



Graph Transformation: Foundations and Applications to Software Engineering

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Based on book with Gabriele Taentzer Philipps Universität Marburg, Germany

Graphs and Graph Transformations

Graphs represent

- Class and object structures
- Processes and architectures
- Diagrams, 3D objects

formally, visually and directly (little need for encoding)

Graph transformations model

- Behaviour and reconfiguration
- Creation, manipulation, evolution
- Operational semantics

in an *intuitive* rule-based way, with *formal semantics* for *execution* and *analysis*

Outline

Foundations

- Graphs
- Graph Transformations
- Semantics and Control
- Analysis Techniques

Two Types of Applications

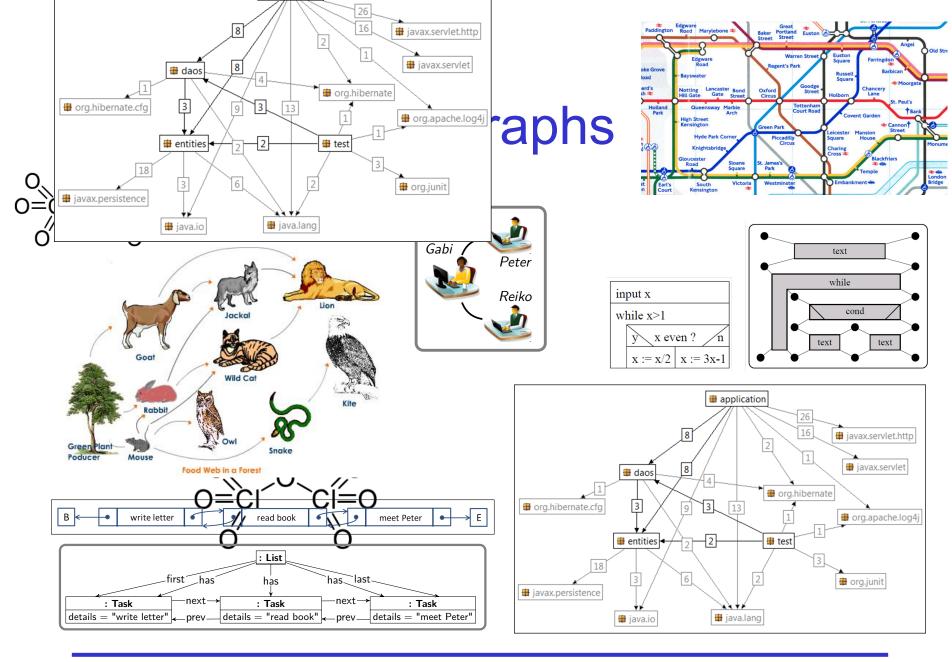
Graph Transformation-Based

- Software Engineering
 - → GTBSE
- Language Engineering
 - → GTBLE

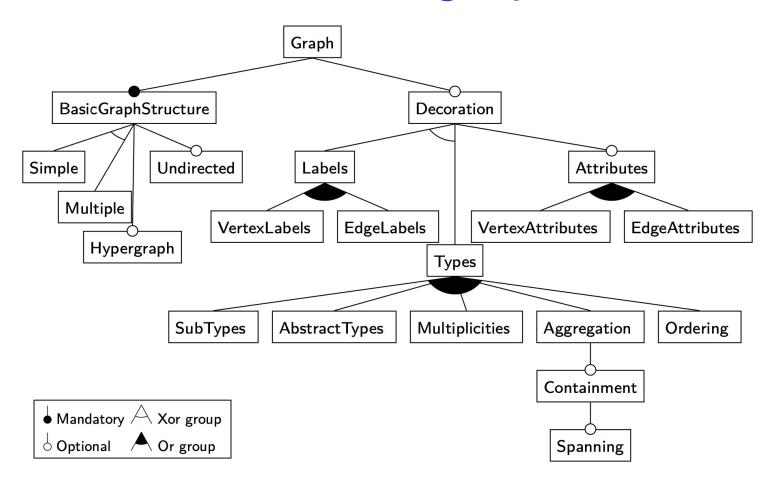


Part 1: Foundations

- Graphs
- Graph Transformations
- Semantics and Control
- Analysis



Variants of graphs

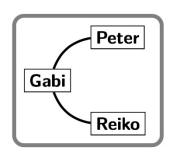


GTBSE or GTBLE?

Directed or undirected Labelled or typed

Undirected labelled graph

- Nodes set V, Edges $E \subseteq P_2(V)$
- Labels L = {Peter, Gabi, Reiko}
- lab_V: V → L

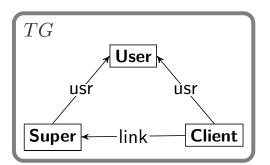


c1: Client —usr→peter: User link s: Super —usr→gabi: User link c2: Client —usr→reiko: User

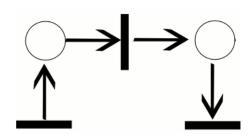


Directed typed graph

- Node, edge labels form type graph TG
- Instance graph G with homomorphism type preserving graph structure



GTBLE: Petri nets



Create

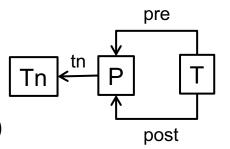
- Type graph for Petri nets
- Instance graph for the example

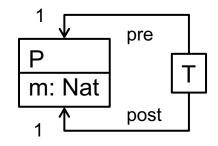
Design decisions

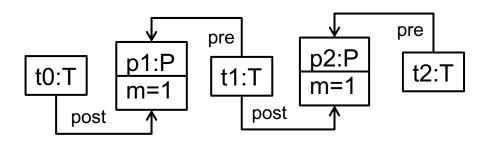
- Bipartite (or unordered hyper) graphs as typed graphs
- Node type Tn to represent tokens (left) vs attribute m for marking (right)

Marked graphs (right)

 add constraints for singleton pre and post sets







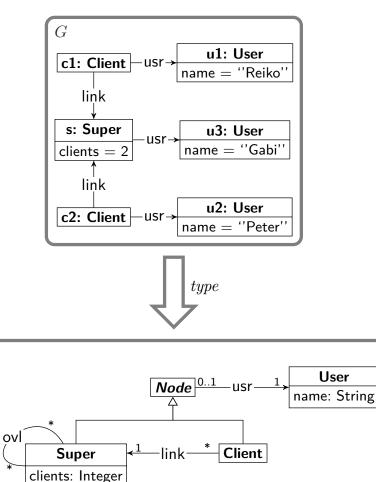
Typed attributed graphs

with subtyping

Structured P2P model

- Client and Super Nodes serve Users
- Super nodes form overlay network, support limited number of Clients

→ GTBSE



TG

Formalise this, ...

Type, instance graphs

- Graphs and graph homomorphisms
- → Slice category

Attributed graphs

- Graphs related to fixed data algebra
- → Symbolic attributed graphs



Part 1: Foundations

- Graphs
- Graph Transformations
 - Basic rules and transformations
 - Global application conditions
 - Advanced features
- Semantics and Control
- Analysis
- Henshin Tool Demo

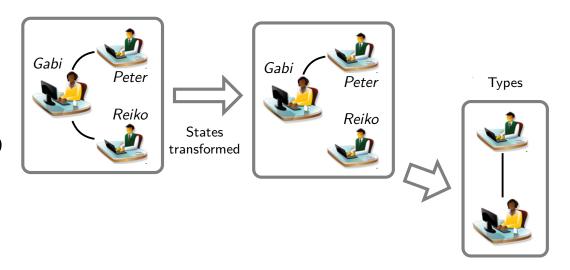
Rules generalize transformations

Specify changes by

- Difference between pre and post graphs
 - → Deleted: L\R
 - → Created: R \ L
- Context required for the changes to happen
 - \rightarrow Preserved: L \cap R

Rule should describe

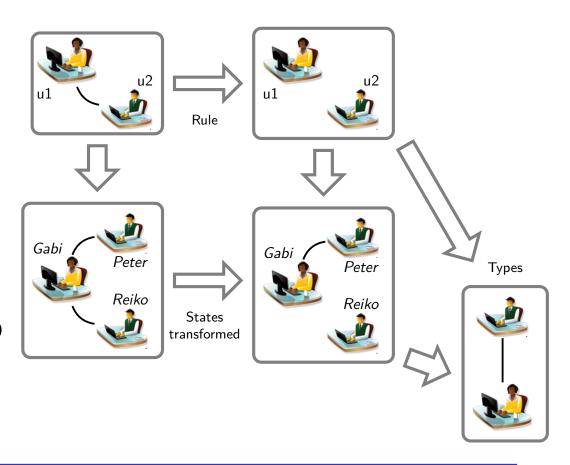
- What changes, and how?
- In which situations?



Rules generalize transformations

Specify changes by

- Difference between pre and post graphs
 - → Deleted: L\R
 - → Created: R \ L
- Context required for the changes to happen
 - \rightarrow Preserved: L \cap R

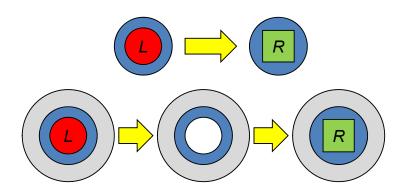


Two ways to Embouropiacoment is graph into context

Gluing approaches

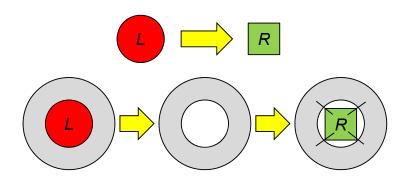
Interface shared between L and R provides gluing points

our choice



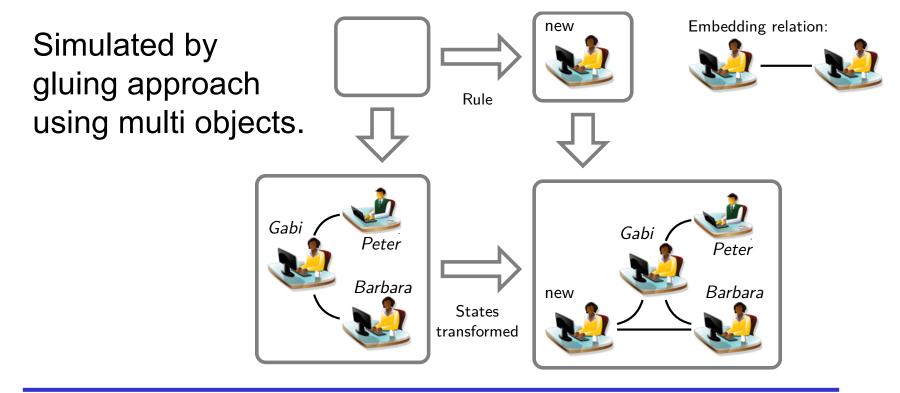
Connecting approaches

R is linked to context by edges created according to embedding relation

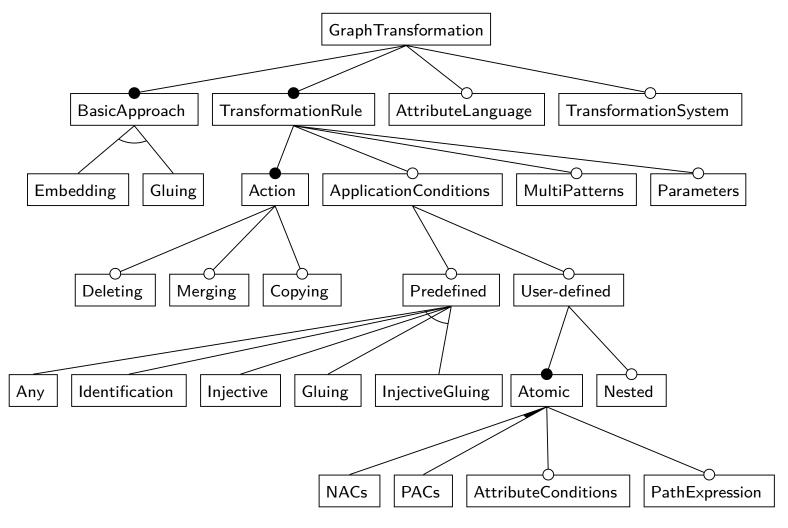


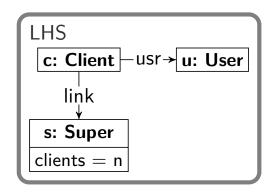
Example: Connecting Approach

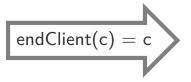
More powerful, but mathematically less tractable.

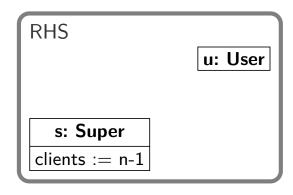


Graph Transformation Concepts





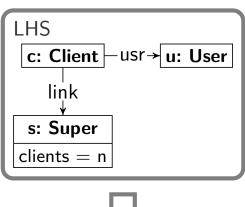


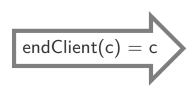


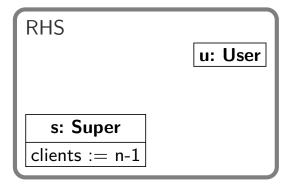
Rule Features

- Left- and right-hand side
- UML-like object notation
 - c:Client, u: User, s: Super
- Attribute conditions
 - clients = n
- Attribute assignments
 - clients := n-1

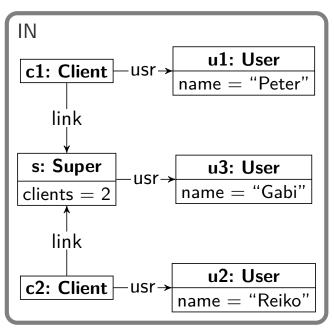
- Rule signature
 - endClient(c: Client) : Client
- Rule label
 - endClient(c) = c







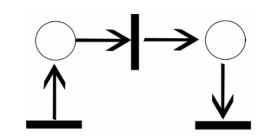




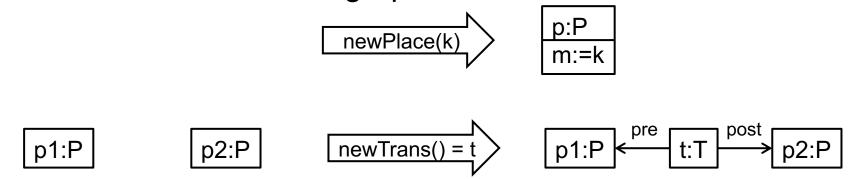
What should be the result?

- Elements of IN deleted?
- New elements created?
- Attribute values changed?

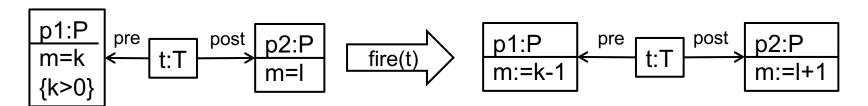
GTBLE: Rules for Nets



Rules to create marked graphs



Firing rule for marked graphs



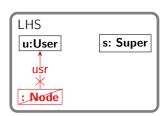
GTBSE: P2P Rules

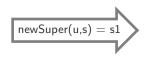
Network formation

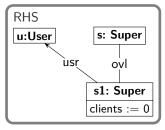
- Node creation and promotion
- Linking and redirecting clients

Rule Concepts

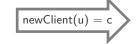
- Creation
- Deletion
- Parameters
- NACs
- Attribute conditions



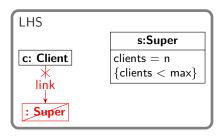






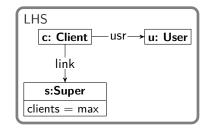




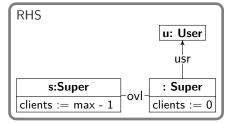


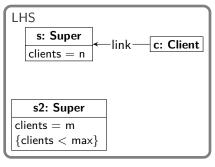




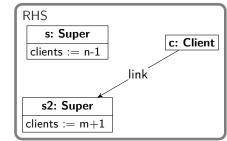












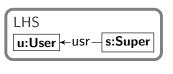
GTBSE: P2P Rules

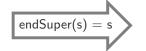
Network deconstruction

- Disconnect nodes
- Terminate nodes

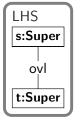
Rule Concepts

- Creation
- Deletion
- Parameters
- NACs
- Attribute conditions

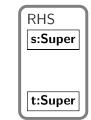


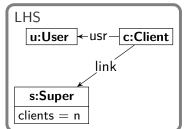


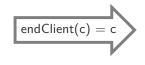


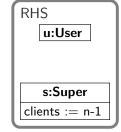


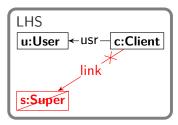










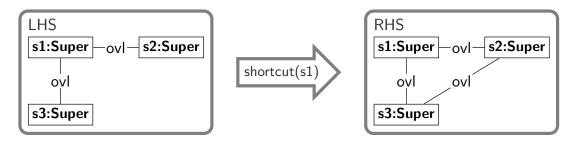






Global application conditions

What could possibly be wrong with this rule?

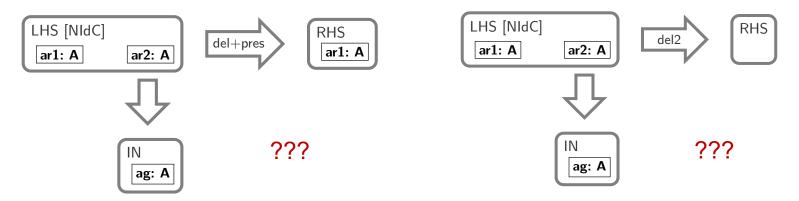


→ With a match mapping s2 and s3 to the same node in the input graph, the rule would create a loop.

Injective matches: Each element in the left-hand side is mapped to a separate element in the input graph.

Global application conditions: Identification condition

Allowing for non-injective matches, what happens here?



Identification condition: elements to be deleted are kept apart from

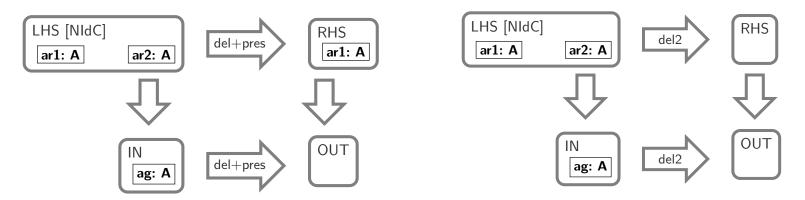
other elements

and

each other

Global application conditions: Identification condition

Allowing for non-injective matches, what happens here?



Identification condition: elements to be deleted are kept apart from

other elements

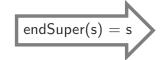
and

each other

Global application conditions: Dangling condition

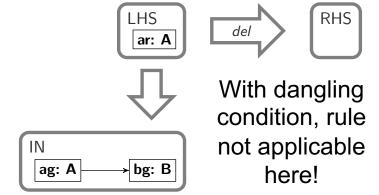
Deletion in unknown contexts

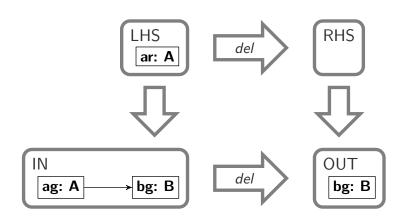




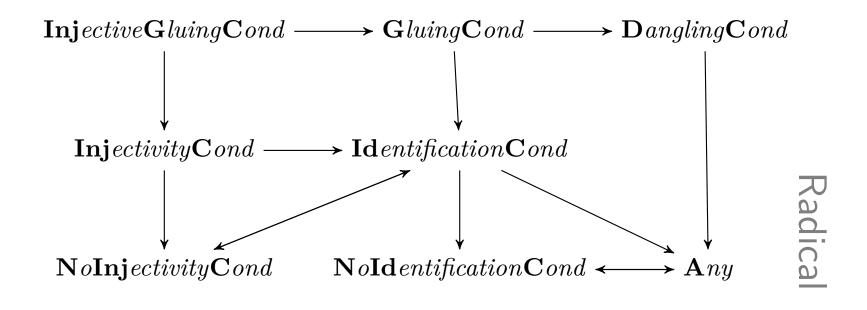


What happens to any connections of s with other Super nodes?





Global application conditions: Overview



Formalise this, ...

Type, instance graphs

- Graphs and graph homomorphisms
- → Slice category

Rules, transformations

- → Rules as spans or partial morphisms
- →DPO, SPO, SqPO

Attributed graphs

- Graphs related to fixed data algebra
- → Symbolic attributed graphs

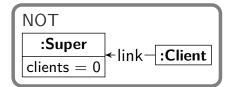
Advanced features: Graph constraints

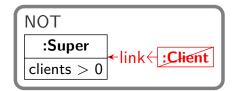
"The *clients* attribute of a *Super* node is 0 iff there are no *Client* nodes connected to it."

Logically

self.clients=0 iff self.link->isEmpty()

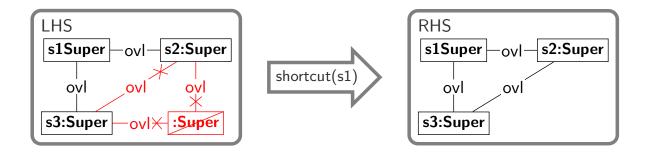
or graphically





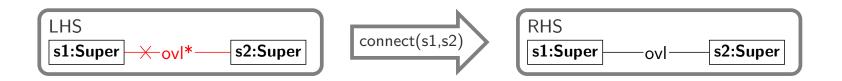
Advanced features: Negative application conditions

"Apply shortcut rule only if nodes *s2* and *s3* are not connected, neither directly nor via a 3rd node."



Advanced features: Path Expressions

"Apply connect only if there is no path of *ovl* edges between *s1* and *s2*."



Advanced features: Multi Objects and Patterns

LHS [DC]

With the dangling condition, we cannot delete a *Super* node without explicitly deleting all link

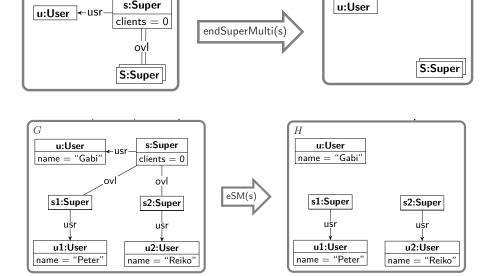


RHS

without explicitly deleting all links to all its peers.

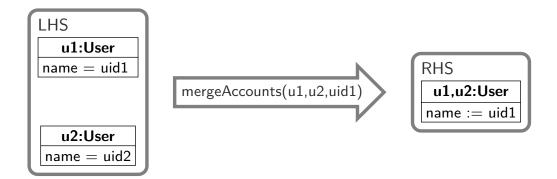
This is possible using a rule with multi object.

S:Super matches the set of all super nodes linked to s: {s1, s2}



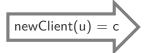
Advanced features: Merging

"Merge two user accounts into one."



Integrated Rule Notation



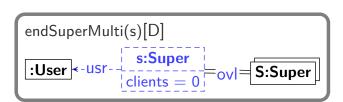


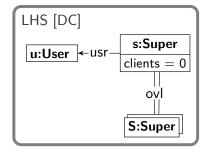


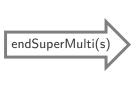
Integrate left- and right-hand side into a single graph

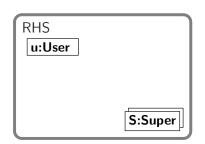
```
newClient(u) = c
\vdots \text{Node} \rightarrow usr \rightarrow u: User \leftarrow usr \leftarrow usr \leftarrow c: Client \\ \vdots \\ \leftarrow usr \\ \vdots \\ \leftarrow usr \leftarrow
```

 Use colours and labels to distinguish different roles

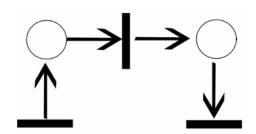




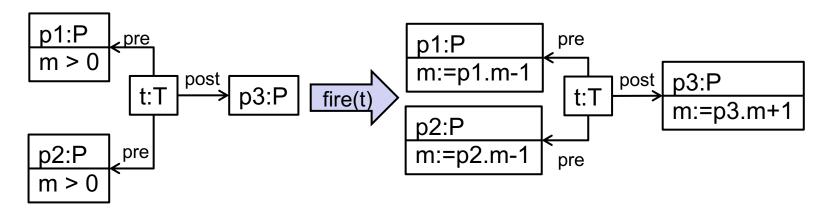




GTBLE: Firing rules for general Petri nets

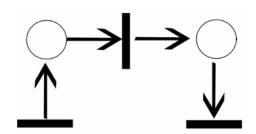


Rule for 2:1 transitions (2 places pre, 1 place post set)

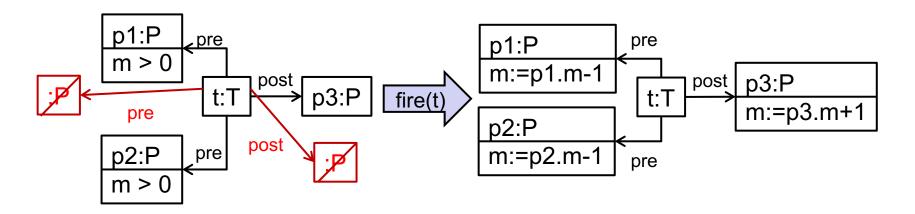


How to ensure this is only applicable to 2:1 transitions?

GTBLE: Firing rules for general Petri nets



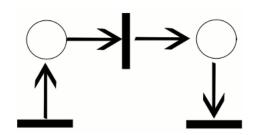
Rule for 2:1 transitions (2 places pre, 1 place post set)



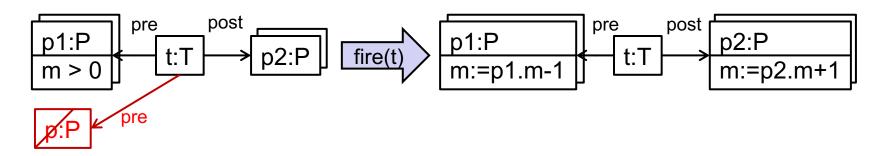
How to ensure this is only applicable to 2:1 transitions?

→ Only matches if there are exactly 2 pre and 1 post places.

GTBLE: Firing rule for general Petri nets



Firing rule for arbitrary transitions



Matches sets of

- {p1:P | p1.m > 0 and p1 in pre set of t}
 - NAC ensures there is no p:P in pre set not in that set
- {p2:P | p2 in post set of t}

Updates attributes m of all matched P nodes

Formalise this, ...

Type, instance graphs

- Graphs and graph homomorphisms
- → Slice category

Attributed graphs

- Graphs related to fixed data algebra
- → Symbolic attributed graphs

Rules, transformations

- → Rules as spans or partial morphisms
- →DPO, SPO, SqPO

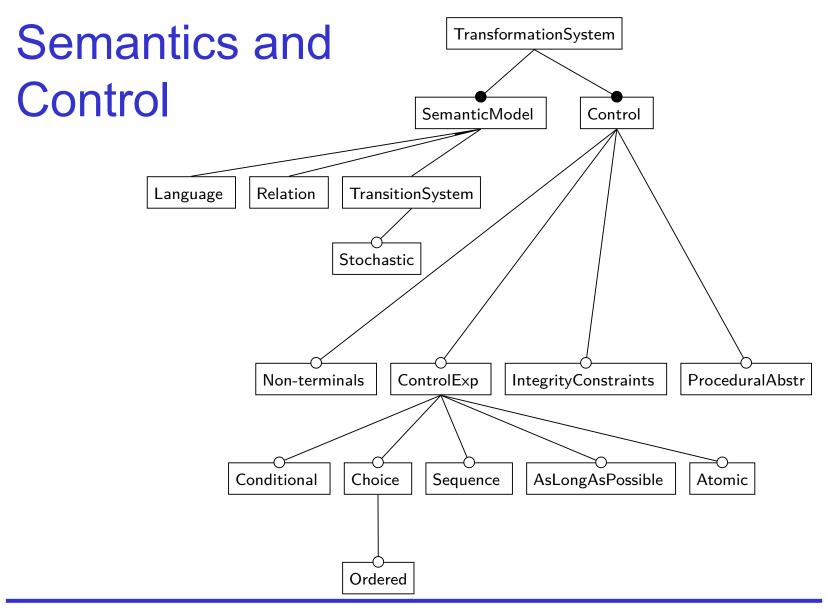
NACs, multi-objects

- → Forbidden matches
- → Amalgamation



Part 1: Foundations

- Graphs
- ✓ Graph Transformations
- Semantics and Control
 - Graph transformation systems, grammars, and derivations
 - Graph languages, relations and transition systems
 - Rule-based programming: textual and diagrammatic
- Analysis
- Henshin Tool Demo



Transformation systems, grammars, and semantic models

Graph transformation system GTS = (TG, R)

Graph grammar $GG = (TG, R, G_0)$

E.g. type graph and rules of the P2P model with start graph as initial network configuration

Semantic Models

- Language: set of all graphs reachable from G₀
- Relation: mapping from input to output graphs
- Transition system: reachable graphs as states, transformation steps as transitions

Relation: Compute Spanning Tree of

Super Nodes

Extended type graph

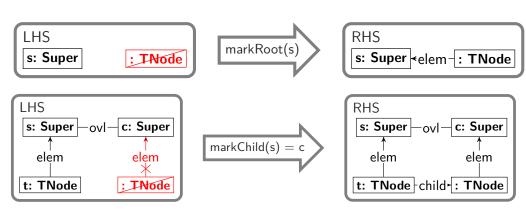
Node * usr 1 User name: String

Super link * Client

Child * elem*

TNode

Rules to mark root and children of marked nodes

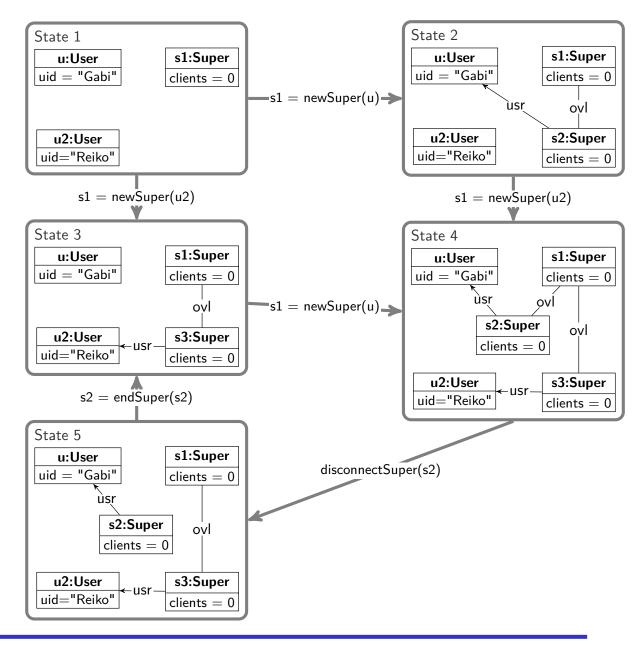


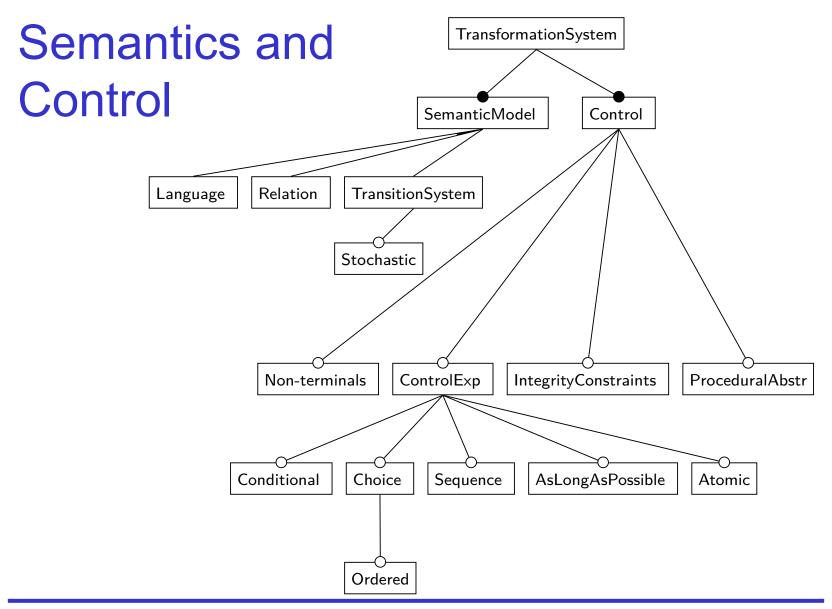
Initial set I: all unmarked instances graph of TG

Final set F: all graphs where all Super nodes are marked

P2P LTS (partial)

- Graphs as states
- Steps as transitions
- Rule names with args as labels





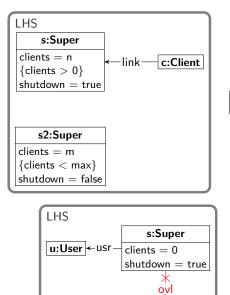
Rule-based programming:

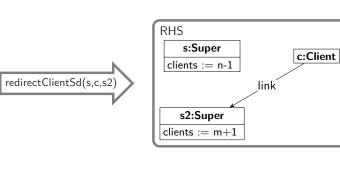
textual

Control expressions, Procedural abstraction, ...

```
LHS
                                                     RHS
                          startShutdown(s)
                                                          s:Super
     s:Super
shutdown = false
                                                      shutdown := true
 LHS
                                                    RHS
      s:Super
                                                    s:Super
  shutdown = true
                        disconnectSuperSd(s
        ovl
                                                    s1:Super
     s1:Super
```

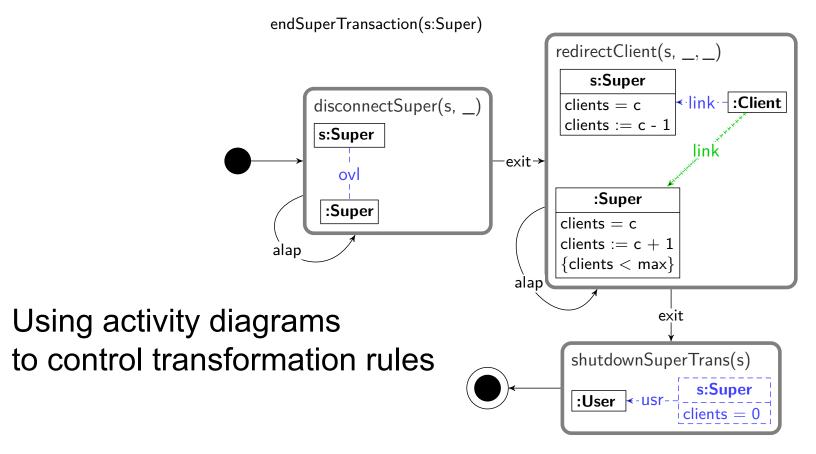
```
unit endSuper(s1: Super) {
   atomic {
     alap {
        disconnectSuper(s1,_)
     }
     alap {
        redirectClient (s1,__,_)
     }
     shutdownSuper(s1)
   }
}
```





s2:Super

Rule-based programming: diagrammatic





Part 1: Foundations

- Graphs
- ✓ Graph Transformations
- Semantics and Control
- Analysis
 - Properties: analysis problems
 - Techniques: solutions
 - Mapping problems to solutions
- Henshin Tool Demo

Properties

Languages

- membership
- inclusion
- instance generation
- non-ambiguity

Relations

- functionality (uniqueness)
- totality
- injectivity
- surjectivity
- correctness

Trans Systems

- reachability
- invariants
- deadlocks
- planning, optimisation
- temporal, prepost properties
- termination
- confluence

Techniques

Static

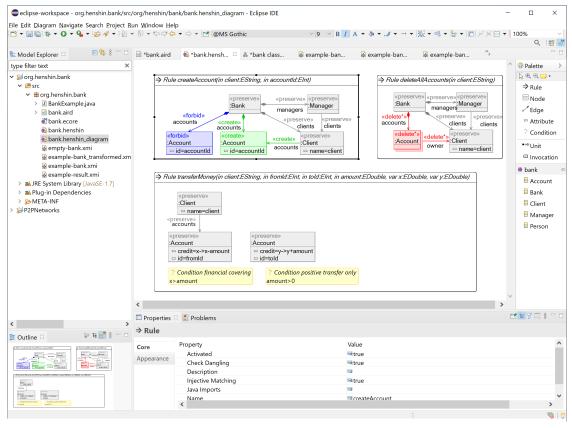
- Conflict and Dependency Analysis
- Termination Analysis
- Constraint Verification and Enforcement

Dynamic

- Model checking
- Graph parsing

	Conflict and dep-	Termination	Constraint	Model	Graph
	endency analysis	analysis	verification	checking	parsing
$\overline{Language}$					
Membership				X	X
Inclusion			\mathbf{X}	X	X
Instance generation	X		\mathbf{X}	X	
Non-ambiguity	X	X		X	
Relation					
Functional behaviour	X	X		X	X
Totality		X		X	
Injectivity	X	X		X	X
Surjectivity		X		X	
Correctness	X		X	X	
Transition system					
Reachability				X	X
Invariants			\mathbf{X}	X	
Deadlocks				X	
Planning &					
optimisation				X	
Temporal properties			\mathbf{X}	X	
Termination		X		X	
Confluence	X	X		X	_

Tools: Henshin



EMF model transformation

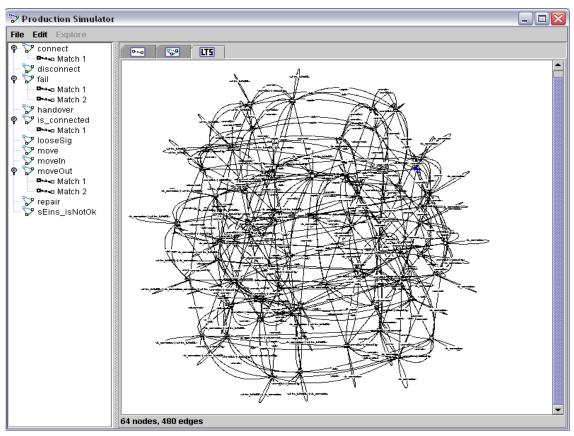
Conflict and dependency analysis

Constraint verification

Model checking by translation

https://www.eclipse.org/henshin/

Tools: Groove



Graph transformation

Visualisation of rules and matches

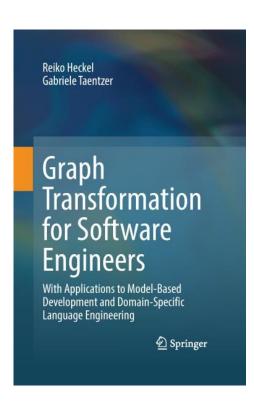
Native model checking in CTL and LTL

https://groove.ewi.utwente.nl



Part 1: Foundations

- ✓ Graphs
- ✓ Graph Transformations
- ✓ Semantics and Control
- Analysis



Further Reading



www.graph-transformation-for-softwareengineers.org

- Author copy
- Exercises
- Slides

Part II – Graph Transformation in Software Engineering

- Detecting Inconsistent Requirements in a Use-Case-Driven Approach
- 6. Service Specification and Matching
- Model-Based Testing
- Reverse Engineering: Inferring Visual Contracts from Java Programs
- Stochastic Analysis of Dynamic Software Architectures
- Advanced Modelling-Language Definition: Integrating Metamodelling with Graph Transformation
- 11. Improving Models and Understanding Model Changes
- 12. Translating and Synchronising Models

Part I – Foundations of Graph Transformation

- 1. Graphs for Modelling and Specification
- 2. Graph Transformation Concepts
- Beyond Individual Rules: Usage Scenarios and Control Structures
- 4. Analysis and Improvement of Graph Transformation Systems