prefdata,
  ☒ proposal:mortgageproposal
  cost:moneyvalue

s&r askProposal
  ⚙ idData:usrdata,
  income:moneyvalue,
  ☒ proposal:mortgageproposal
  loanData:loandata,
  accountIncluded:bool,
  insuranceRequired:bool

s&r getInsurance
  ⚙ idData:usrdata,
  loanData:loandata,
  ☒ insuranceData:insurancedata

s&r openAccount
  ⚙ idData:usrdata,
  loanData:loandata,
  ☒ accountData:accountdata

s&r signOutLoan
  ⚙ insuranceData:insurancedata,
  accountData:accountdata,
  ☒ contract:loancontract

snd confirmation
  ⚙ contract:loancontract

ask getLenders(prefdata):setids

tll regContract(loandata,loancontract)
  
```

SLA VARIABLES

CHARGE:[0..100]

ORCHESTRATION

local s:[INITIAL, WAIT_PROPOSAL, WAIT_DECISION,
PROPOSAL_ACCEPTED, SIGNING, FINAL],

...

transition GetAccount

```

triggeredBy openAccount☒
guardedBy s=PROPOSAL ACCEPTED
effects needAccount'=false
  ∧ -needInsurance ⊃ s'=SIGNING
  ∧ accountData=openAccount.accountData
sends -needInsurance ⊃ signOutLoanⓂ
  ∧ signOutLoan.insuranceData=insuranceData
  ∧ signOutLoan.accountData=accountData
  
```

transition GetInsurance

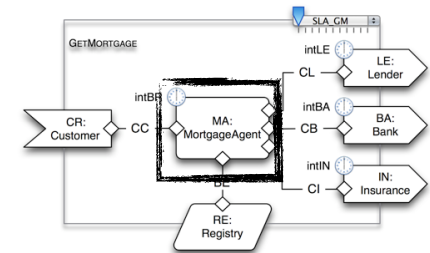
```

triggeredBy getInsurance☒
guardedBy s=PROPOSAL ACCEPTED
effects needInsurance'=false
  ∧ -needAccount ⊃ s'=SIGNING
  ∧ insuranceData=getInsurance.insuranceData
sends -needAccount ⊃ signOutLoanⓂ
  ∧ signOutLoan.insuranceData=insuranceData
  ∧ signOutLoan.accountData=accountData
  
```

transition Conclude

```

triggeredBy signOutLoan☒
guardedBy s=SIGNING
effects s'=FINAL
sends confirmationⓂ
  ∧ confirmation.contract=signOutLoan.contract
  ∧ regContract(askProposal.loanData,signOutLoan.contract)
  
```



The specifications (the business protocols)

BUSINESS PROTOCOL Bank is

INTERACTIONS

```
r&s newMortgageAccount
  ⌚ idData:usrdata,
    loanData:loandata,
  ☒ accountData:accountdata
```

BEHAVIOUR

```
initiallyEnabled newMortgageAccount⌚?
newMortgageAccount.Reply after newMortgageAccount☒!
```

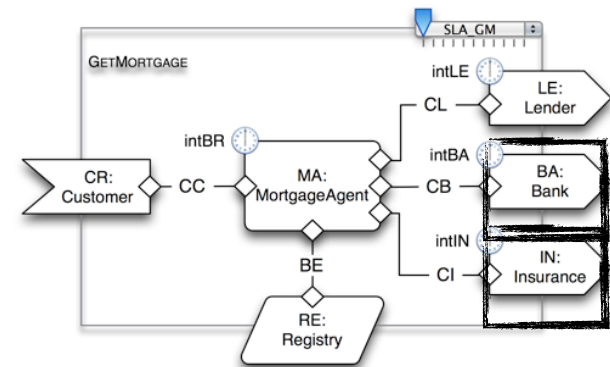
BUSINESS PROTOCOL Insurance is

INTERACTIONS

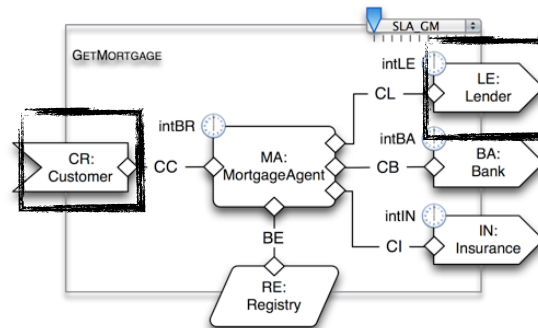
```
r&s newMortgageInsurance
  ⌚ idData:usrdata,
    loanData:loandata,
  ☒ insuranceData:insurancedata
```

BEHAVIOUR

```
initiallyEnabled newMortgageInsurance⌚?
newMortgageInsurance.Reply after newMortgageInsurance☒!
```



The specifications (the business protocols)



BUSINESS PROTOCOL Customer is

INTERACTIONS

r&s getProposal

- ⌚ idData:usrdata,
income:moneyvalue,
preferences:prefdata,
- ✉ proposal:mortgageproposal
cost:moneyvalue

snd confirmation

- ⌚ contract:loancontract

SLA VARIABLES

CHARGE:[0..100]

BEHAVIOUR

initiallyEnabled getProposal⌚?
 getProposal.cost ≤ 750 * (CHARGE / 100 + 1) **after** getProposal✉!
 getProposal✓? **ensures** confirmation⌚!

BUSINESS PROTOCOL Lender is

INTERACTIONS

r&s requestMortgage

- ⌚ idData:usrdata,
income:moneyvalue,
- ✉ proposal:mortgageproposal
loanData:loandata,
accountIncluded:bool,
insuranceRequired:bool

r&s requestSignOut

- ⌚ insuranceData:insurancedata,
accountData:accountdata,
- ✉ contract:loancontract

BEHAVIOUR

initiallyEnabled requestMortgage⌚?
 requestMortgage✓? **enables** requestSignOut⌚?
 requestSignOut✉.Reply **after** requestSignOut✉?

END SPECIFICATIONS

A view of ensemble

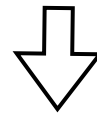
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from scratch
use case -> SRML

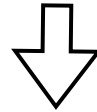
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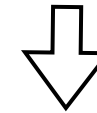
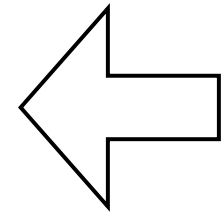
Define requirements



Define modules + EX-Is



Define internal structure
(textual notation)



Define behaviours

The textual definition (nodes)

MODULE GETMORTGAGE **is**

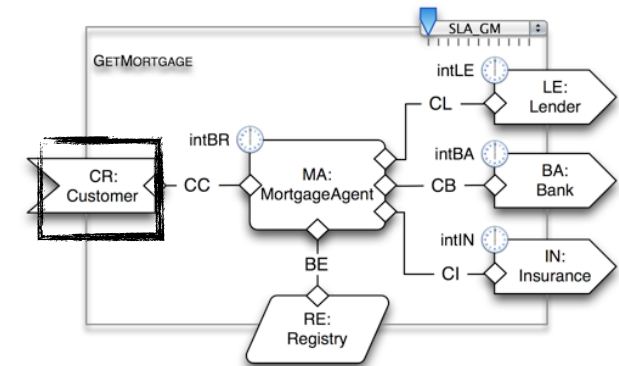
DATATYPES

sorts: usrdata, prefdata,
moneyvalue, mortgageproposal,
loandata, loancontract,
insurancedata, accountdata,
setids, bool, nat

PROVIDES

CR: Customer

CR Customer	MA MortgageAgent
r&s getProposal 🔔 idData income preferences ✉ proposal cost	r&s getProposal 🔔 idData income preferences ✉ proposal cost
snd confirmation 🔔 contract	snd confirmation 🔔 contract
SLA VARIABLES CHARGE	SLA VARIABLES CHARGE



The textual definition (nodes and internal policies)

REQUIRES

LE: Lender
 intLE⌚trigger: getproposal🔔?
BA: Bank
 intBA⌚trigger: default
IN: Insurance
 intIN⌚trigger: default

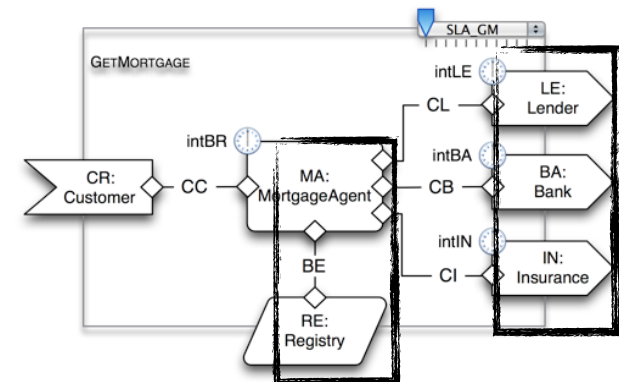
COMPONENTS

MA: MortgageAgent
 intMA⌚init: s=INITIAL
 intMA⌚term: s=FINAL

USES

RE: Registry

...



The textual definition (external policies)

EXTERNAL POLICY

SLA VARIABLES

MA.CHARGE, MA.getProposal^{*},

LE.ServiceId, LE.COST, LE.requestMortgage^{*}

CONSTRAINTS

C1 relates the charge with the deadline of getProposal

C₁: {MA.CHARGE, MA.getProposal^{*}}

$$\text{def}(c, t) = \begin{cases} 1 & \text{if } t \leq 10 * c \\ 1 + 2 * c - 0.2 * t & \text{if } 10 * c < t \leq 5 + 10 * c \\ 0 & \text{otherwise} \end{cases}$$

e.g., if I give 2 hours and the charge is more than 20 then I am happy

C2 enforces that the chosen lender is in the set of trusted lenders stored in the registry

C₂: {LE.ServiceId}

$$\text{def}(s) = \begin{cases} 1 & \text{if } s \in \text{MA.lenders} \\ 0 & \text{otherwise} \end{cases}$$

e.g., if I give 2 hours and the charge is between 10 and 20 then the satisfaction depends on c and t

C3 ensures that the deadline for proposal can be met wrt the delays

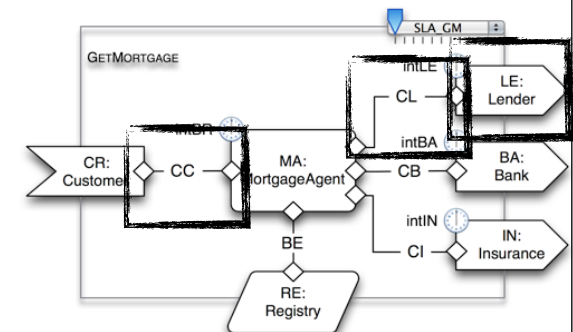
C₃: {MA.getProposal^{*}, LE.requestMortgage^{*}},

$$\text{def}(t1, t2) = \begin{cases} 1 & \text{if } t2 > t1 + \text{CC.Delay} + \text{CL.Delay} \\ 0 & \text{otherwise} \end{cases}$$

C4 ensures that the cost is inversely proportional to the satisfaction and directly proportional to the deadline of requestMortgage

C₄: {LE.COST, LE.requestMortgage^{*}}

$$\text{def}(c, t) = \begin{cases} \frac{1}{c} + \frac{t}{100} & \text{if } c < 500 \\ 0 & \text{otherwise} \end{cases}$$



The textual definition (wires part1)

WIRES

MA MortgageAgent	C₄	BE	d₄	RE Registry
ask getLenders	S ₁	Straight. A(prefdata)R(setids)	R ₁	rpl getLenders
tll regContract	S ₁	Straight. T(loandata,loancontract)	R ₁	prf registerContract

MA MortgageAgent	C₁	CB	d₁	BA Bank
s&r openAccount Ⓐ idData loanData ☒ accountData	S ₁ i ₁ i ₂ o ₁	Straight. I(usrdata, loandata) O(accountdata)	R ₁ i ₁ i ₂ o ₁	r&s newMortgageAccount Ⓐ idData loanData ☒ accountData

MA MortgageAgent	C₁	CI	d₁	IN Insurance
s&r getInsurance Ⓐ idData loanData ☒ insuranceData	S ₁ i ₁ i ₂ o ₁	Straight. I(usrdata, loandata) O(insurancedata)	R ₁ i ₁ i ₂ o ₁	r&s newMortgageInsurance Ⓐ idData loanData ☒ insuranceData

...

The textual definition (wires part2)

c₁	CC	d₁	MA MortgageAgent
S ₁ i ₁ i ₂ i ₃ o ₁ o ₂	Straight. I(usrddata, moneyvalue,prefdata) O(mortageproposal, moneyvalue)	R ₁ i ₁ i ₂ i ₃ o ₁ o ₂	r&s getProposal Ⓐ idData income preferences ⓧ proposal cost
R ₁ i ₁	Straight O(loancontract)	S ₁ i ₁	snd confir- mation Ⓐ contract

END MODULE

This was the end of the module
The module refers to a number of specifications
(as seen before)

A view of ensemble

use case diagrams

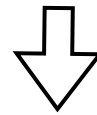
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CBD

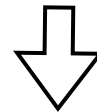
from scratch
reuse
refinement

PEPA
Logics of Interactions

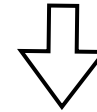
Define requirements



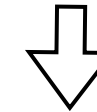
Define modules + EX-Is



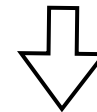
Define internal structure



Define behaviours



Analyse



Deploy

SRML2PEPA

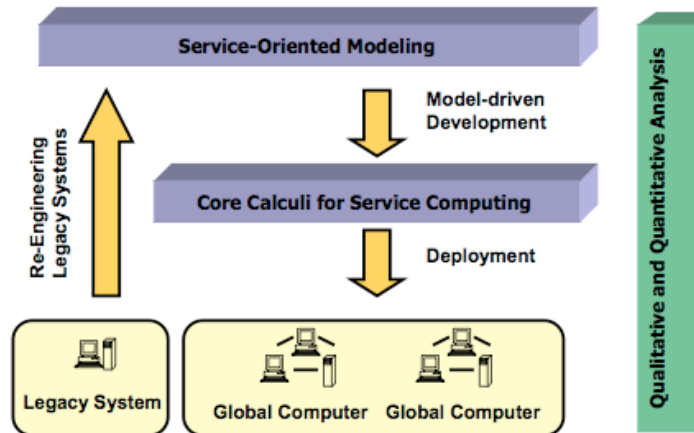


A Formal Approach to Model Time Properties in Service-Oriented Systems

Bocchi, Fiadeiro, Gilmore, Abreu, Solanki, Vankayala

<http://www.cs.le.ac.uk/people/jfiadeiro/Papers/SRML-T.pdf>

SRML



PEPA

- PEPA is a process algebra for stochastic quantitative analysis whose building entities are (1) components and (2) activities (a,r) where a is an action and r is a rate
- There are tools for PEPA (Eclipse Plugin) that allow to make quantitative analysis
- We provided an encoding from SRML modules to PEPA terms to perform quantitative analysis on SRML modules

SRML -> PEPA

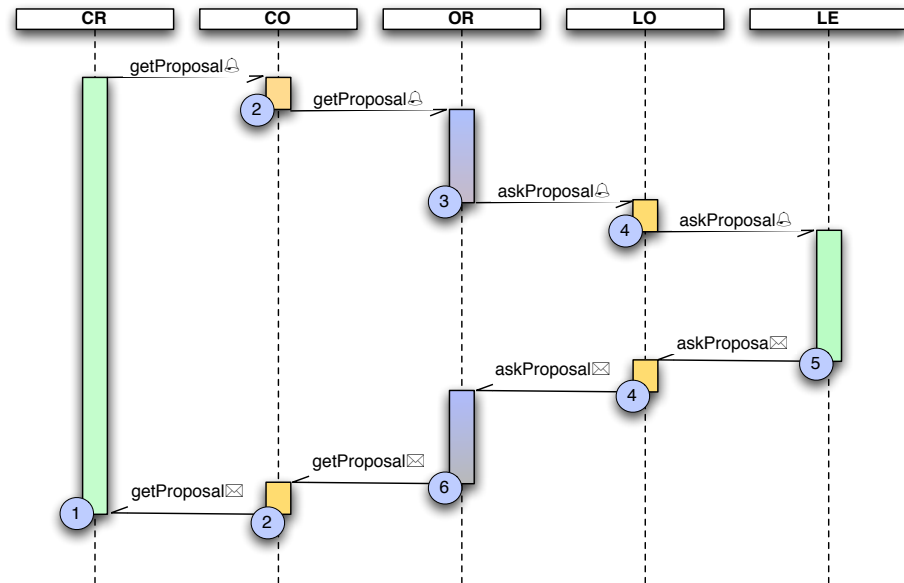
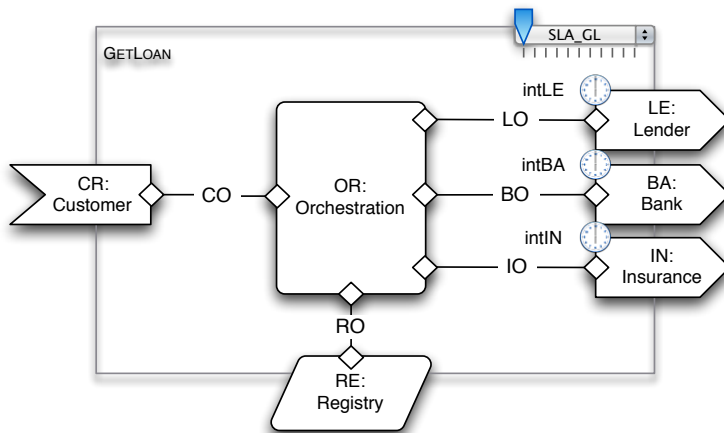
- ⦿ The aim is to enable quantitative analysis of SRML modules
- ⦿ SRML describes behaviours terms of interaction events
- ⦿ We want to determine the delay between couple of interaction events

“is the reply to the event `getProposal`, in the 80% of the cases, received in 7s?”

- ⦿ We want to determine how the single rates influence the overall delay between couples of interaction events

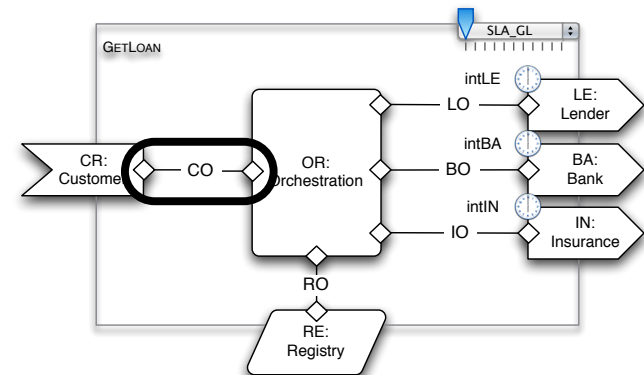
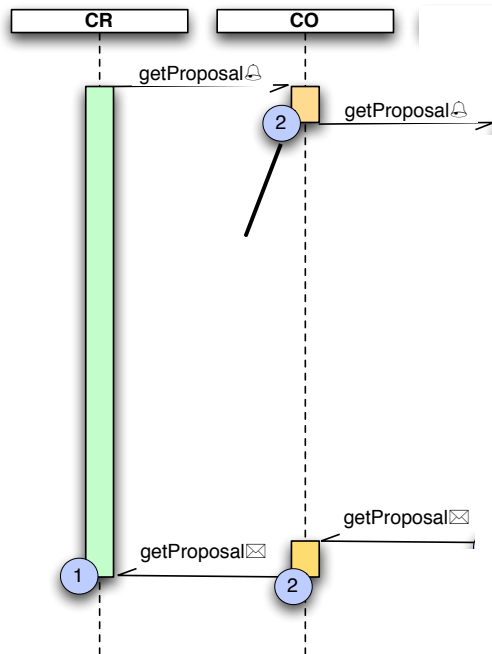
“which rates worth to be improved?”

The aim of the encoding



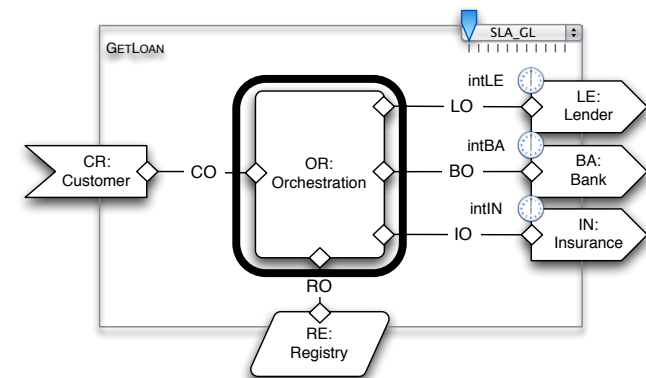
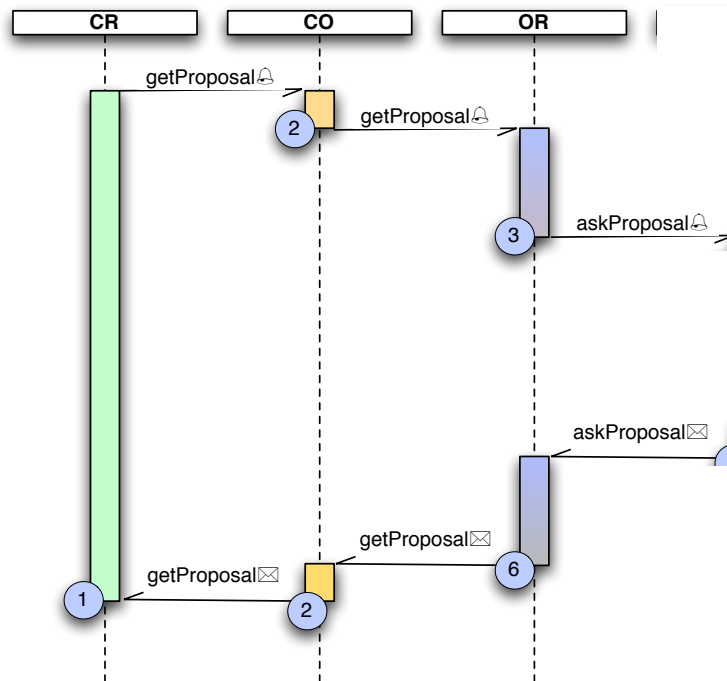
- We want to determine the delay between couples of events in the provides-interface
- For example, we want to analyze the delay between `getProposal`🔔 and `getProposal`✉ in CR ((1) in the sequence diagram)
- In order to analyze time-related properties in a SRML module:
 - we determine which delays occur in a SRML module
 - we encode SRML into PEPA

Delays in SRML



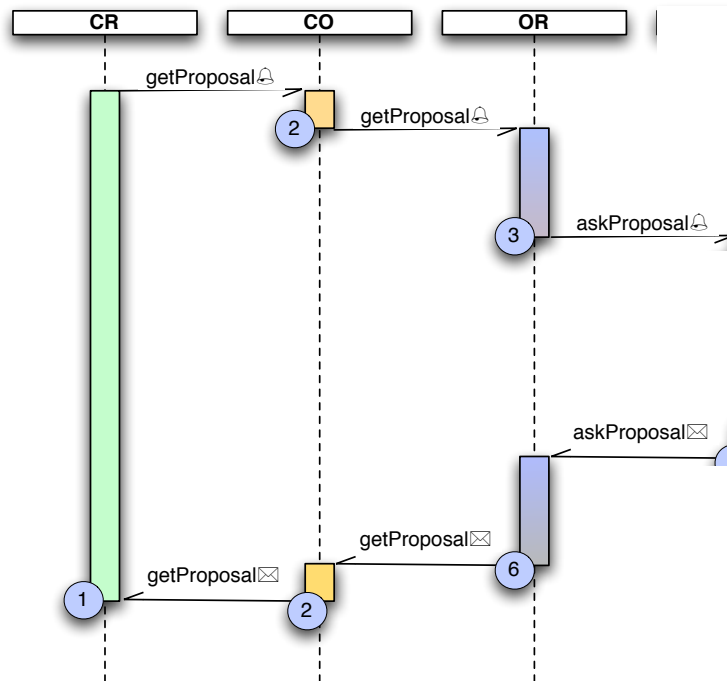
- CO transfers `getProposal` from CR to OR
- Wires can take some amount of time to transfer events between nodes
 - because of bandwidth/capacity
 - because of the execution of an interaction protocol
- Each wire has its own TRANSFER RATE
- For example, the delay of wire CO is `CO.transferRate`

Delays and dependencies



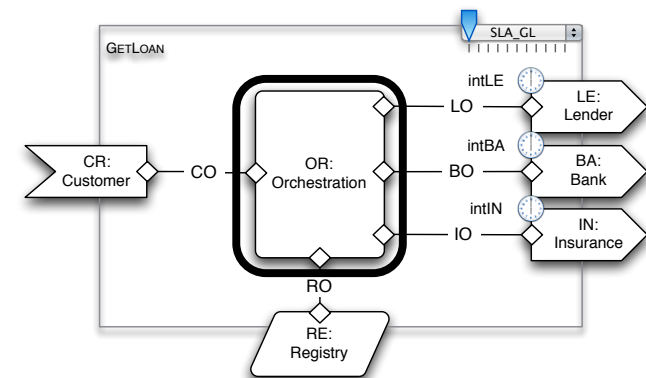
- The event `getProposal` (with a bell icon) is stored in OR's buffer, waiting to be processed
- Each component is associated to a PROCESSING RATE, which represents the delay of processing a received event
- All the events received by a specific component are affected by the same processing rate
- For example the delay of OR for processing events is $\text{processingRate}(\text{OR})$

Delays and dependencies



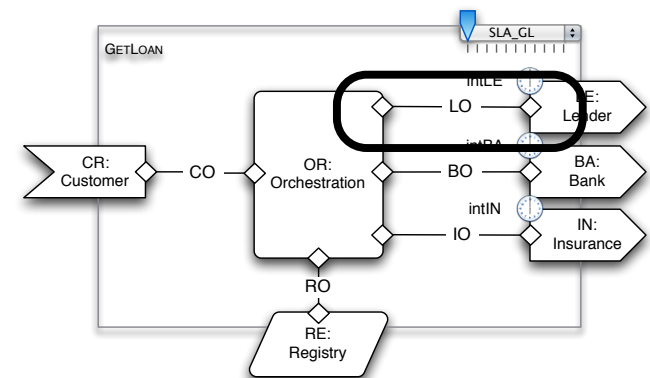
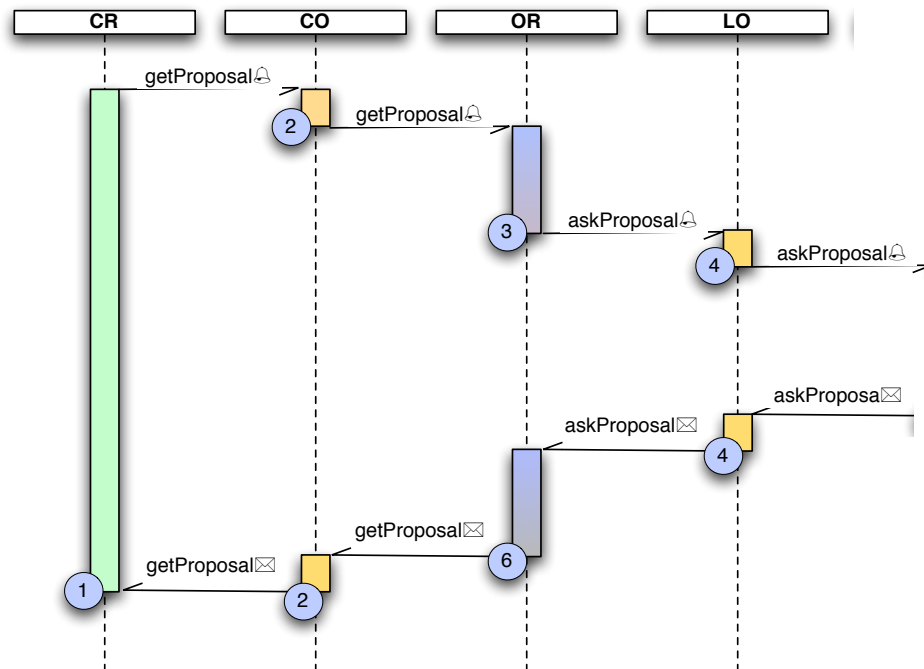
transition P1

triggeredBy getProposal
 guardedBy s=INITIAL
 effects s'=WAIT_PROPOSAL
 sends askProposal



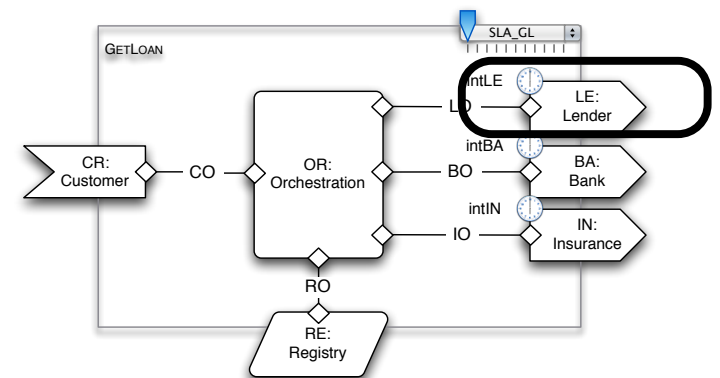
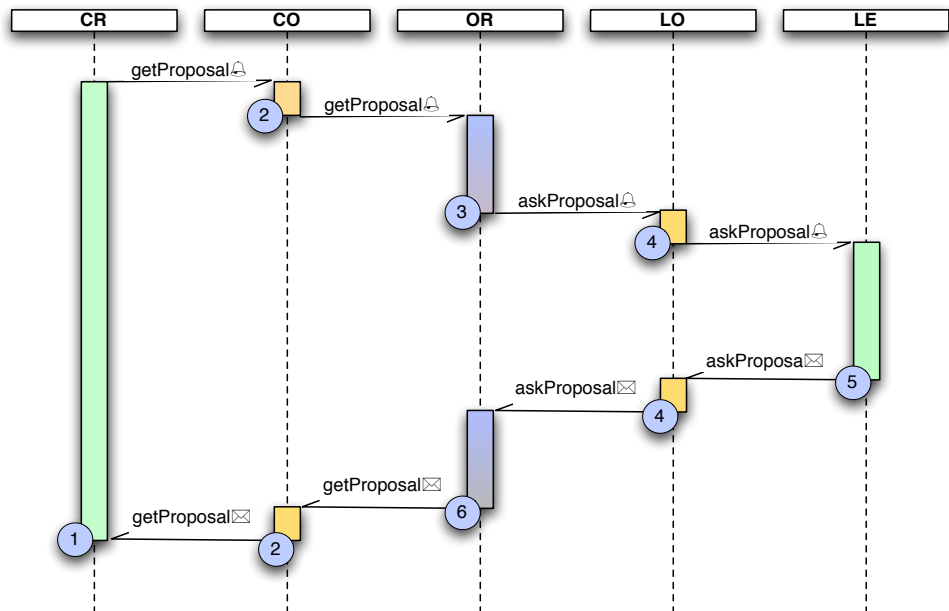
- When the event is processed it can be either executed or discarded
- The execution of an event corresponds to the execution of a transition.
- Each transition is associated to an EXECUTION RATE which represents the time taken to compute the reaction to an event, for example to execute transition P2
- The execution rate is different for any transition of the same component (e.g., $\text{executionRate}(\text{OR})(\text{P1})$)

Delays and dependencies



- Again there is a delay due to a wire
- (4) represents `transferRate(LO)`

Delays and dependencies



- when a requires interface is discovered and bound at run time, we have a delay which we call COMPOSITION RATE
- For example, the composition rate of LE is $\text{compositionRate}(\text{LE})$
- Also, each requires-interface is associated to a RESPONSE RATE representing the time taken by an external interface to reply to an event
- we associate a responseRate to every r&s (delay between getProposal -event and askProposal -event)
- (5) is $\text{compositionRate}(\text{LE}) + \text{responseRate}(\text{LE})(\text{askProposal}, \text{askProposal})$

PEPA Eclipse Plugin

- After having extended a SRML module with delays, associating one or more rate to each component, EX-I, and wire, we can encode the module into a PEPA term

The screenshot shows the Eclipse IDE with the PEPA Eclipse Plugin installed. The main editor displays the PEPA term for 'MyPEPA2.pepa'. The left sidebar shows the project structure, and the right sidebar shows the state space analysis results.

PEPA Term:

```
// rate specifications
startRate=1;
coTransferRate=1;
orProcessingRate=4;
orp1ExecutionRate=2;
orp2ExecutionRate=1.5;
clTransferRate=1.5;
leResponseTime=2;
//leCompositionRate=0.1;

CR=(crgetProposalB,startRate).CR1;
CR1=(crgetProposalE,T).CR;

T0 = (crgetProposalB,T).C01b + (orgetProposalE,T).C01e;
C01b=(orgetProposalB,coTransferRate).C0;
C01e=(crgetProposalE,coTransferRate).C0;

OR=(orgetProposalB,T).OR1;
OR1=(p1,orProcessingRate).OR2;
OR2=(oraskProposalB,orp1ExecutionRate).OR3;
```

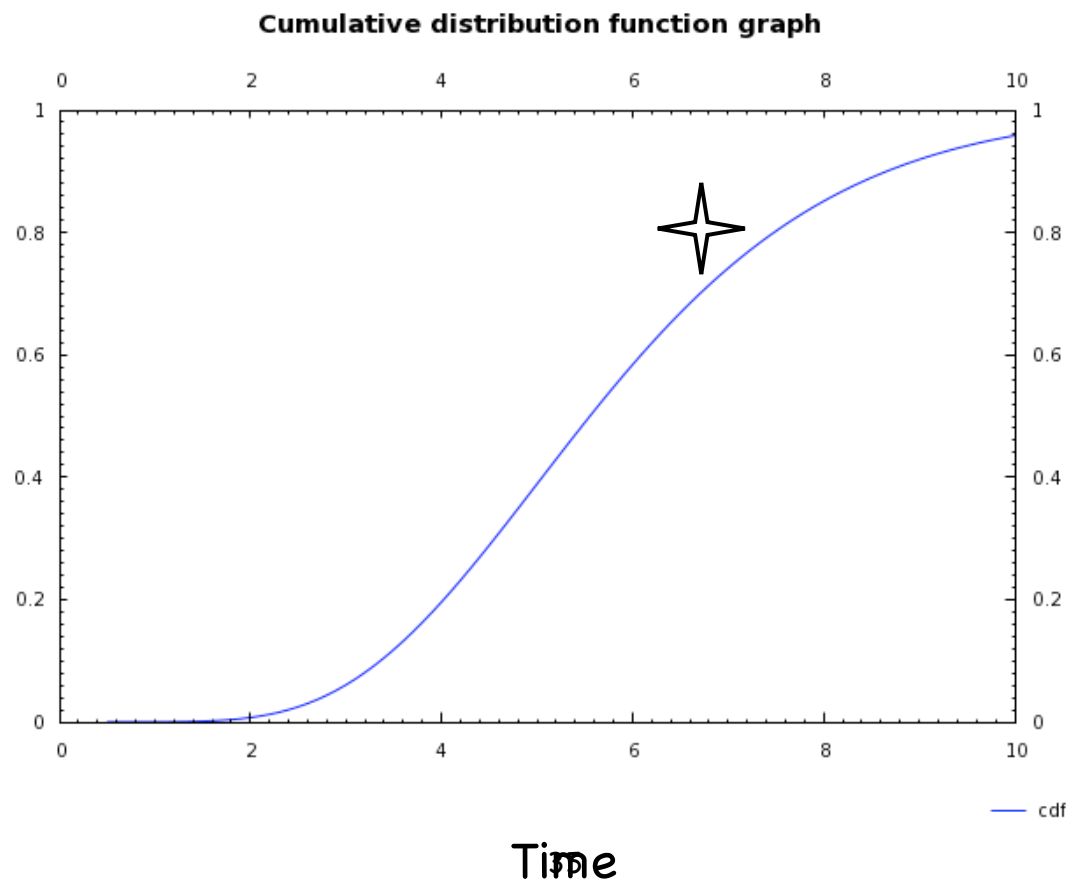
State Space Analysis Results:

State	Component	Rate	Local States
5	CR1 OR3 LE CO	CL1b	0.09756097560975609
6	CR1 OR3 LE1 CO	CL	0.07317073170731707
7	CR1 OR3 LE CO	CL1e	0.14634146341463414
8	CR1 OR4 LE CO	CL	0.036585365853658534
9	CR1 OR5 LE CO	CL	0.09756097560975609
10	CR1 OR LE	CO1e	0.14634146341463414

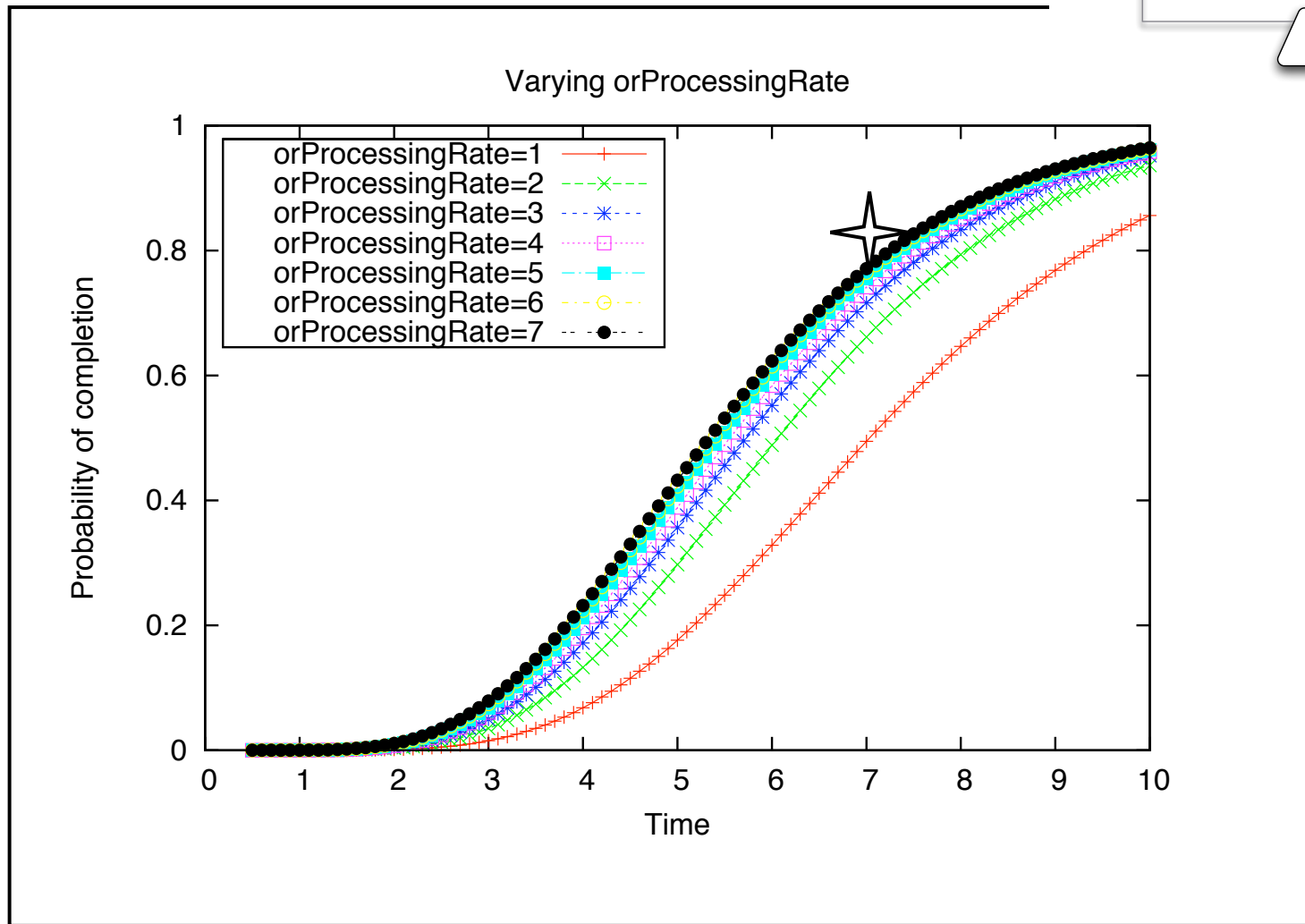
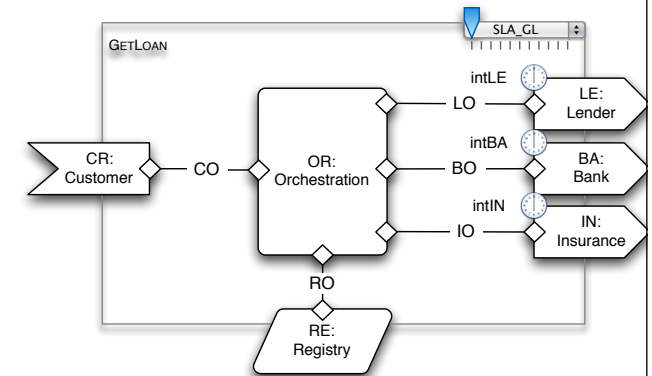
Passage Time Analysis

- Using the Passage Time Analysis we analyse for which rates the following holds: "In 80% of the cases, the delay between `getProposal`🔔 and `getProposal`📧 has an upper bound of 7s".

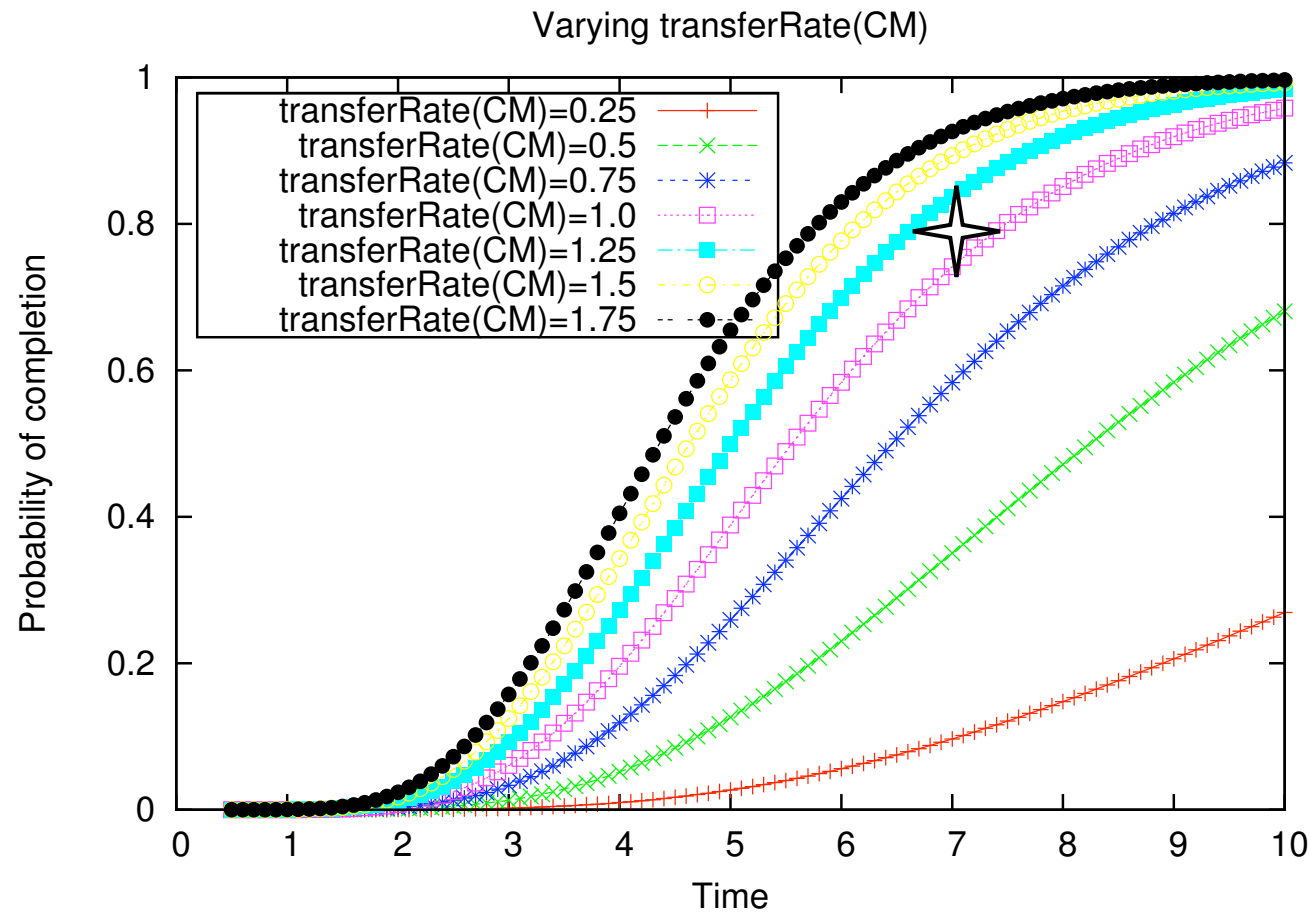
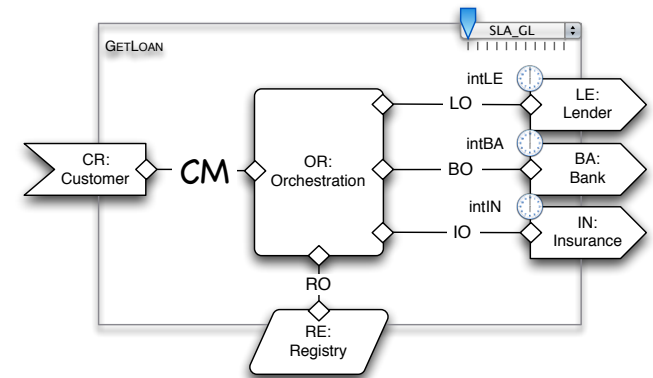
Probability



Passage Time Analysis



Passage Time Analysis



A view of ensemble

use case diagrams

from scratch
use case -> SRML

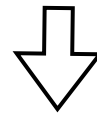
from scratch
CBD

from scratch
reuse
refinement

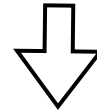
PEPA
Logics of Interactions

SCA

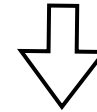
Define requirements



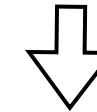
Define modules + EX-Is



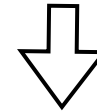
Define internal structure



Define behaviours



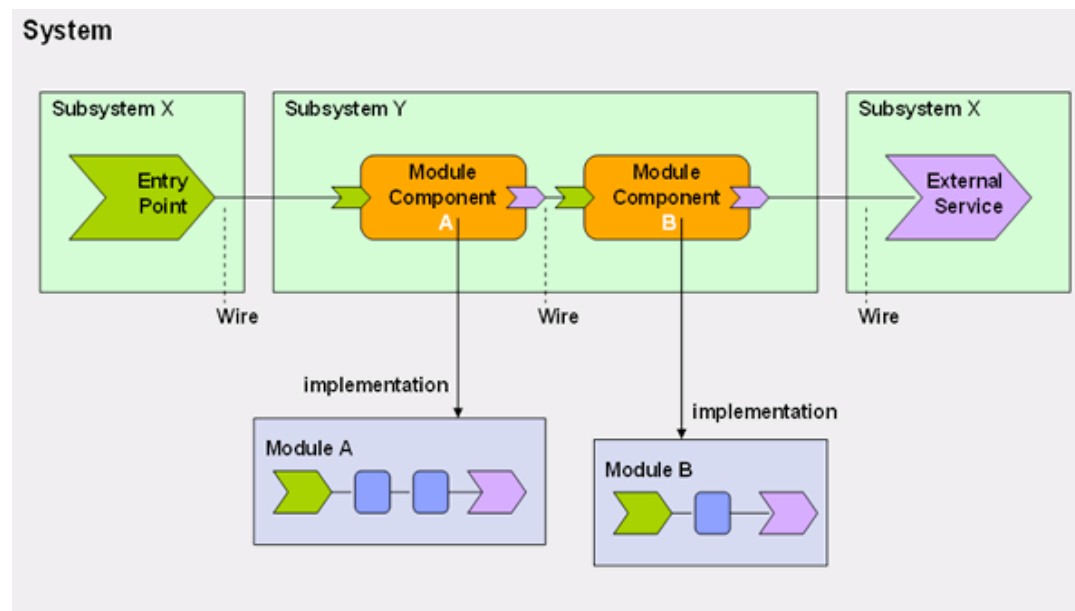
Analyse



Deploy

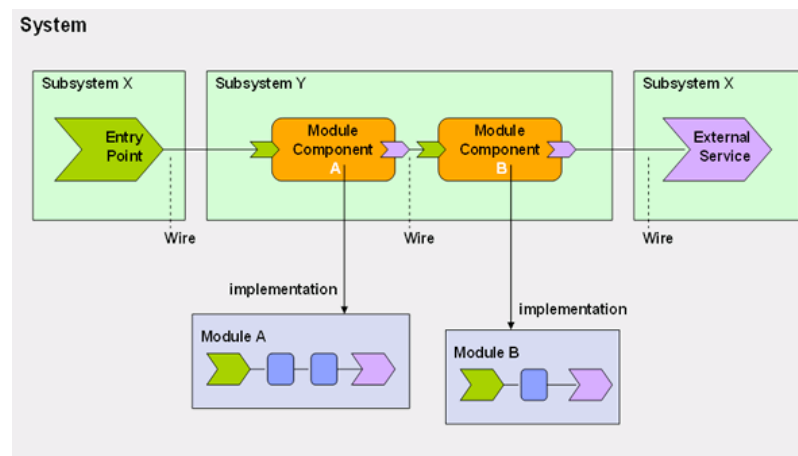
SCA

- The Service Component Architecture (SCA) is a recent set of specifications, proposed by an industrial consortium including major vendors like IBM, ORACLE, BEA, etc.
- SCA describe a middleware-independent model for building over SOAs.
- Similarly to SCA, SRML provides primitives for modelling, in a technology agnostic way, business processes as assemblies of
 - (1) tightly coupled components that may be implemented using different technologies (including wrapped-up legacy systems, BPEL, Java, etc.)
 - (2) loosely coupled, dynamically discovered services.

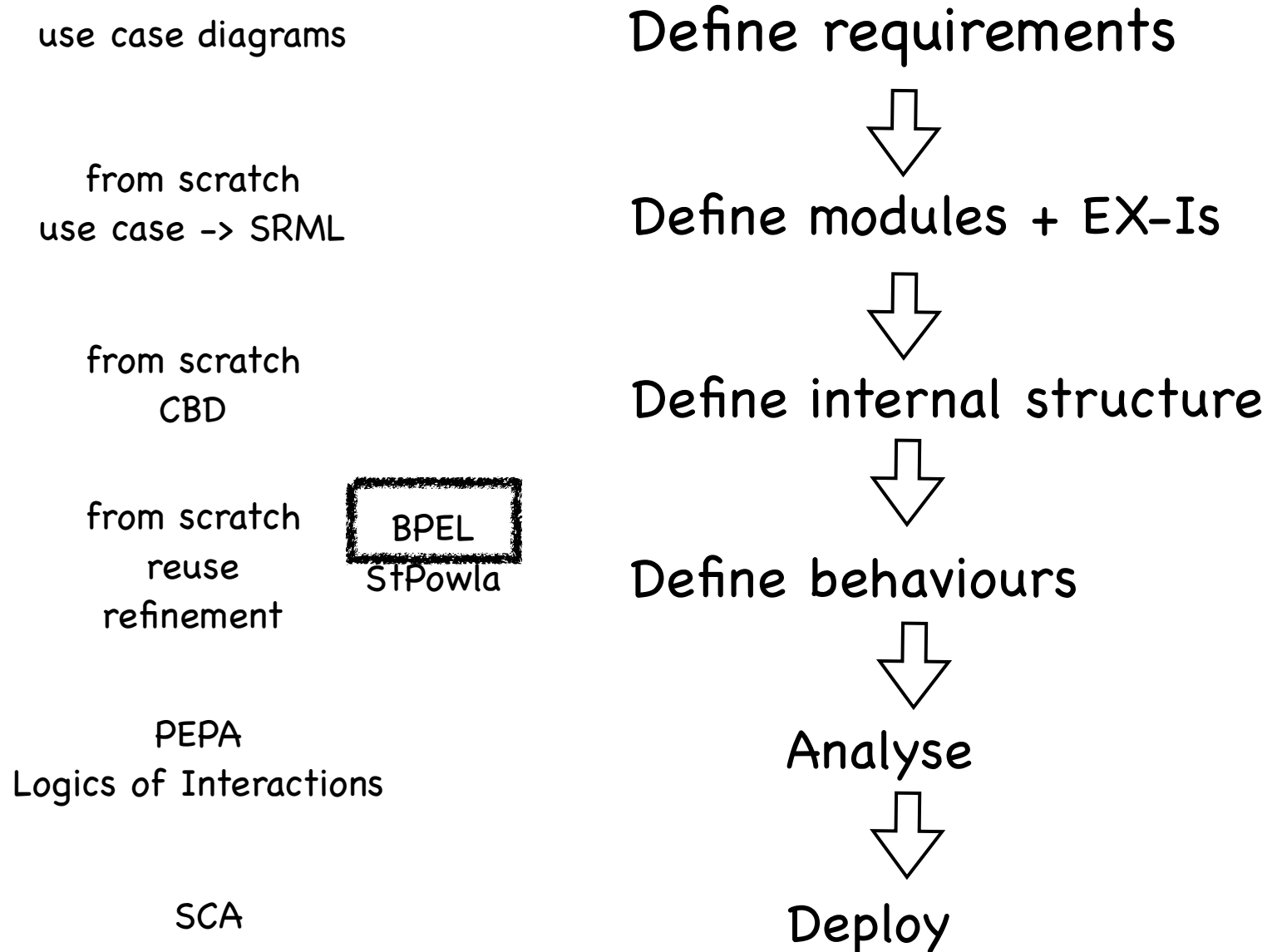


SCA

- ⦿ Differently from SRML, SCA is not a modelling language but a framework for modelling the structure of a service-oriented software artefact and for its deployment.
- ⦿ SCA abstracts from the business logic provided by components in the sense that it does not provide a means to model the behavioural aspects of services.
- ⦿ SRML is, instead, a modelling language that
 - ⦿ provides the primitives to specify such behavioural aspects.
 - ⦿ relies on a mathematical framework for reconfiguration, behavioural interfaces and SLA



A view of ensemble



BPEL2SRML



From BPEL to SRML: a formal transformational approach

Bocchi, Hong, Lopes and Fiadeiro, WSFM 2008

<http://www.cs.le.ac.uk/people/jfiadeiro/Papers/BPEL2SRML.pdf>

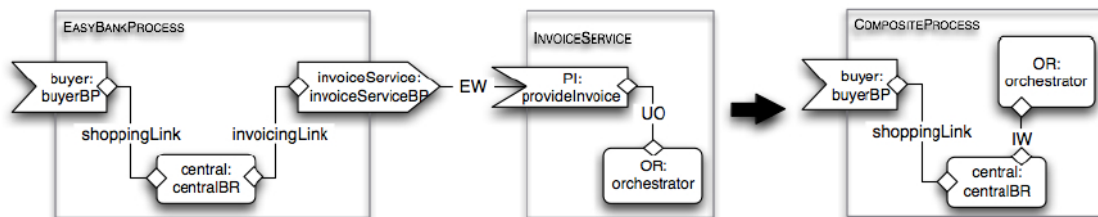
- WS-BPEL is an OASIS specification for defining business processes of Web services
- We defined an encoding to extract SRML models out of existing BPEL processes
- In fact we defined an encoding of (part of) BPEL into SRML
- Aspects of session/fault/configuration management still have to be added

BPEL Tag/Construct	Tool	Encoding
Invoke , Receive, Reply, Assign (BA)	✓	✓
Wait, Empty, Exit (BA)	✗	✓
Throw (BA)	✗	✗
Sequence, Switch (SA)	✓	✓
Flow, While (SA)	✗	✓
Control Links, Scopes, Correlation Sets	✗	✗

Model Extraction

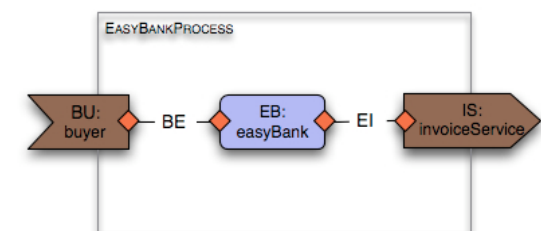
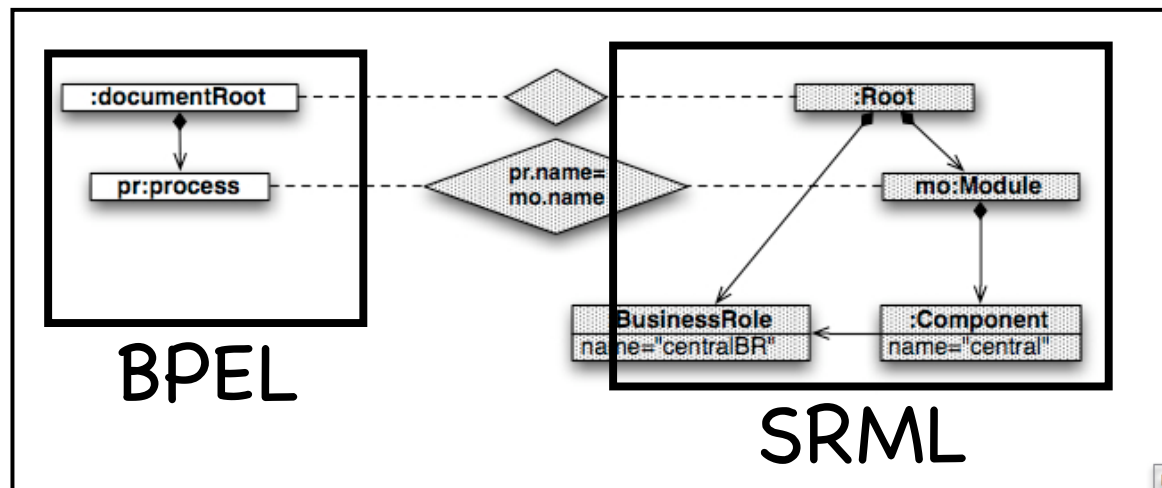
BPEL and SRML: Advantages

- The main aim is not to provide BPEL with a formal semantics
- The aim is
 - to enable extraction of models,
 - to provide a library of models (SRML components),
 - to allow the models deriving from existing processes to be used to define other SRML models. The components may be more than one and they can be
 - derived by BPEL processes,
 - defined from scratch,
 - derived from any language for which an encoding into SRML exists.
- To allow the ensemble to be analysed within one formal framework.



The encoding BPEL2SRML

- The encoding has been done as follows...
 - EMF tree for WSDL/BPEL derived from XSD and Eclipse BPEL project
 - EMF tree for SRML being refined (while implementing the SRML editor)
 - Design of transformation rules (structure and behaviour)



- A BPEL process is encoded into a SRML module with one component

The encoding BPEL2SRML

- For every transition A we define a variable ra ("a is ready") and fa ("a has finished"), which are initially false.
- The encoding of a simple activity (receive)

<pre><receive partnerLink="pl" portType="pt" operation="op" variable="v" createInstance=.../></pre>		<pre>transition transition_A triggeredBy pt.op? guardedBy ra $\wedge$ $\neg$exit effects $\neg ra'$ $\wedge$ fa' $\wedge$ v.pl'=pt.op.v.pl $\wedge$... $\wedge$ v.pn'=pt.op.v.pn</pre>
---	--	--

- The encoding of a structured activity (sequence)

```
<sequence
  name="X">
  activity A
  activity B
</sequence>
```

```
transition transition_X
triggeredBy (rx  $\vee$  fa  $\vee$  fb)  $\wedge$   $\neg$ exit
guardedBy (rx  $\vee$  fa  $\vee$  fb)  $\wedge$   $\neg$ exit
effects (rx  $\supset$  ra'  $\wedge$   $\neg$ rx)  $\wedge$  (fa  $\supset$  rb'  $\wedge$   $\neg$ fa')  $\wedge$  (fb  $\supset$  fx'  $\wedge$   $\neg$ fb')
```

The transition is executed the first time when X is triggered

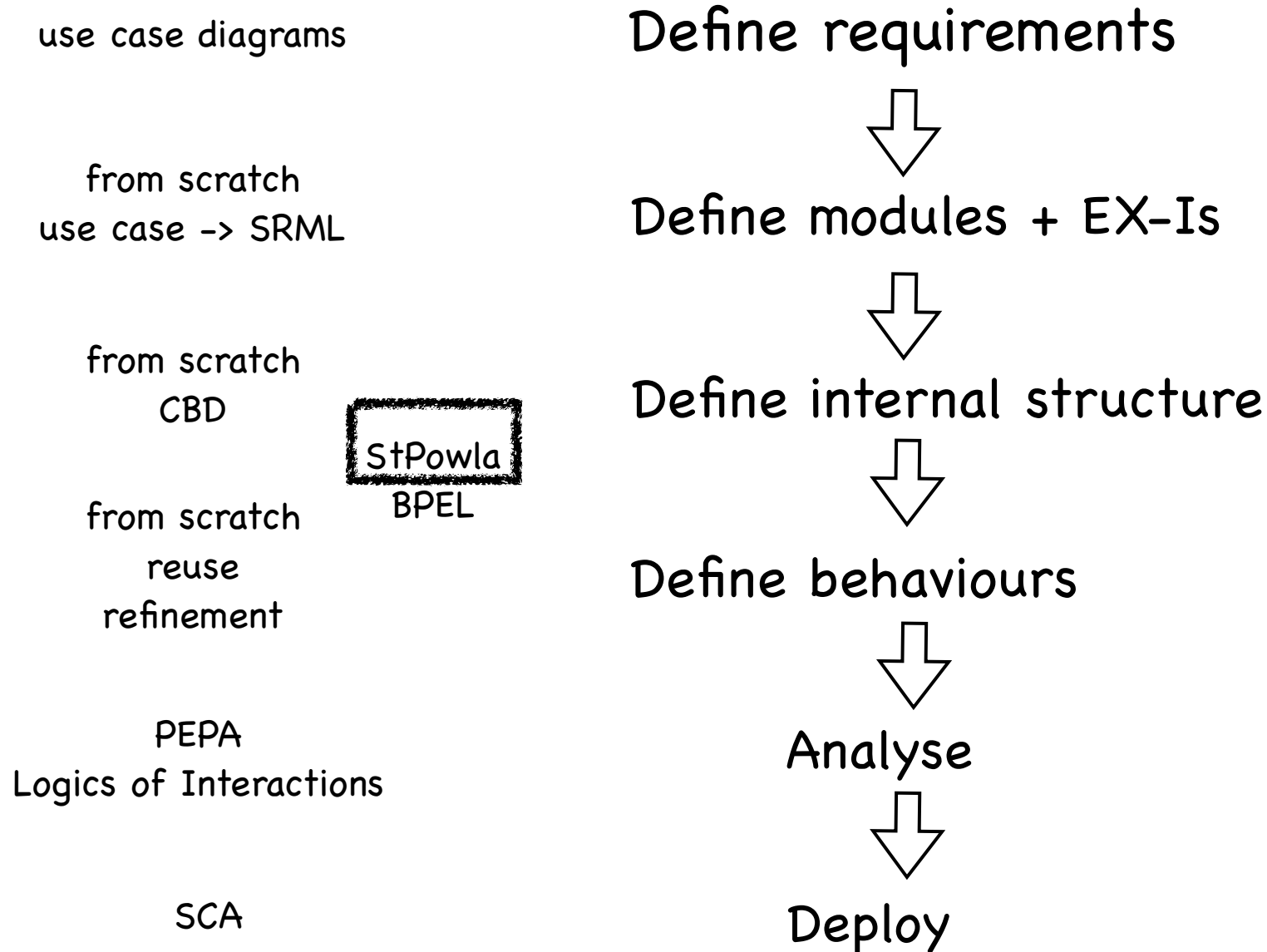
The transition is executed the second time when A terminates

The effect is to enable the execution of A

The effect is to trigger B

When the last activity terminates, the end of X is notified

A view of ensemble



StPowla2SRML

StPowla

[S Gorton, C Montangero, S Reiff-Marganiec and L Semini. StPowla: SOA, Policies and Workflows. WESOA 2007]

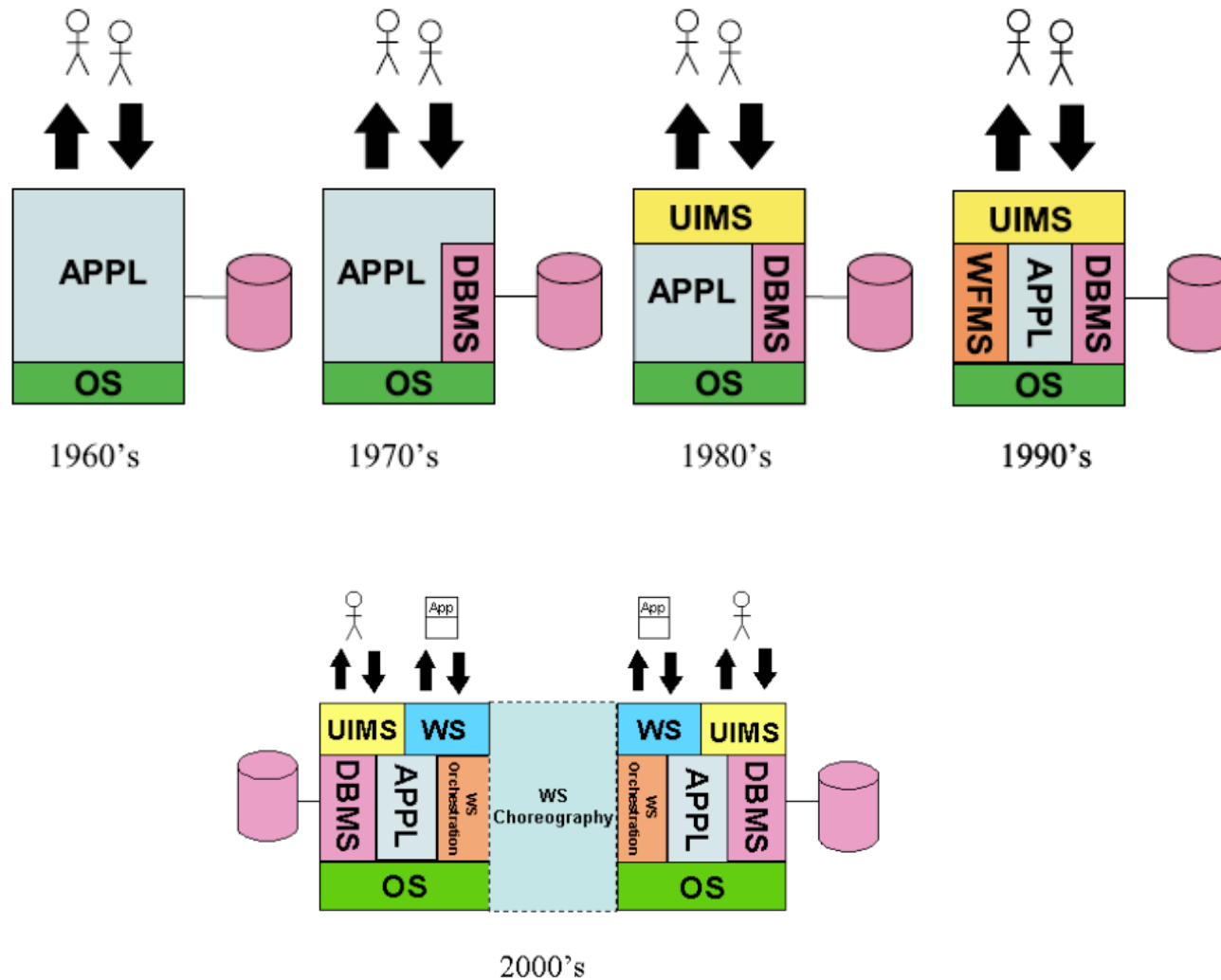
is a service-targeted, policy-oriented, workflow approach

- workflows
- reconfiguration through policies
- StPowla has been encoded into SRML in order to
 - provide StPowla with a formal framework
 - add a higher level of modelling in SRML



Engineering Service Oriented Applications: From StPowla Processes to SRML Models
Bocchi, Gorton, Reiff-Marganiec, FASE 2008
<http://www.cs.le.ac.uk/people/srm13/publications/fase08.pdf>

A bit of history...



Workflow

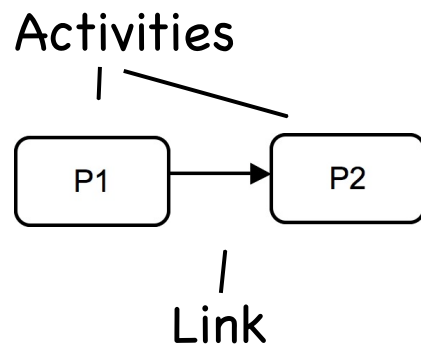
- ☉ The automation of a business process, in whole or part, during which documents, information or tasks are passed from one participant to another for action, according to a set of procedural rules.

Business Process

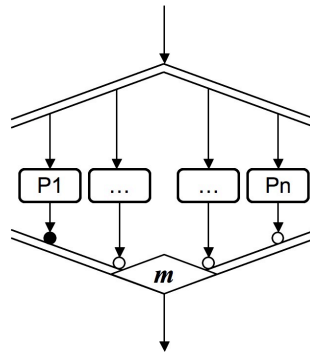
- ☉ A set of one or more **linked** procedures or **activities** which collectively realise a business objective or policy goal, normally within the context of an organisational structure defining functional roles and relationships.

from WPMC (Workflow Management Coalition) Glossary
(<http://www.aiai.ed.ac.uk/project/wfmc/ARCHIVE/DOCS/glossary/glossary.html>)

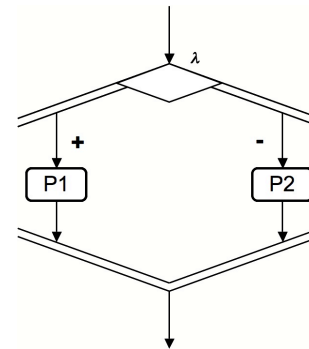
A sample notation...



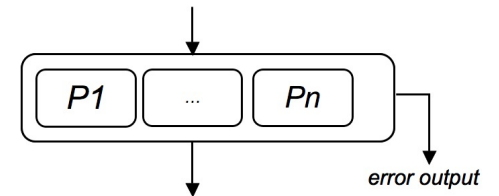
sequence



split join

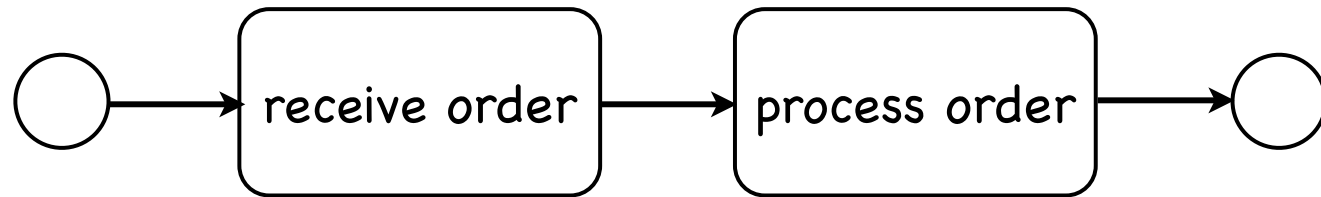


condition join



random preference

Example



- 🕒 Business Modelling

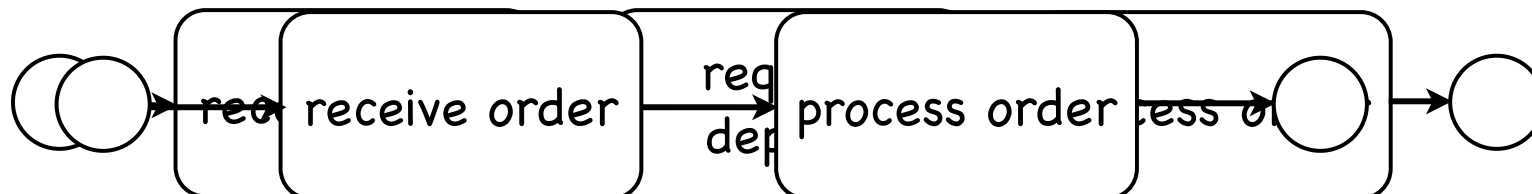
- 🕒 concerned with the ordering of tasks in an execution model
- 🕒 kept at a high-level for end-users to create

Business Policies

- There exists a need for flexibility: customization of a core model to handle variability in domain
- Business models are subjected to overarching constraints (e.g., business rules, global or enterprise constraints) expressed as business policies

```
getDepositIfLargeOrder  
  appliesTo receiveOrder  
  when task_completion  
    if receiveOrder.orderValue > £250000 do  
      insert(requestDeposit, receiveOrder, true)
```

request
deposit

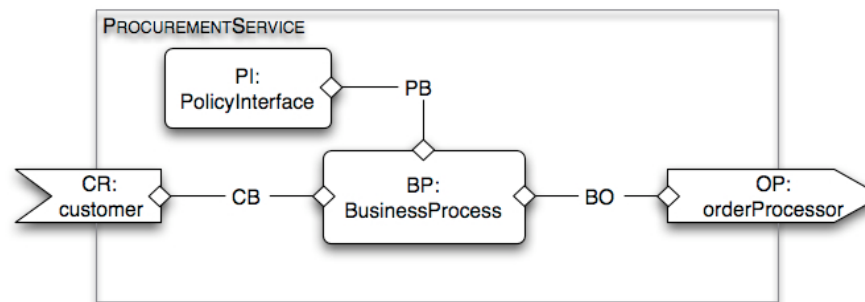


StPowla: Reconfiguration Functions

- ⑤ - fail() -
declares the current task to have failed (i.e., discard further task processing and generates the task_failure event)
- ⑤ - abort() -
aborts the current task and progresses to the next task, generating the task_abort event
- ⑤ - block(s,p) -
waits until predicate p is true before commencing scope s
- ⑤ - insert(x,y,z) -
inserts task or scope y into the current workflow instance after task x if z is true, or in parallel with x if z is false
- ⑤ - delete(x) -
deletes task/scope x from the current workflow instance

StPowla2SRML

- ⑤ The internal structure of the SRML module is organised in two components: one implementing the business process and one implementing the policy interface
- ⑤ The policy interface determines when a policy requires a reconfiguration and notifies the business process component



- ⑤ BP has one interaction for each of the reconfiguration functions...

BUSINESS ROLE BusinessProcess is

INTERACTIONS

r&s delete[i:natural]

🔔 task:taskId

r&s insert[i:natural]

🔔 task:taskId

newTask:taskId

c:condition

...

r&s fail[i:natural]

🔔 task:taskId

r&s abort[i:natural]

🔔 task:taskId

r&s block[i:natural]

🔔 task:taskId

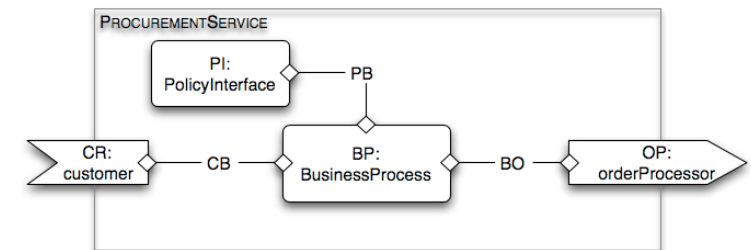
c:condition

StPowla2SRML

- The business role of BP has one or more transition that model the reaction to each reconfiguration function
- The the transitions for the delete task reconfiguration function are presented below:

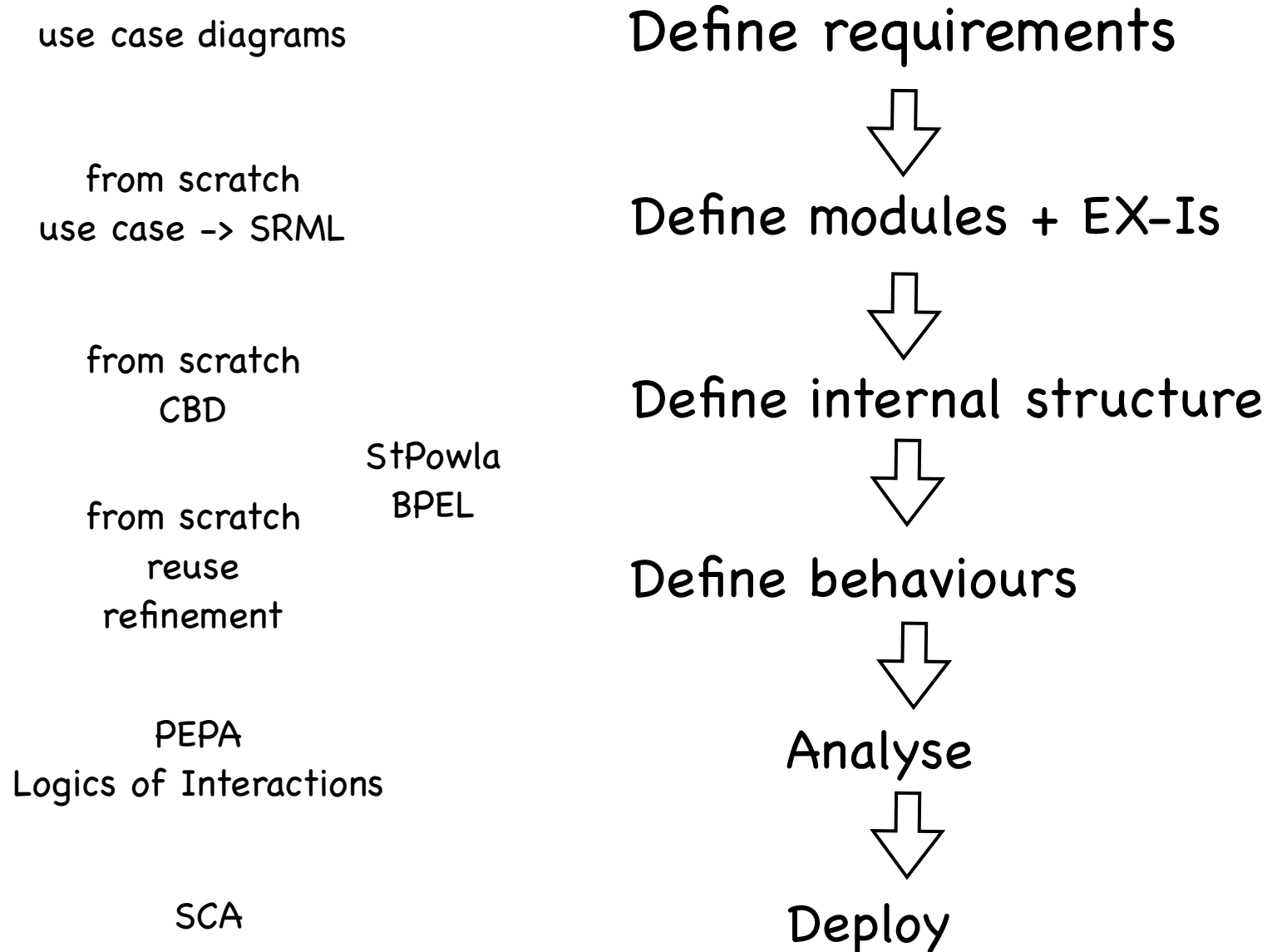
```
transition policyHandler_delete_1  
  triggeredBy delete[i]🔔?  
  guardedBy state[delete[i]🔔.task]=toStart  
  effects policy[delete[i]🔔.task]'
```

```
transition policyHandler_delete_2  
  triggeredBy start[x]  
  guardedBy P_delete[i]🔔? ∧ P_delete[i]🔔.task=x  
  effects ¬start[x]' ∧ done[x]' ∧ state[x]'=done  
  sends delete[i]📧!
```



- The encoding of the workflow constructs is similar to the one of the encoding from BPEL

A view of ensemble



A view of ensemble

