

A Frog, A Mug, Gluing, and Logical Relations

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or, possibly, in the Wise words of the late Eric Morecombe:

... “but not necessarily,” in that order ...

AMP: “Exponentials in Heyting Algebra gluing are a bit of a mystery.”

The aim of this talk is to de-mystification.

- ▶ Brief review of applications of gluing.
- ▶ Artin Gluing of CCCs to *Set*.
- ▶ Joke.
- ▶ A version of Logical Relations Gluing. Relating AG and LRG.
- ▶ Frog.
- ▶ Gluing of HAs to $\mathcal{P}(X)$.
- ▶ Mug.
- ▶ A 2nd version of Logical Relations Gluing. Relating it to LRG.
- ▶ Conclusions.

Roughly speaking, a glued category $\mathcal{GL}(\Gamma)$ is one built from two highly structured categories, $\mathcal{C}$ and $\mathcal{D}$, and a functor $\Gamma: \mathcal{D} \rightarrow \mathcal{C}$ with few structure preserving properties. What's cool is that $\mathcal{GL}(\Gamma)$ often enjoys similar structure to $\mathcal{C}$ and $\mathcal{D}$. There are many applications

- ▶ Proving the Existence and Disjunction Properties of Intuitionistic Predicate Calculus.
- ▶ Proving “similar” properties in richly typed logics of quite different flavours.
- ▶ Proving Conservative Extension Results for Equational Type Theories.
- ▶ Proving results about Normalisation

Let $\Gamma: \mathcal{D} \rightarrow \mathbf{Set}$ be a functor where $\mathcal{D}$ is a ccc. Suppose also that Γ preserves finite products. We define a category $\mathcal{GL}(\Gamma)$, the “Artin gluing of $\mathcal{D}$ to $\mathbf{Set}$ along Γ ”. This has objects (X, f, D) where $X \in \mathbf{Set}$, $D \in \mathcal{D}$ and $f: X \rightarrow \Gamma D$ is a function. And morphisms

$$(\theta, \phi): (X, f, D) \rightarrow (X', f', D')$$

where

$$\begin{array}{ccc} X & \xrightarrow{f} & \Gamma D \\ \theta \downarrow & & \downarrow \Gamma \phi \\ X' & \xrightarrow{f'} & \Gamma D' \end{array}$$

The exponential $(X, f, D) \Rightarrow (X', f', D')$ is $(\boxed{P}, p_2, D \Rightarrow D')$ given by the pullback

$$\begin{array}{ccc}
 \boxed{(X \Rightarrow X') \times_{X \Rightarrow \Gamma D'} \Gamma(D \Rightarrow D')} & \xrightarrow{p_2} & \Gamma(D \Rightarrow D') \\
 \downarrow p_1 & & \downarrow \lambda(\Gamma ev \circ \cong \circ (id \times f)) \\
 X \Rightarrow X' & \xrightarrow{id \Rightarrow f'} & X \Rightarrow \Gamma D'
 \end{array}$$

where

$$\Gamma(D \Rightarrow D') \times \Gamma D \xrightarrow{\cong} \Gamma((D \Rightarrow D') \times D) \xrightarrow{\Gamma ev} \Gamma D'$$

... Joke ...

“Is **everything** you do professionally, related to glue?”

“Is **everything** you do professionally, related to glue?”

The answer I gave is “yes”.

“Is **everything** you do professionally, related to glue?”

If you find this tricky to believe, I'd better give a proof

“Is **everything** you do professionally, related to glue?”

Proof: I spend quite a bit of time writing papers about gluing. Thus it suffices to show that, when I'm not writing such papers, I'm doing something connected with glue.

“Is **everything** you do professionally, related to glue?”

For the most part, when not writing about glue, I'm **stuck** (on a proof). Q.E.D.

We define a category $\mathcal{Gl}(\Gamma)$, the “logical relations gluing of $\mathcal{D}$ to $\mathcal{Set}$ along Γ ”.

- The objects of $\mathcal{Gl}(\Gamma)$ are triples $(X, \triangleleft, D)$ where

$$\triangleleft \subseteq X \times \Gamma D$$

- A morphism $(\theta, \phi): (X, \triangleleft, D) \rightarrow (X', \triangleleft', D')$ is given by function $\theta: X \rightarrow X'$ and a morphism $\phi: D \rightarrow D'$ in $\mathcal{D}$ for which given $x \in X$ and $d \in \Gamma D$

$$x \triangleleft d \quad \implies \quad \theta x \triangleleft' (\Gamma \phi) d$$

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The exponential $(X, \triangleleft, D) \Rightarrow (X', \triangleleft', D')$ is

$$(X \Rightarrow X', \triangleleft \Rightarrow \triangleleft', D \Rightarrow D')$$

where given $f \in X \Rightarrow X'$ and $F \in \Gamma(D \Rightarrow D')$

$$f(\triangleleft \Rightarrow \triangleleft')F \Longleftrightarrow$$

$$(\forall x \in X) (\forall d \in \Gamma D)$$

$$(x \triangleleft d \Longrightarrow fx \triangleleft' ((\Gamma ev) \circ \cong)(F, d))$$

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Notice that there is an inclusion functor

$$I: \mathcal{GL}(\Gamma) \rightarrow \mathcal{Gl}(\Gamma)$$

since given any (X, f, D) with $f: X \rightarrow \Gamma D$, then trivially $f \subseteq X \times \Gamma D$.

Taking (X, f, D) and (X', f', D') in $\mathcal{GL}(\Gamma)$, the exponential in $\mathcal{Gl}(\Gamma)$ is

$$(X \Rightarrow X', f \Rightarrow f', D \Rightarrow D')$$

where $f \Rightarrow f' \subseteq (X \Rightarrow X') \times \Gamma(D \Rightarrow D')$, and in fact it is rather nice that

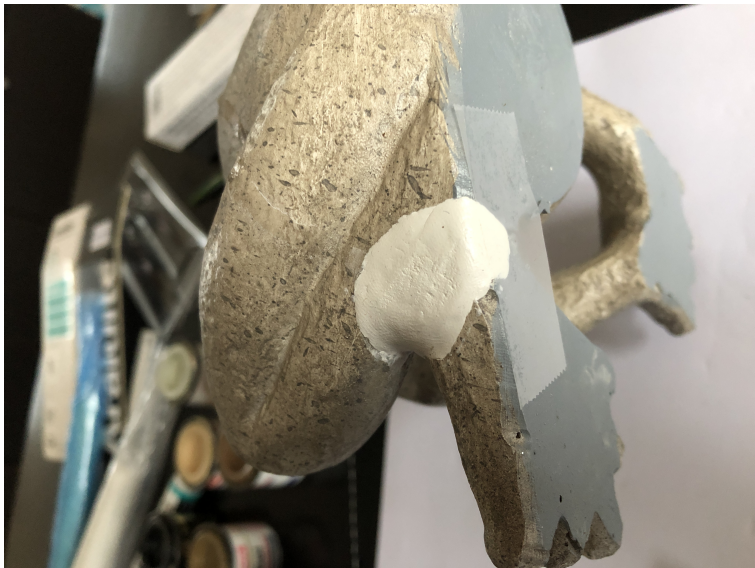
$$f \Rightarrow f' = P \stackrel{\text{def}}{=} (X \Rightarrow X') \times_{X \Rightarrow \Gamma D'} \Gamma(D \Rightarrow D')$$

...Frog ...

Broken Frog



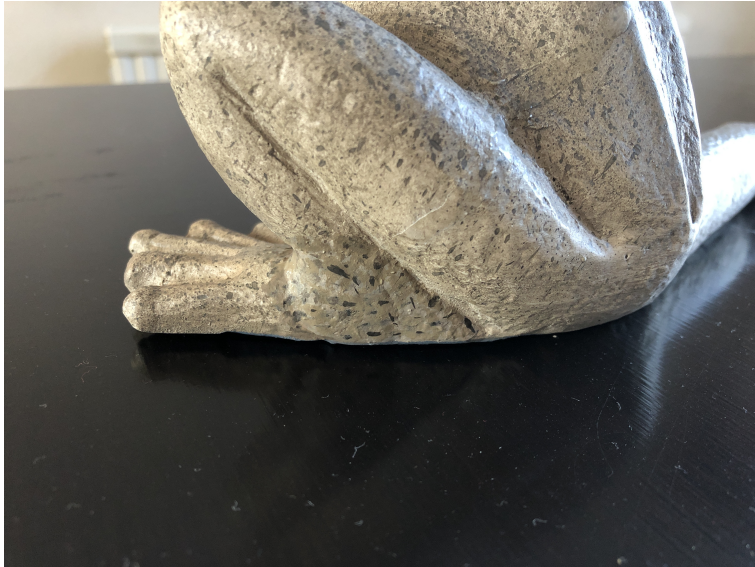
Epoxy Frog



Clappy Glued Frog



Happy Glued Frog



Let $\gamma: H \rightarrow \mathcal{P}(X)$ be a functor (monotone function) where H is a Heyting Algebra. Suppose also that γ preserves finite products (meets).

We define a poset $\mathcal{GL}(\gamma)$, the “gluing of H to $\mathbf{Set}$ along γ ”. This has elements (k, h) where $k \in \mathcal{P}(X)$, $h \in H$ and $k \leq_{\mathcal{P}(X)} \gamma h$. And order instances

$$(k, h) \leq (k', h')$$

just in case

$$\begin{array}{ccc} k & \leq_{\mathcal{P}(X)} & \gamma h \\ \downarrow & & \downarrow \\ \leq_{\mathcal{P}(X)} & & \leq_{\mathcal{P}(X)} \\ \downarrow & & \downarrow \\ k' & \leq'_{\mathcal{P}(X)} & \gamma h' \end{array}$$

The exponential $(k, h) \Rightarrow (k', h')$ is

$$((k \Rightarrow k') \wedge \gamma(h \Rightarrow h'), h \Rightarrow h')$$

Don't forget that $k \leq_{\mathcal{P}(X)} \gamma h$ is simply the subset $k \subseteq \gamma h$.

And that in $\mathcal{P}(X)$, the exponential $k \Rightarrow k'$ is $(X \setminus k) \cup k'$.

The meet is of course $\cap$ in $\mathcal{P}(X)$.

...Mug ...

Broken Mug



Gluing Mug



Fixed Mug



I always thought of the (k, h) as being pairs satisfying a property, namely $k \subseteq \gamma h$.

I spent a little while thinking about a generalisation with objects (k, Φ, h) where $\Phi(k, h)$ is a property.

But the right thing to do is to note that we have an inclusion *function*
 $\iota: k \rightarrow \gamma h$.

And further that this is the *relation* $\Delta_k: k \rightarrow \gamma h$ where Δ_k is the diagonal on k .

We define a category $\mathcal{Gl}(\gamma)$, the “logical relations gluing of H to Set along γ ”.

- ▶ The objects of $\mathcal{Gl}(\gamma)$ are triples $(k, \triangleleft, h)$ where $\triangleleft \subseteq k \times \gamma h$.
- ▶ The order $(k, \triangleleft, h) \leq (k', \triangleleft', h')$ is pointwise. The pointwise order on the first and third components makes the order on the second components type-correct.

The exponential $(k, \triangleleft, h) \Rightarrow (k', \triangleleft', h')$ is

$$(k \Rightarrow k', \triangleleft \Rightarrow \triangleleft', h \Rightarrow h')$$

where given $x \in k \Rightarrow k'$ and $y \in \gamma(h \Rightarrow h')$

$$x(\triangleleft \Rightarrow \triangleleft')y \quad \Longleftrightarrow \quad (x \triangleleft y \quad \Longrightarrow \quad x \triangleleft' y)$$

Notice that there is an inclusion functor

$$I: \mathcal{GL}(\gamma) \rightarrow \mathcal{Gl}(\gamma)$$

since given any (k, h) with $k \leq \gamma h$, then trivially $\Delta_k \subseteq k \times \gamma h$. So $I(k, h) \stackrel{\text{def}}{=} (k, \Delta_k, h)$.

Taking (k, h) and (k', h') in $\mathcal{GL}(\Gamma)$, the exponential (of the images) in $\mathcal{Gl}(\gamma)$ is

$$(k \Rightarrow k', \Delta_k \Rightarrow \Delta_{k'}, h \Rightarrow h')$$

where $\Delta_k \Rightarrow \Delta_{k'} \subseteq (k \Rightarrow k') \times \gamma(h \Rightarrow h')$ and in fact I think it is definitely rather nice that

$$\Delta_k \Rightarrow \Delta_{k'} = (k \Rightarrow k') \wedge \gamma(h \Rightarrow h')$$

- ▶ Artin exponential $\cong$ logical relation: appears in my Applied Categorical Structures 1996 paper, *On Fixpoint Objects and Gluing Constructions* (†).
- ▶ There are many “gluing techniques” results. While they have overlaps, to me they often feel disconnected.
- ▶ It would be good to write a paper that presents the theory of gluing in a single framework.
- ▶ How general can the target categories for Γ be while still facilitating logical relations?
 - ▶ In (†) it is $\omega\mathbf{CPO}^c$.